Slacker Outline

Background

- Containers: lightweight isolation
- Docker: file-system provisioning

Container Workloads

Default Driver: AUFS

Our Driver: Slacker

Evaluation

Conclusion

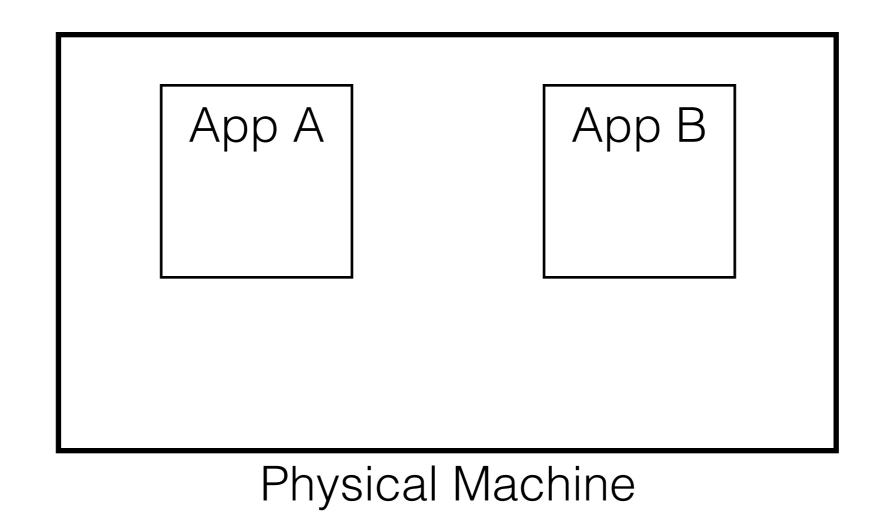
Why use containers?

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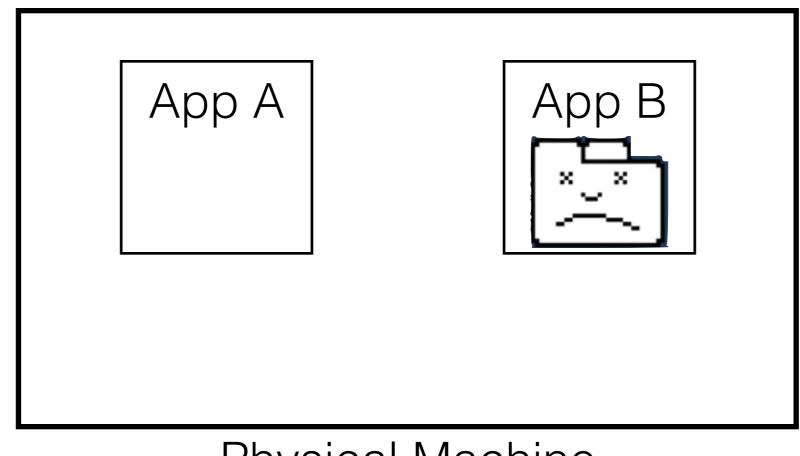
(it's trendy)

Why use containers?

(it's trendy)
(efficient solution to classic problem)

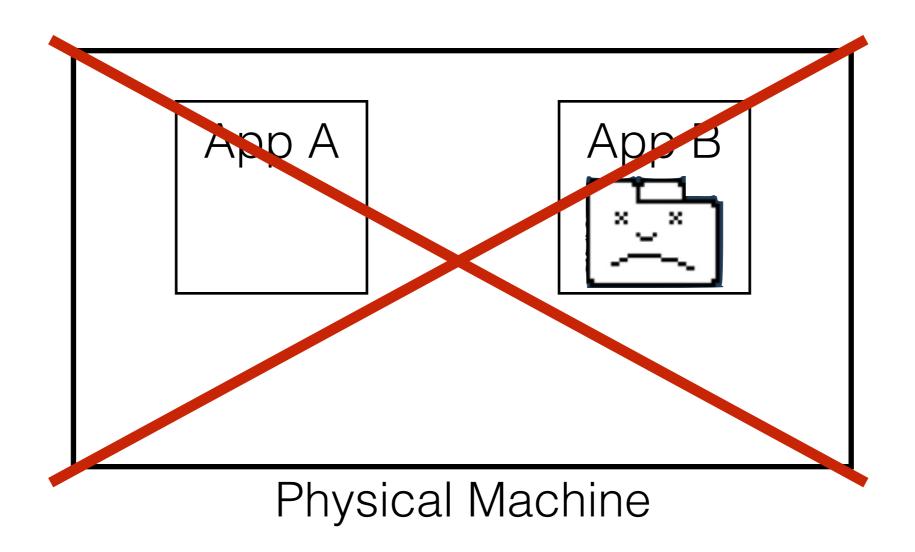


want: multitenancy

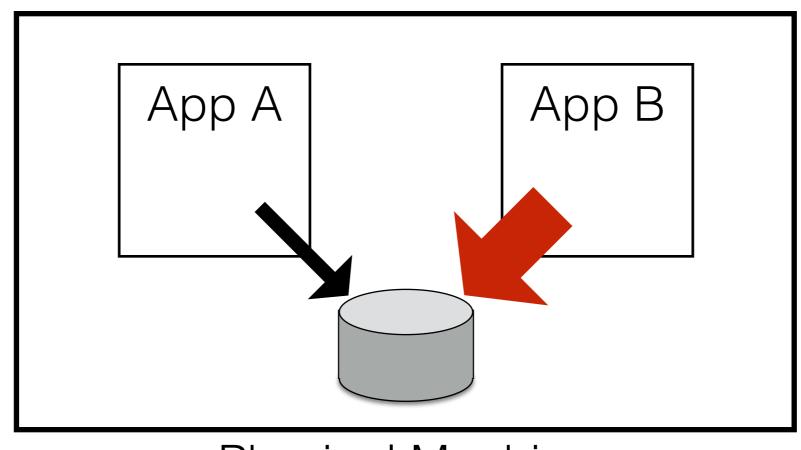


Physical Machine

don't want: crashes

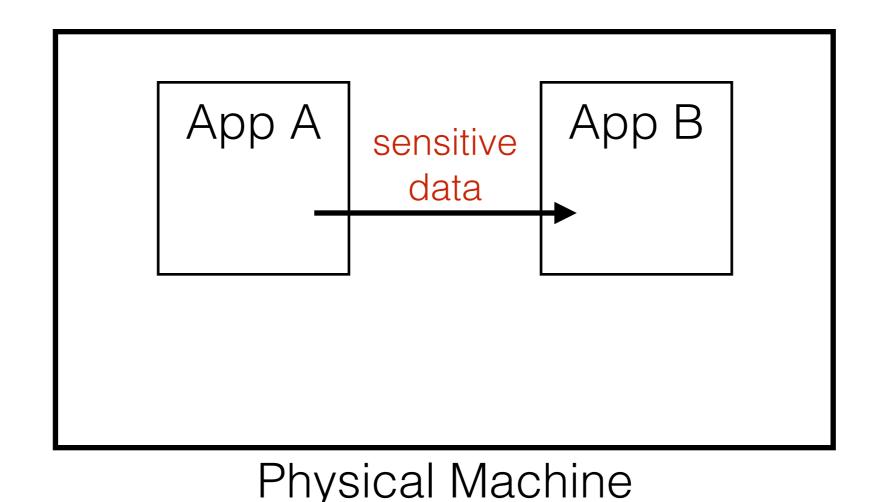


don't want: crashes



Physical Machine

don't want: unfairness



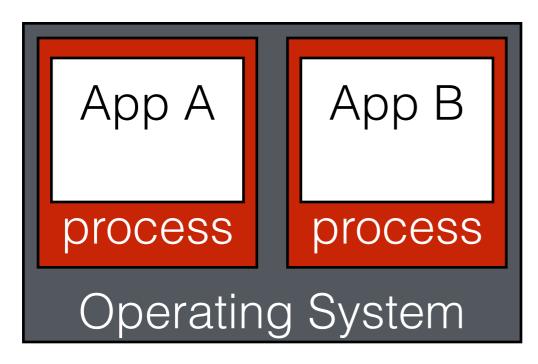
don't want: leaks

Solution: Virtualization

namespaces and scheduling provide illusion of private resources

1st generation: process virtualization

- isolate within OS (e.g., virtual memory)
- fast, but incomplete (missing ports, file system, etc.)



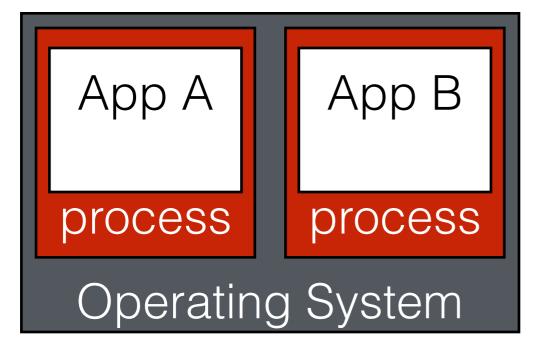
process virtualization

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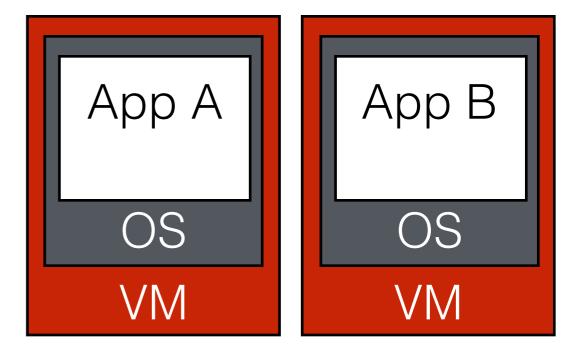
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2nd generation: machine virtualization

- isolate around OS
- complete, but slow (redundancy, emulation)



process virtualization



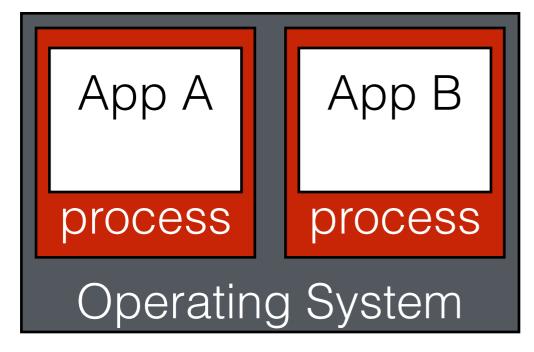
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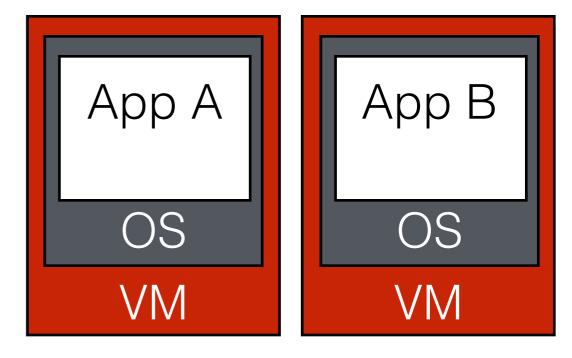
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3rd generation: container virtualization

- extend process virtualization: ports, file system, etc.
- fast and complete

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3rd generation: container virtualization

- extend process virtualization: ports, file system, etc.
- fast and complete????

many storage challenges

New Storage Challenges

Crash isolation

Physical Disentanglement in a Container-Based File System.

Lanyue Lu, Yupu Zhang, Thanh Do, Samer Al-Kiswany, Andrea C. Arpaci-Dusseau, Remzi H. Arpaci-Dusseau. **OSDI '14.**

Performance isolation

Split-level I/O Scheduling For Virtualized Environments.

Suli Yang, Tyler Harter, Nishant Agrawal, Salini Selvaraj Kowsalya, Anand Krishnamurthy, Samer Al-Kiswany, Andrea C. Arpaci-Dusseau, Remzi H. Arpaci-Dusseau. **SOSP '15.**

File-system provisioning

Slacker: Fast Distribution with Lazy Docker Containers.

Tyler Harter, Brandon Salmon, Rose Liu,

Andrea C. Arpaci-Dusseau, Remzi H. Arpaci-Dusseau. FAST '16.

today

Slacker Outline

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Docker Background

Deployment tool built on containers

An application is defined by a file-system image

- application binary
- shared libraries
- etc.

Version-control model

- extend images by committing additional files
- deploy applications by pushing/pulling images

Containers as Repos

LAMP stack example

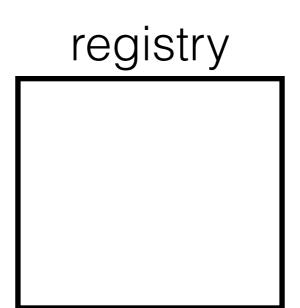
- commit 1: Linux packages (e.g., Ubuntu)
- commit 2: Apache
- commit 3: MySQL
- commit 4: PHP

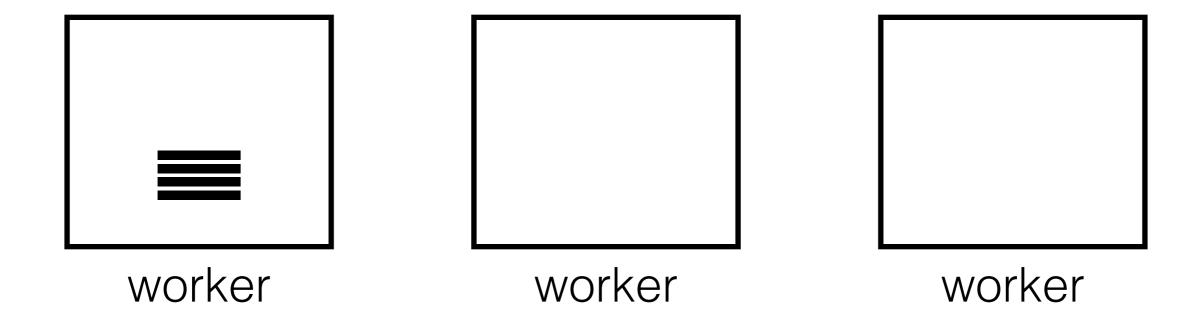
Central registries

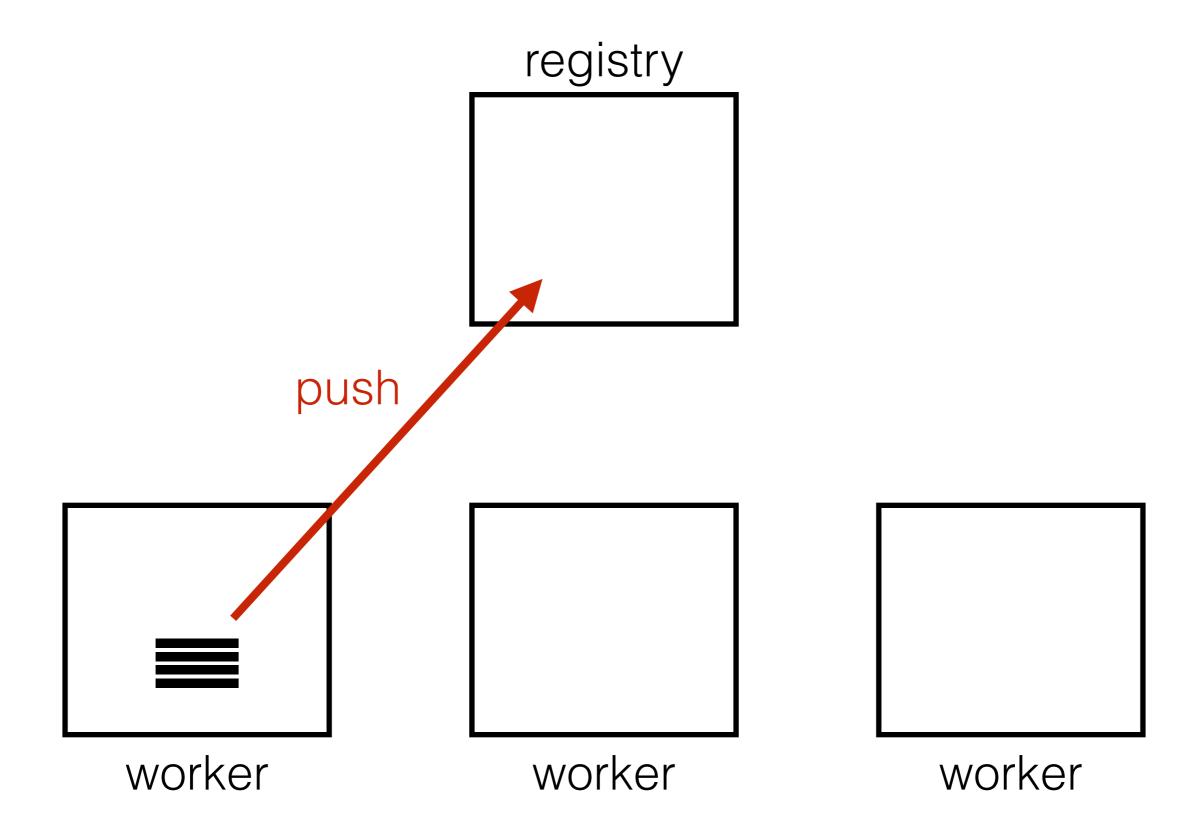
- Docker HUB
- private registries

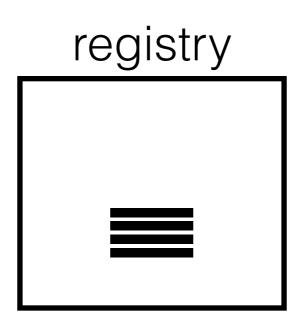
Docker "layer"

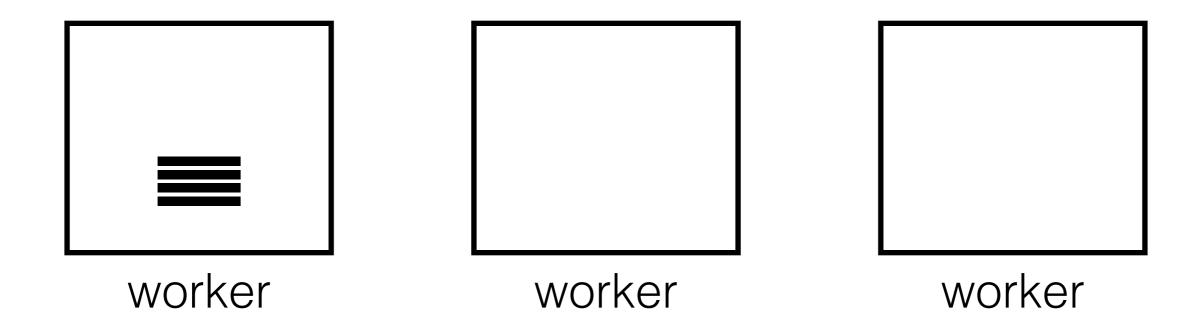
- commit
- container scratch space

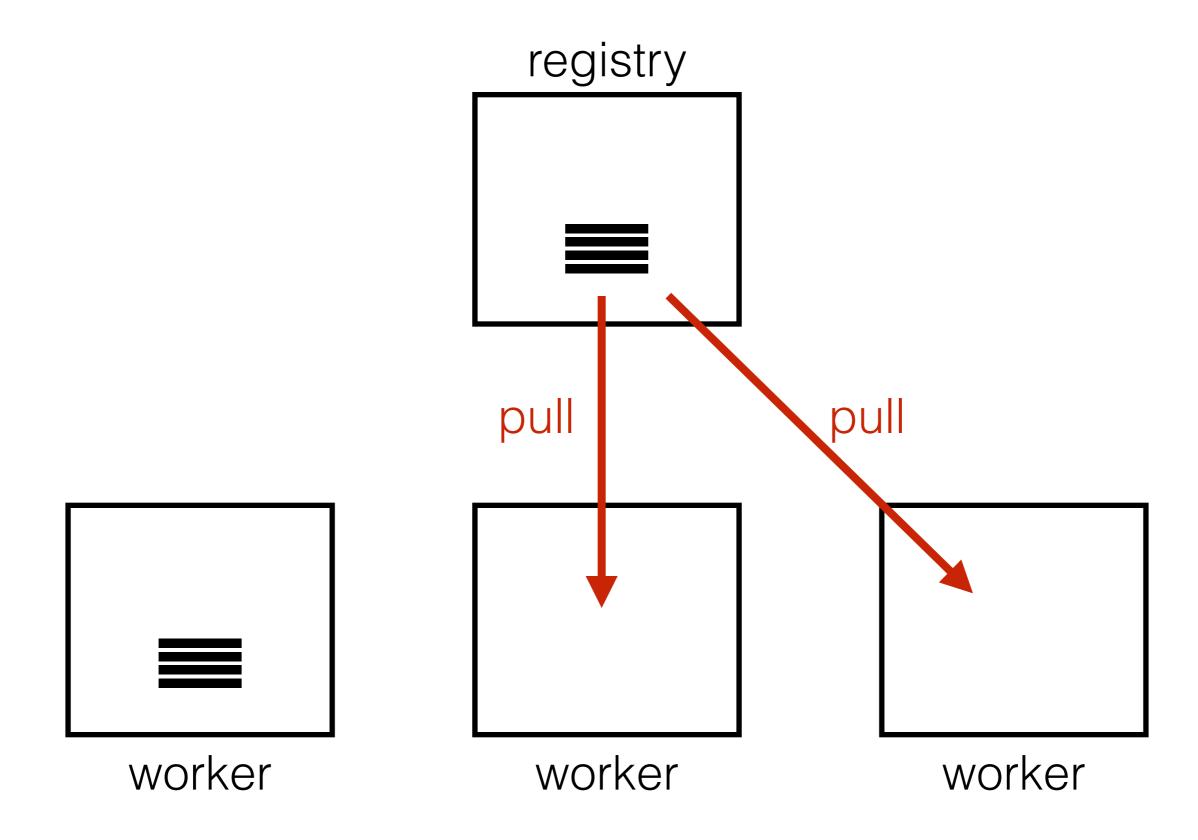


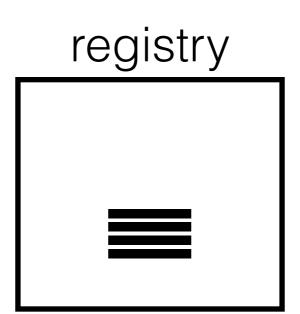


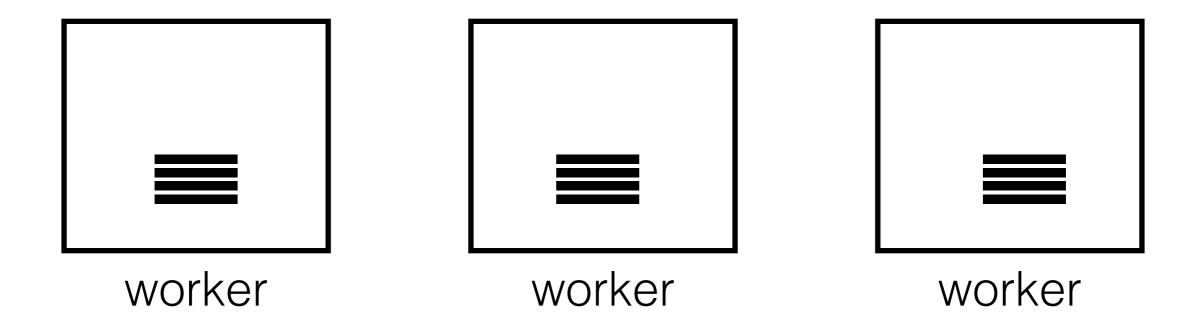


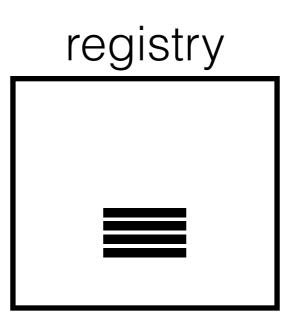


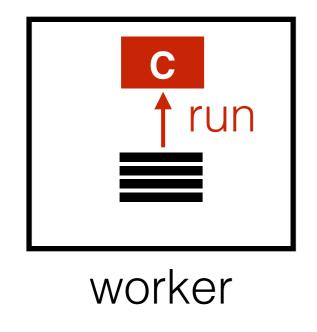


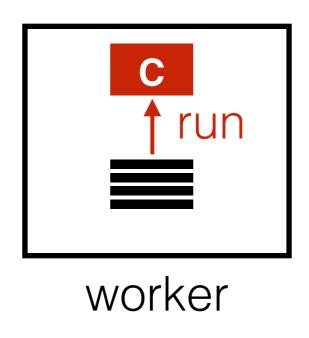


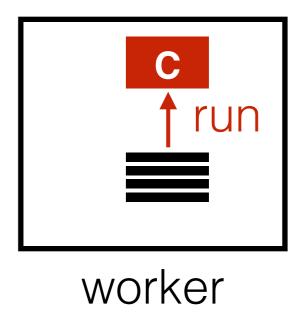




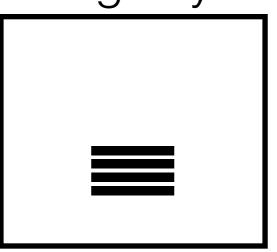




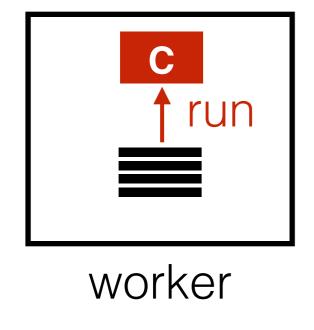


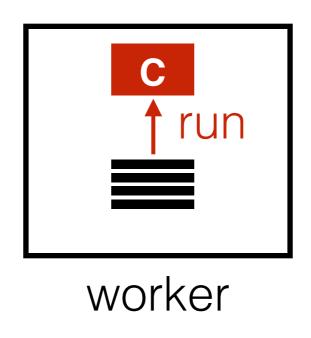


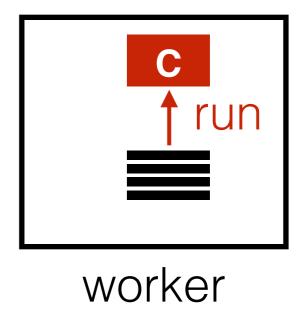




need a new benchmark to measure Docker push, pull, and run operations.







Slacker Outline

Background

Container Workloads

- HelloBench
- Analysis

Default Driver: AUFS

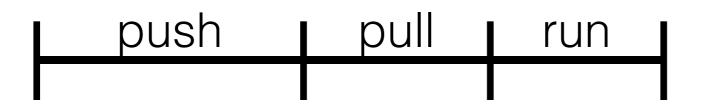
Our Driver: Slacker

Evaluation

Conclusion

Goal: stress container startup

- including push/pull
- 57 container images from Docker HUB
- run simple "hello world"-like task
- wait until it's done/ready

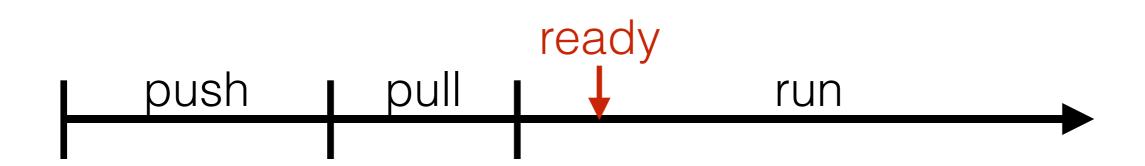


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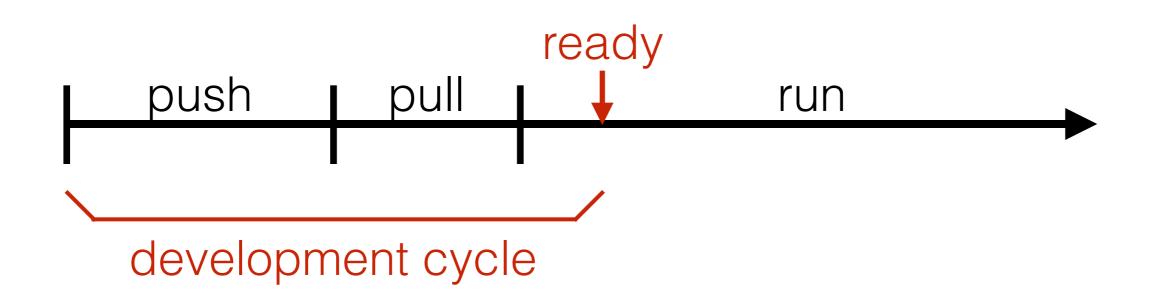


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Development cycle

distributed programming/testing



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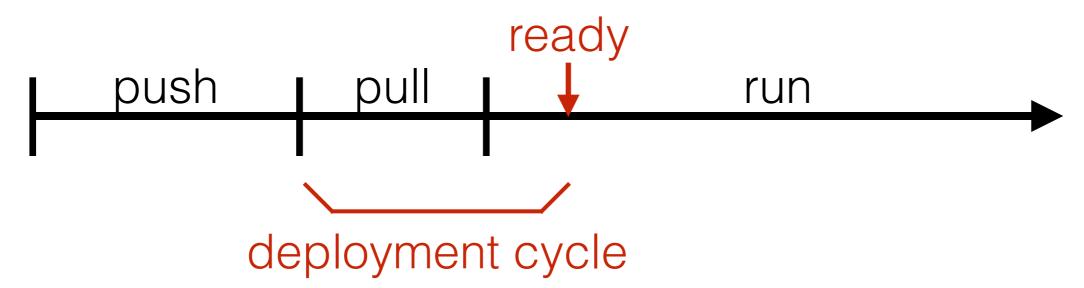
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- run simple "hello world"-like task
- wait until it's done/ready

Development cycle

distributed programming/testing

Deployment cycle

flash crowds, rebalance



Workload Categories

Language

clojure

gcc

golang

haskell

hylang

java

jruby

julia

mono

perl

php

pypy

python

r-base

rakudo-star

ruby

thrift

Linux Distro

alpine

busybox

centos

cirros

crux

debian

fedora

mageia

opensuse

oraclelinux

ubuntu

ubuntu-

debootstrap

ubuntu-upstart

Database

cassandra

crate

elasticsearch

mariadb

mongo

mysql

percona

postgres

redis

rethinkdb

Web Framework

django

iojs

node

rails

Web Server

glassfish

httpd

jetty

nginx

php-zendserver

tomcat

Other

drupal

ghost

hello-world

jenkins

rabbitmq

registry

sonarqube

Slacker Outline

Background

Container Workloads

- HelloBench
- Analysis

Default Driver: AUFS

Our Driver: Slacker

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Questions

How is data distributed across Docker layers?

How much image data is needed for container startup?

How similar are reads between runs?

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HelloBench images

• circle: commit

red: image

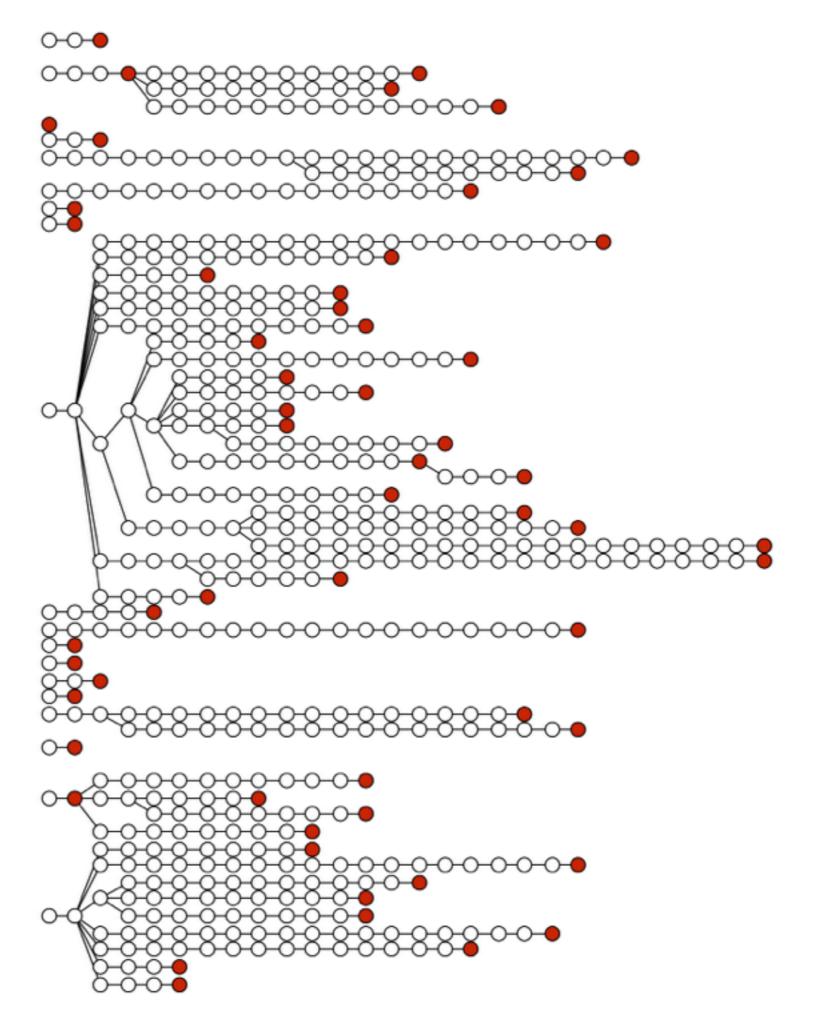


Image Data Depth

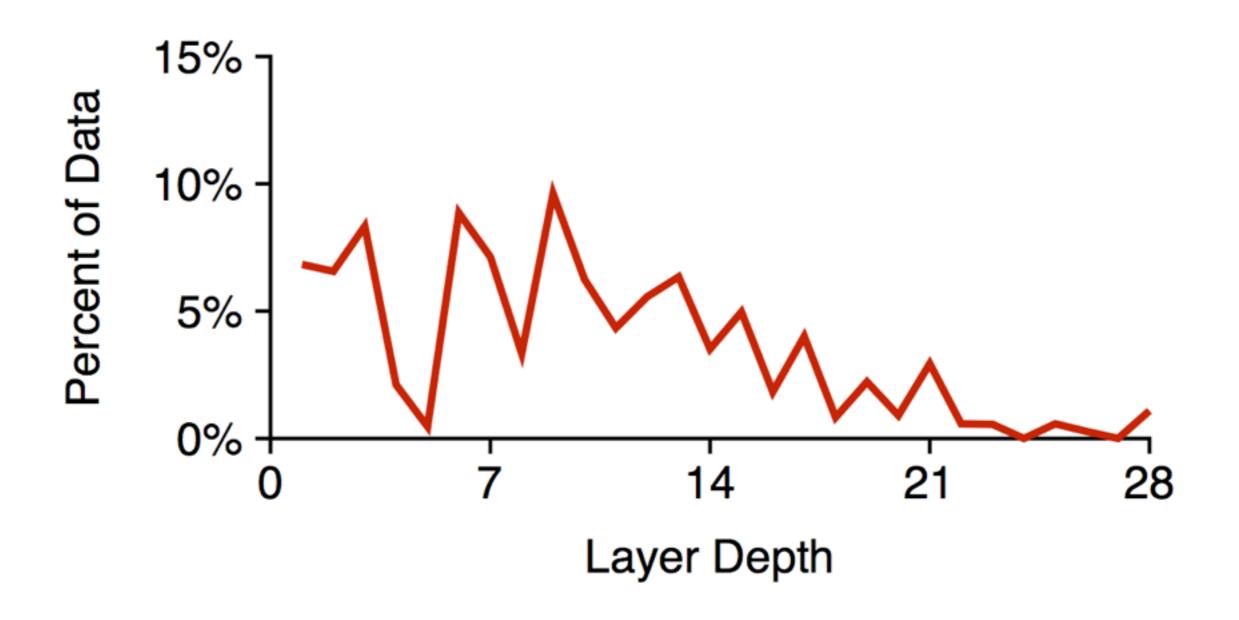
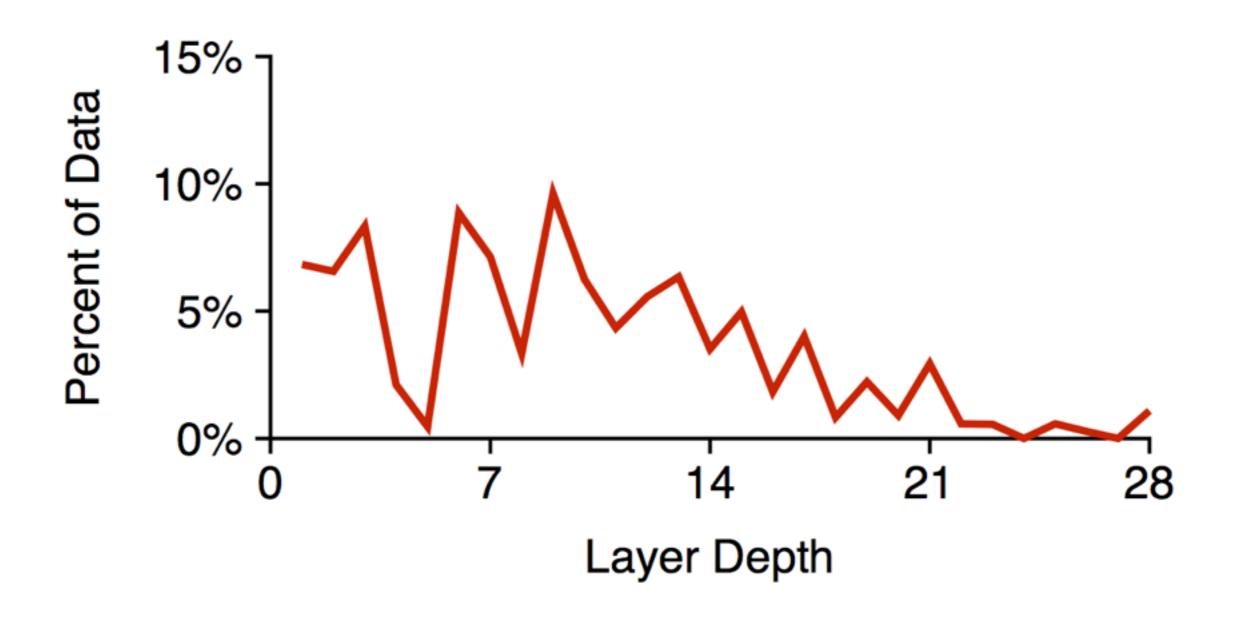


Image Data Depth



half of data is at depth 9+

Questions

How is data distributed across Docker layers?

- half of data is at depth 9+
- design implication: flatten layers at runtime

How much image data is needed for container startup?

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Questions

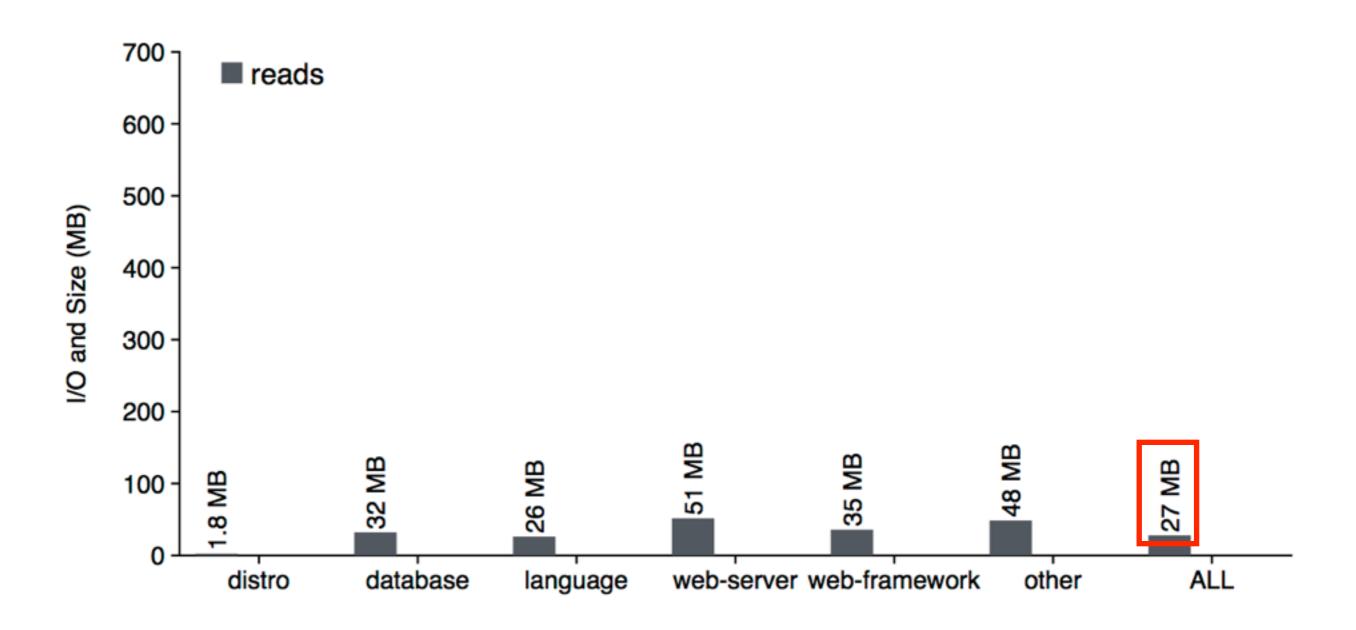
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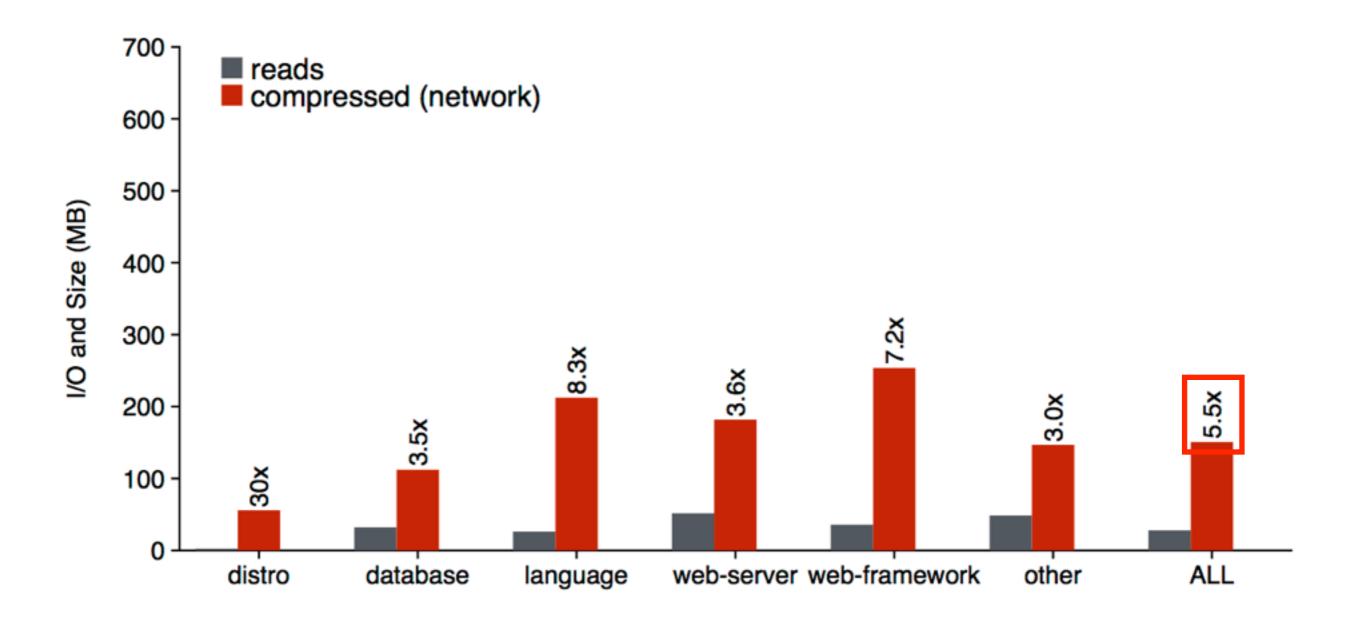
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How similar are reads between runs?

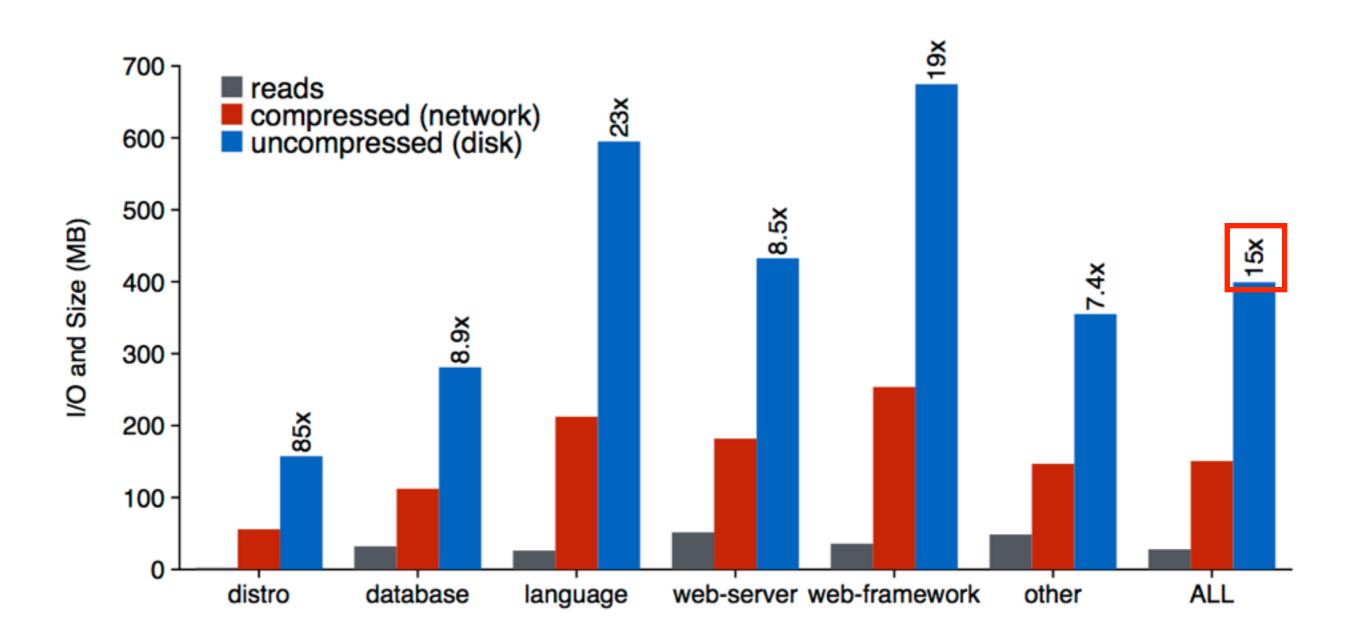
Container Amplification



Container Amplification



Container Amplification



only 6.4% of data needed during startup

Questions

How is data distributed across Docker layers?

- half of data is at depth 9+
- design implication: flatten layers at runtime

How much image data is needed for container startup?

- 6.4% of data is needed
- design implication: lazily fetch data

How similar are reads between runs?

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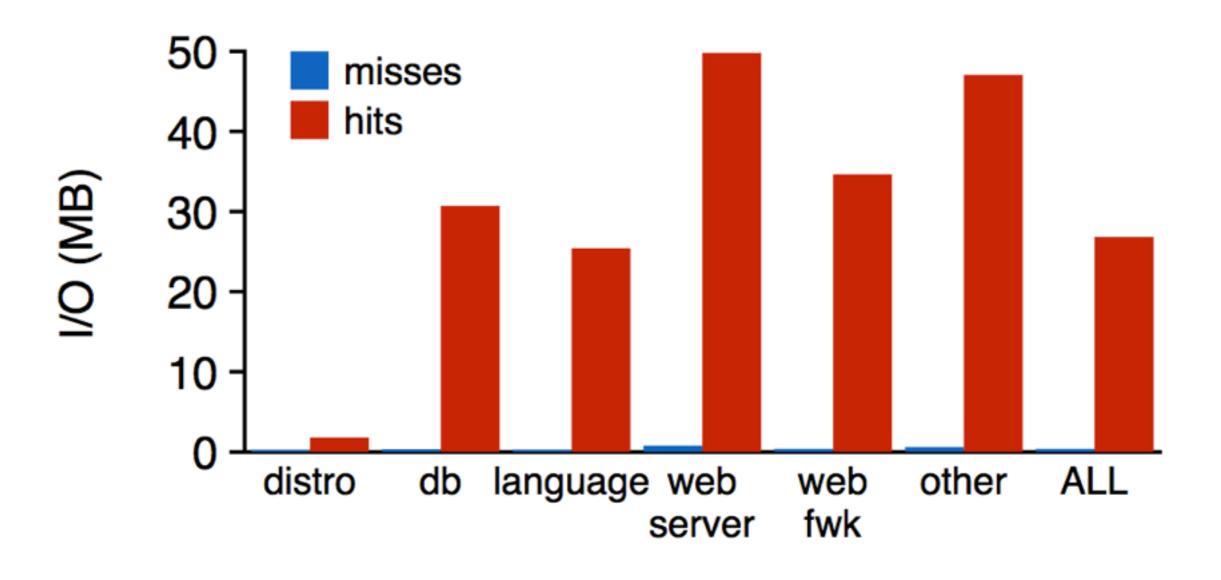
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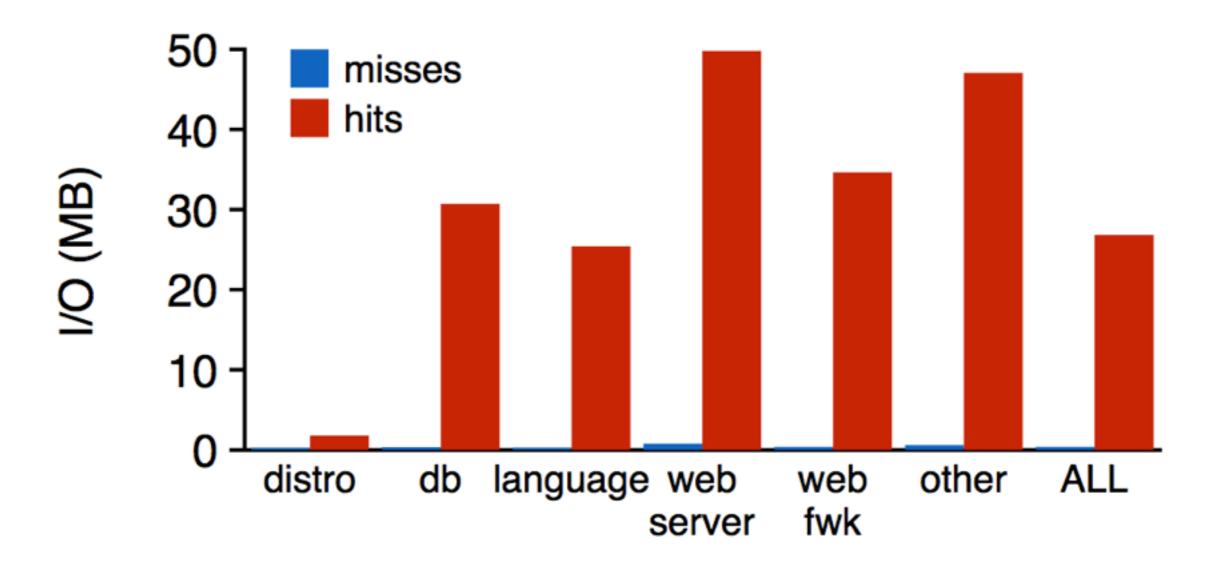
Repeat Runs

measure hits/misses for second of two runs



Repeat Runs

measure hits/misses for second of two runs



up to 99% of reads could be serviced by a cache

Questions

How is data distributed across Docker layers?

- half of data is at depth 9+
- design implication: flatten layers at runtime

How much image data is needed for container startup?

- 6.4% of data is needed
- design implication: lazily fetch data

How similar are reads between runs?

- containers from same image have similar read patterns
- design implication: share cache state between containers

Slacker Outline

Background

Container Workloads

Default Driver: AUFS

- Design
- Performance

Our Driver: Slacker

Evaluation

Conclusion

- stores data in an underlying FS (e.g., ext4)
- layer ⇒ directory in underlying FS
- root FS ⇒ union of layer directories

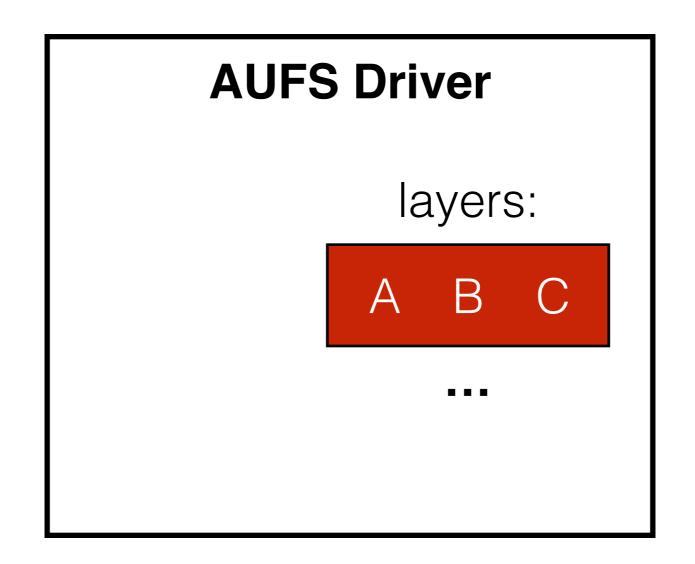
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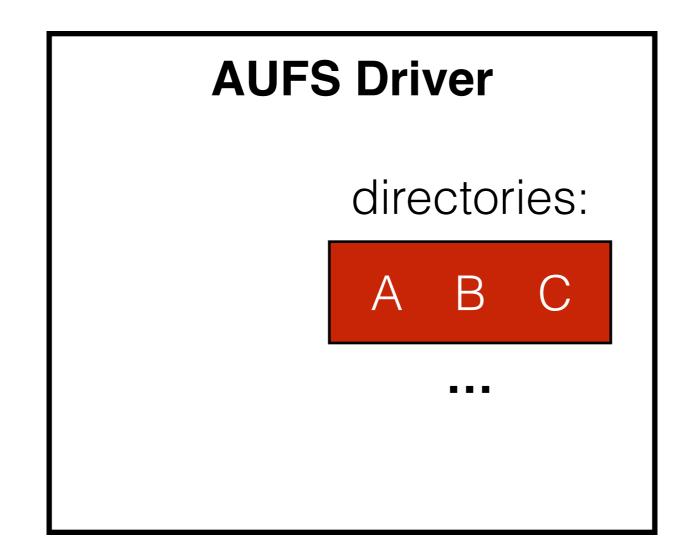
Operations

- push
- pull
- run

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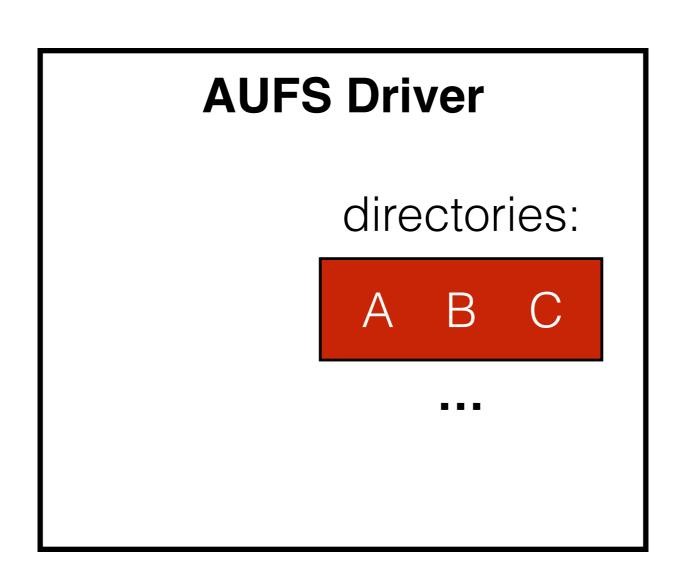
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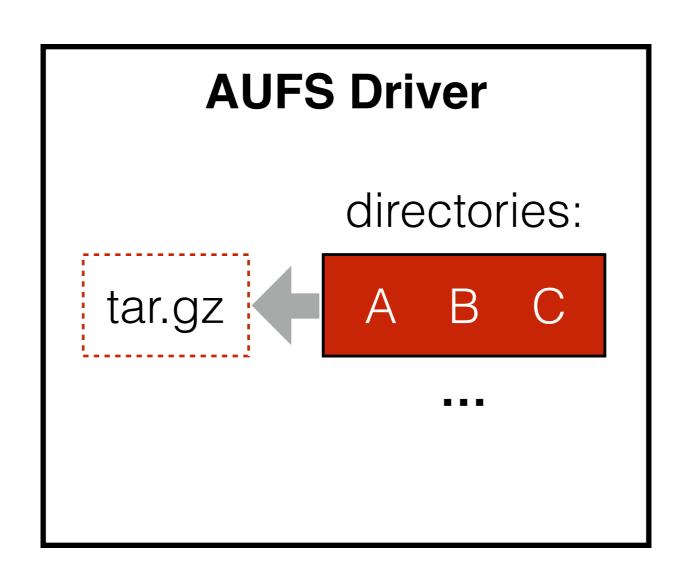
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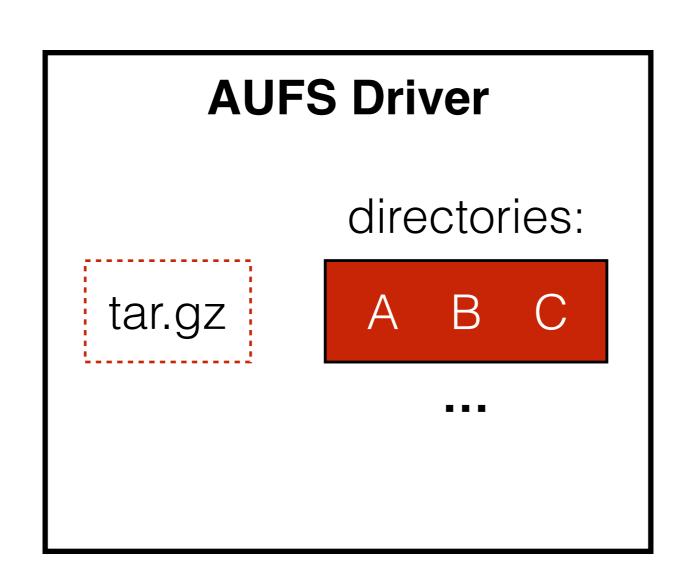
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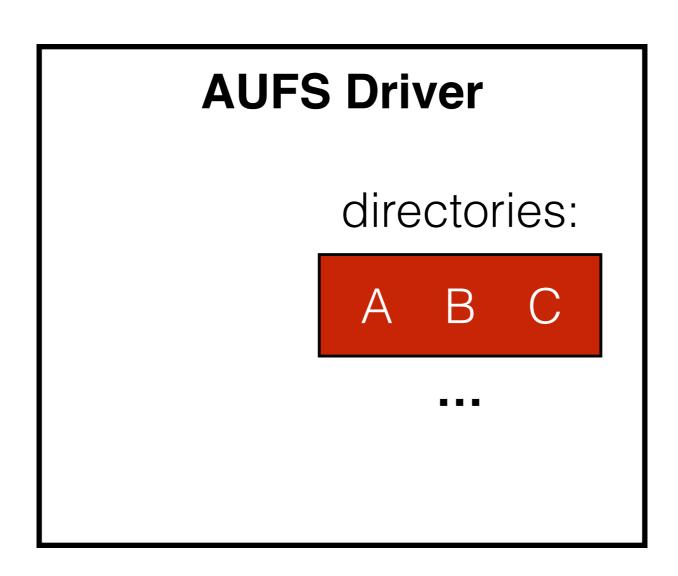
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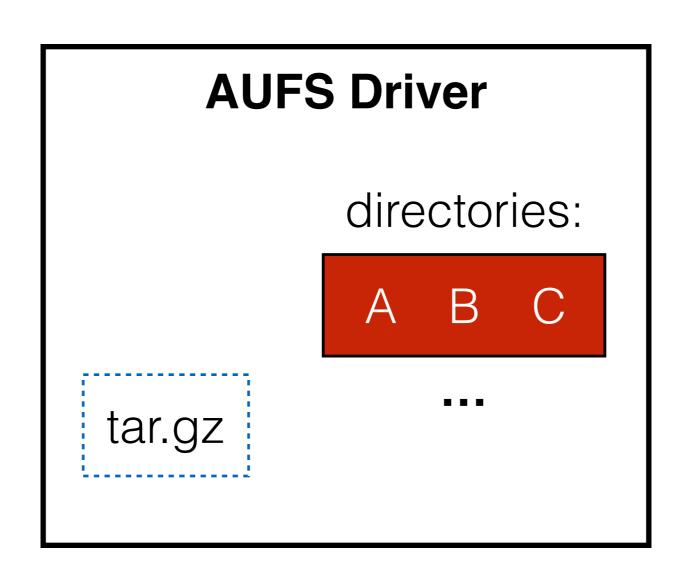
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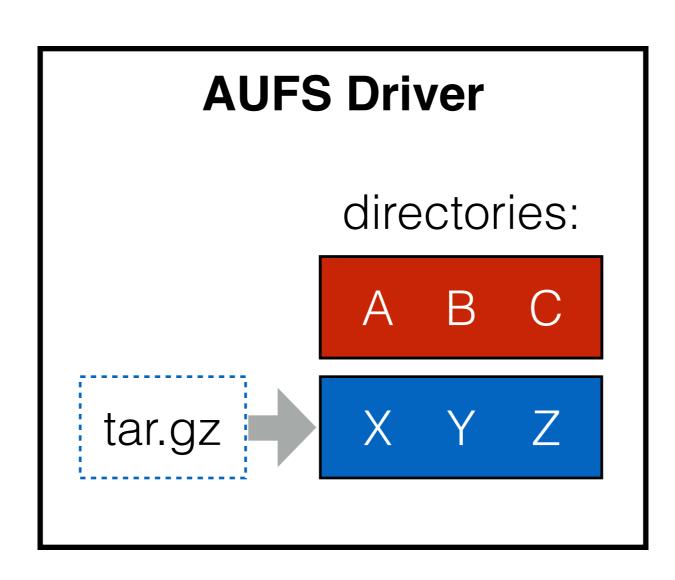
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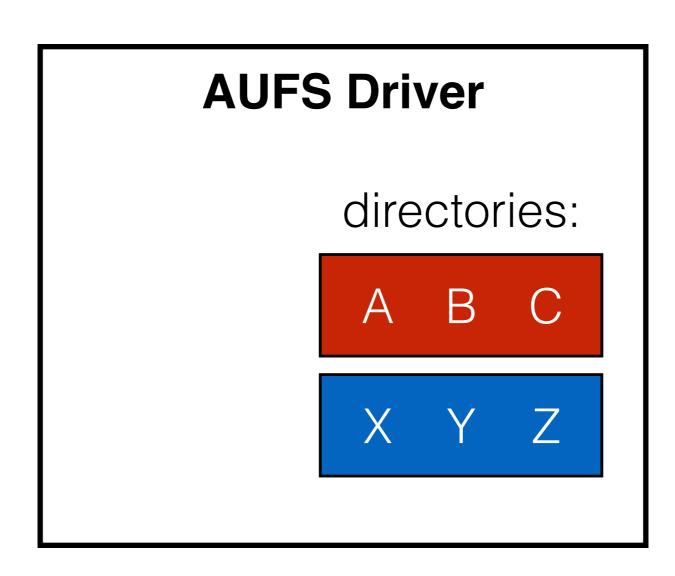
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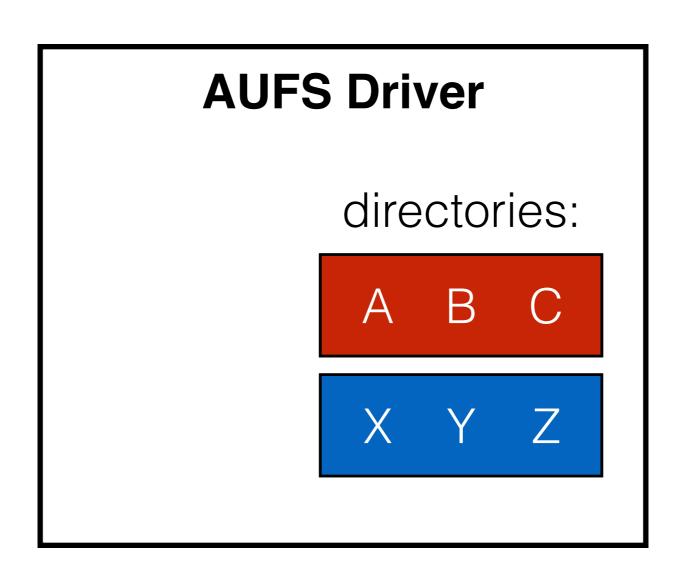
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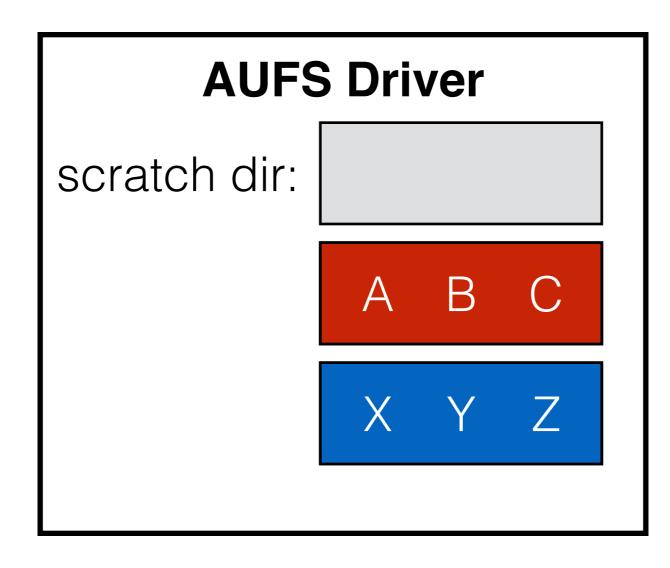
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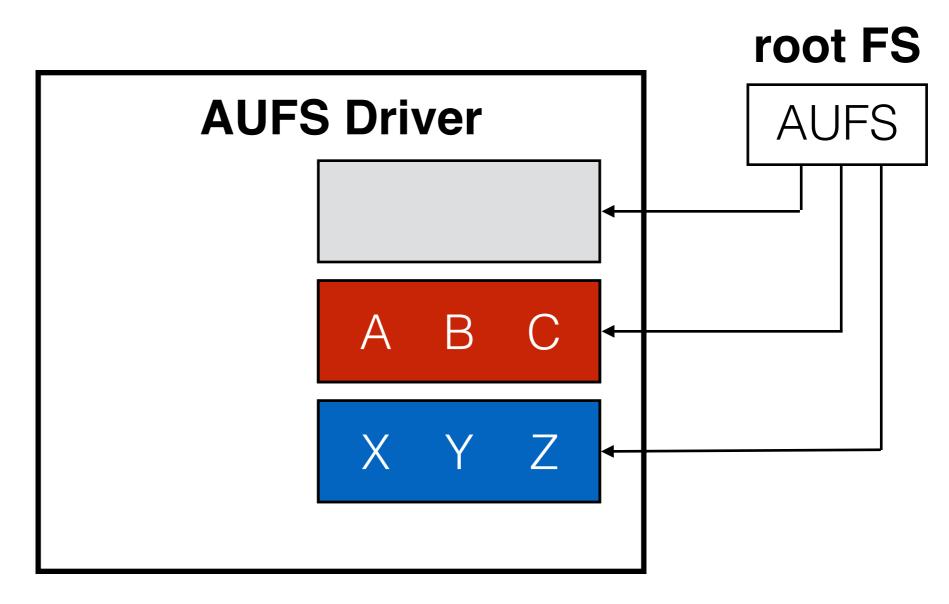
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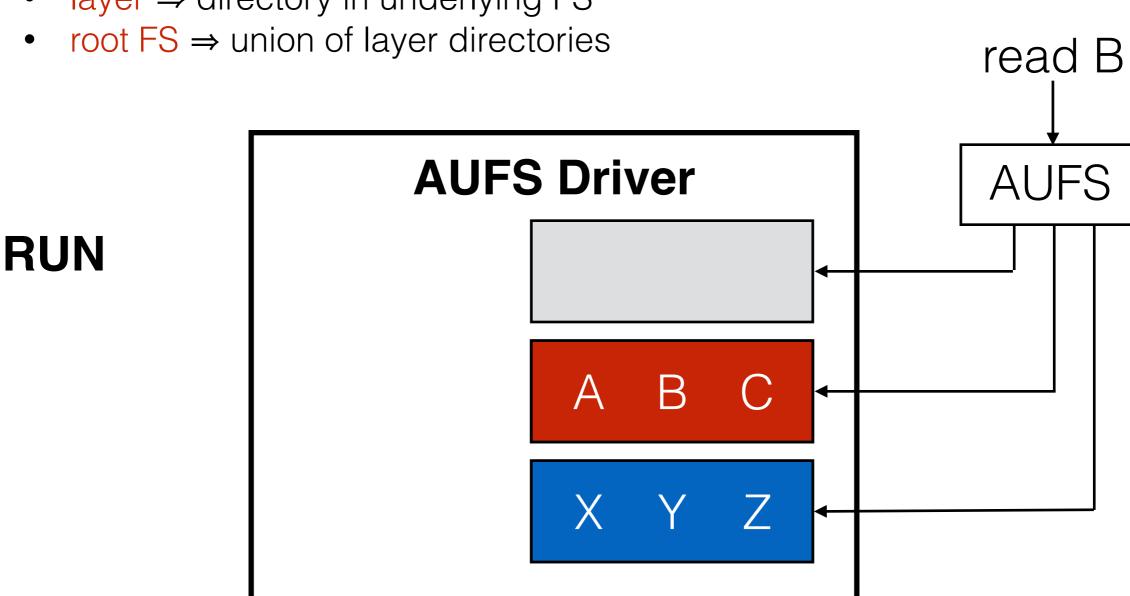


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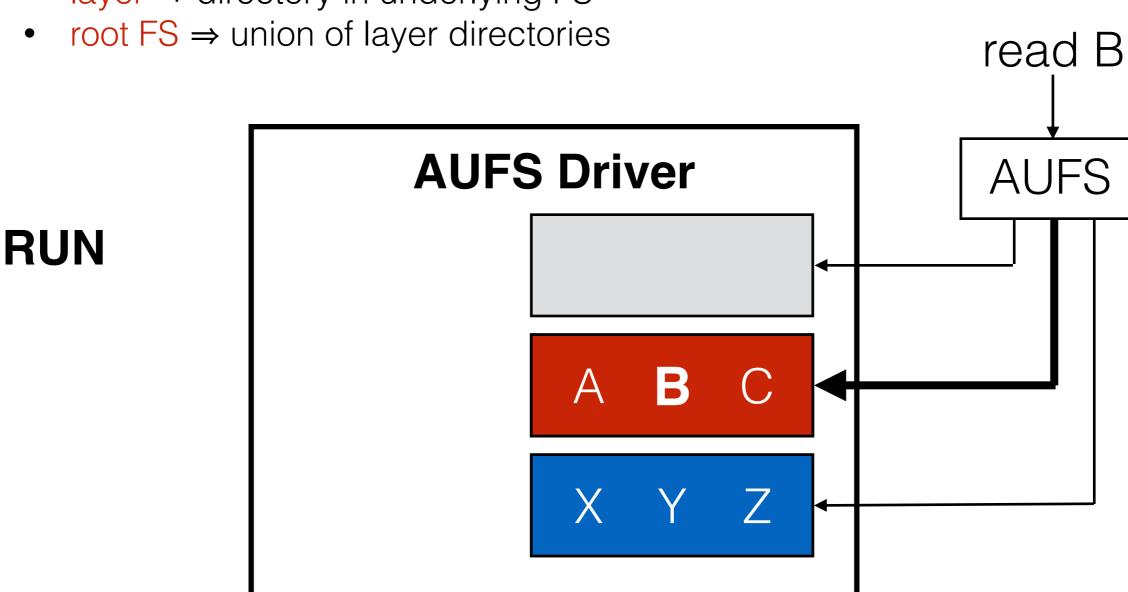
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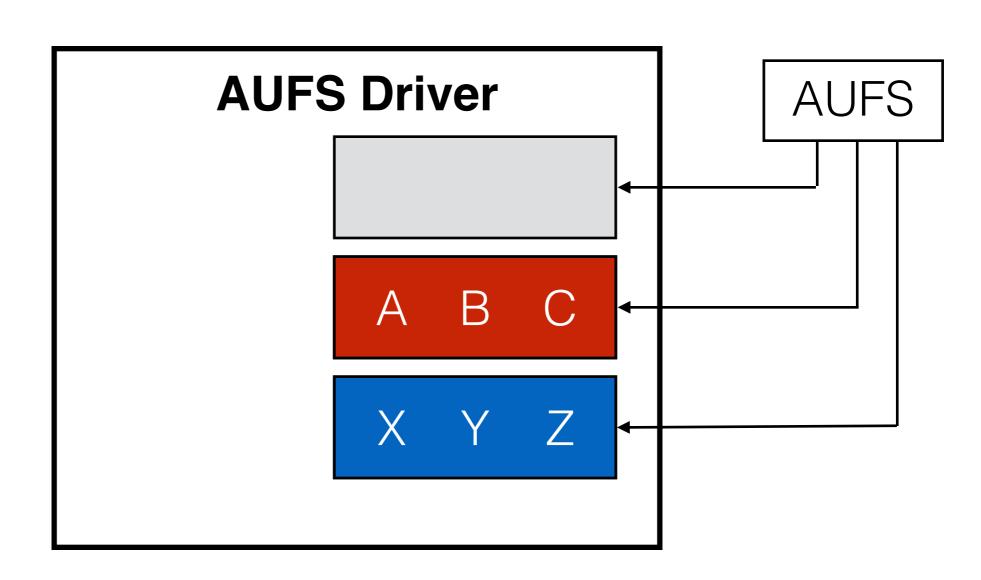


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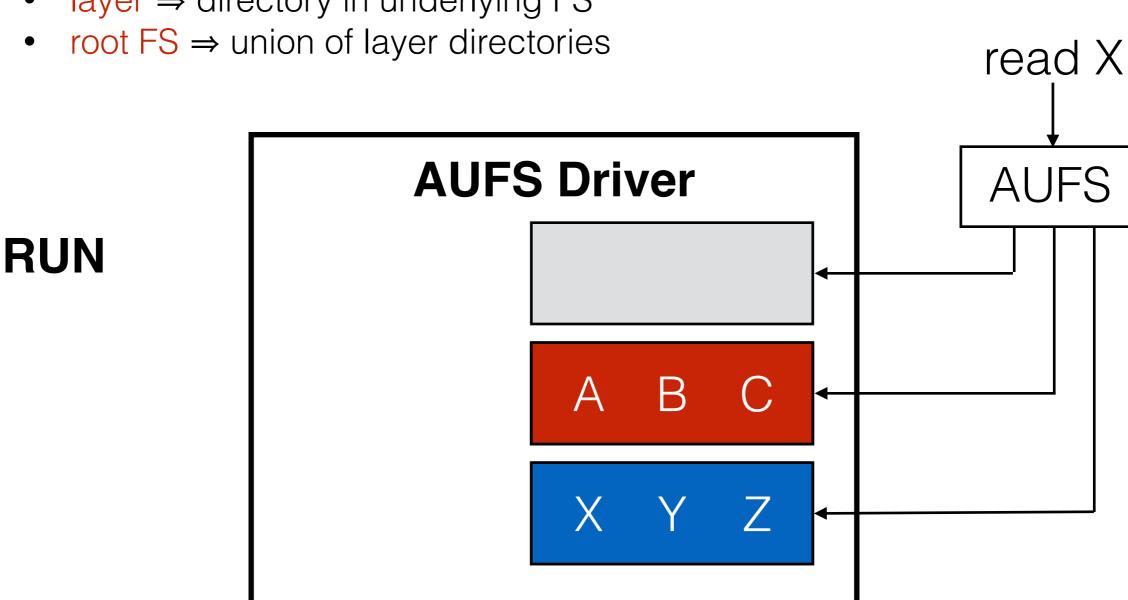


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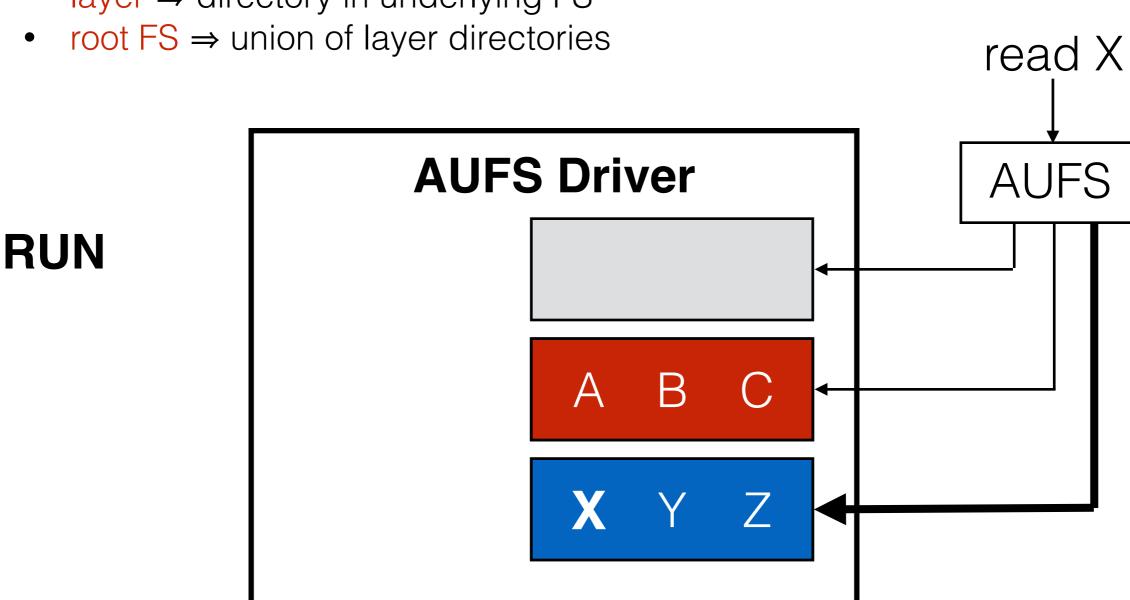
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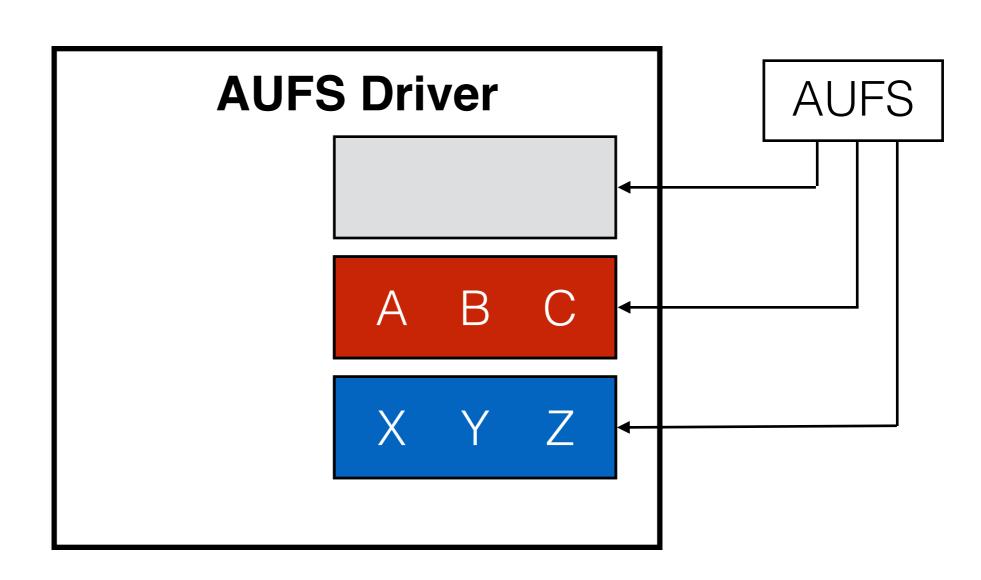


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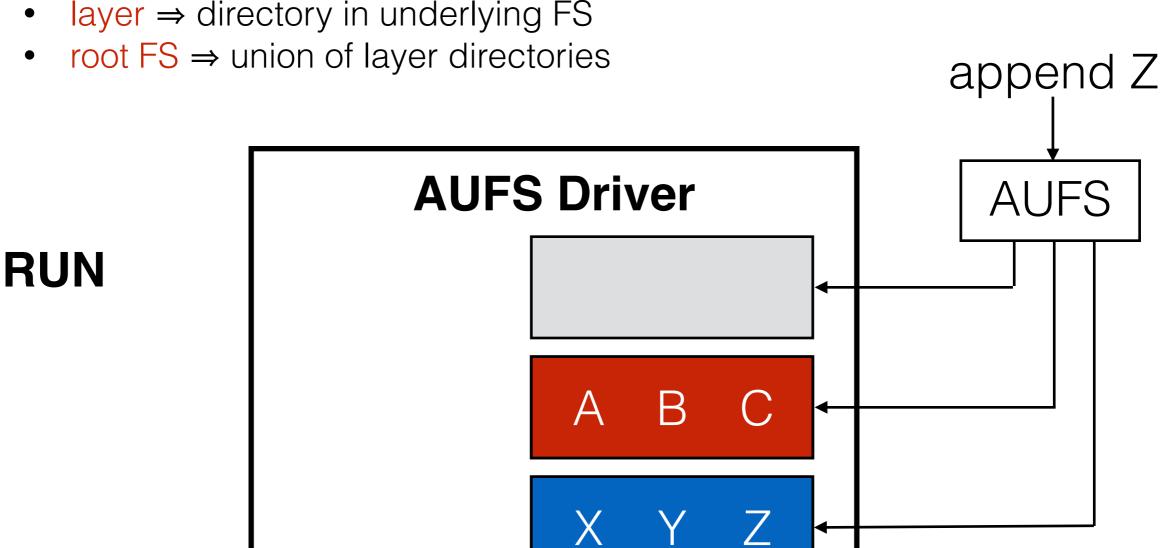


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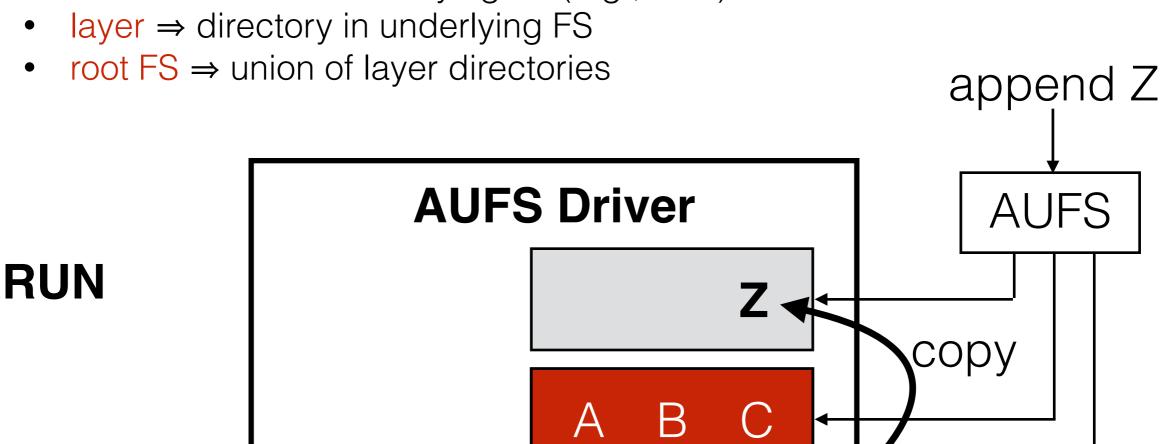


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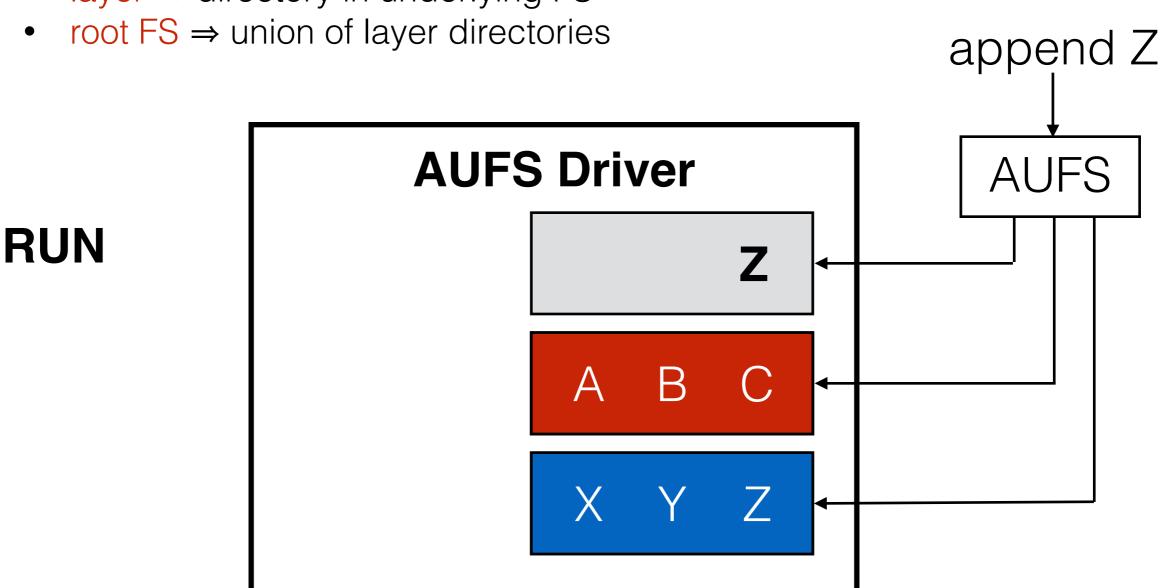
Uses AUFS file system (Another Union FS)

stores data in an underlying FS (e.g., ext4)



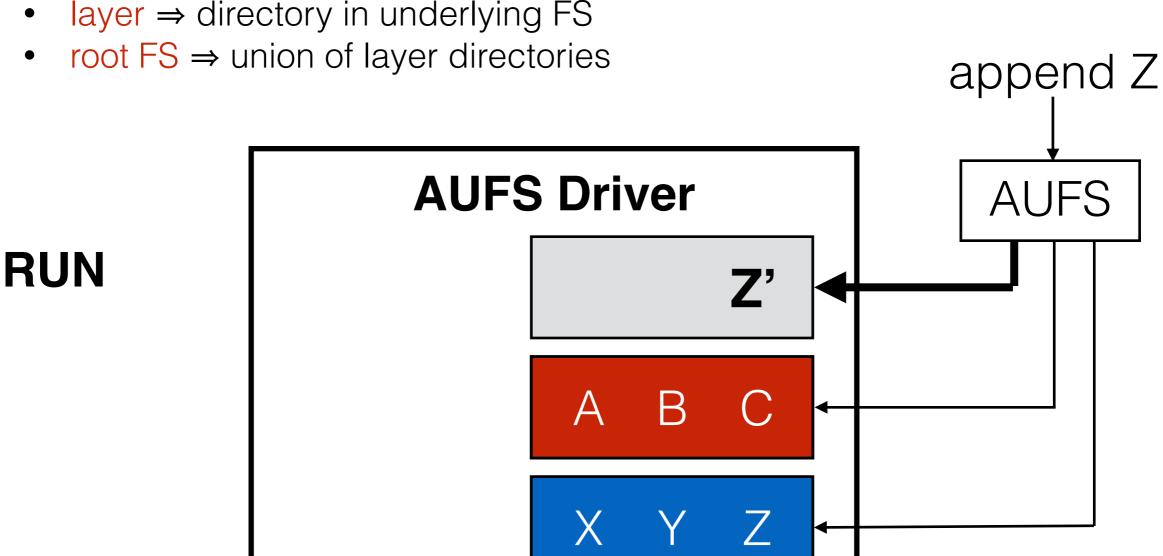
Uses AUFS file system (Another Union FS)

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- layer ⇒ directory in underlying FS



Uses AUFS file system (Another Union FS)

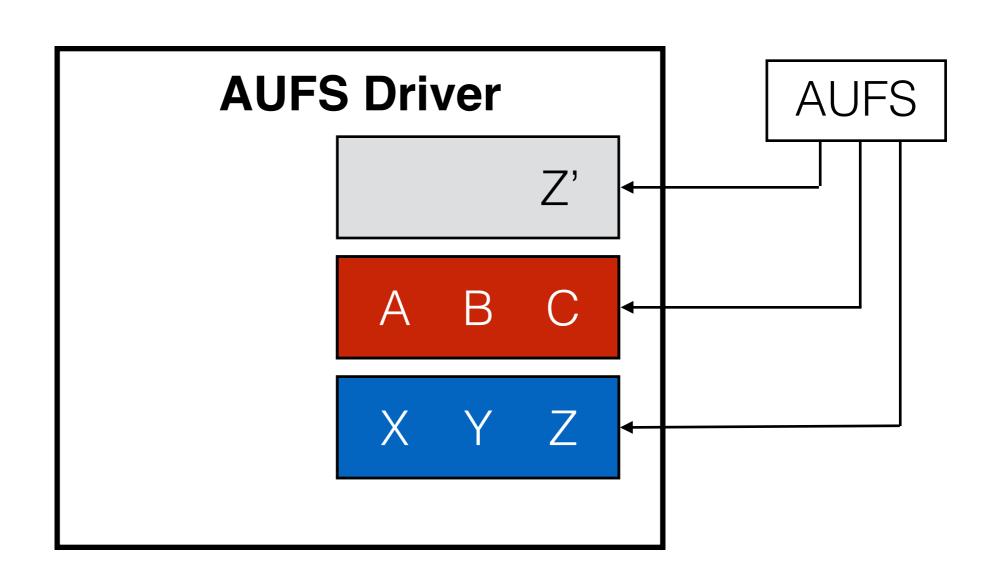
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Uses AUFS file system (Another Union FS)

- stores data in an underlying FS (e.g., ext4)
- layer ⇒ directory in underlying FS
- root FS ⇒ union of layer directories

RUN



Slacker Outline

Background

Container Workloads

Default Driver: AUFS

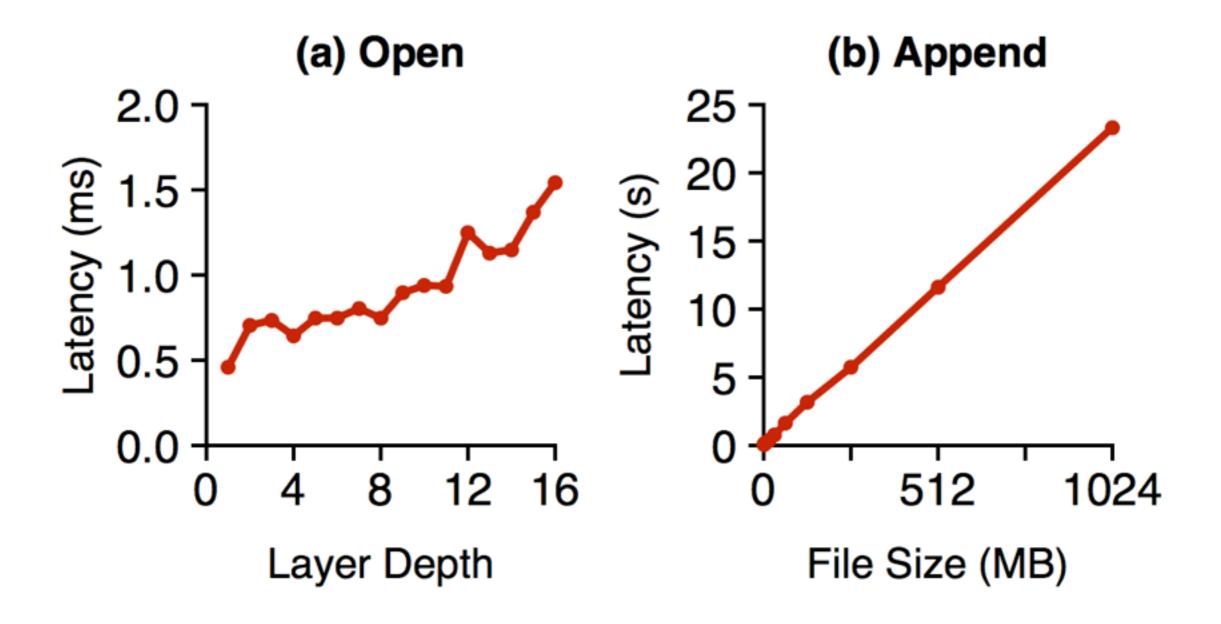
- Design
- Performance

Our Driver: Slacker

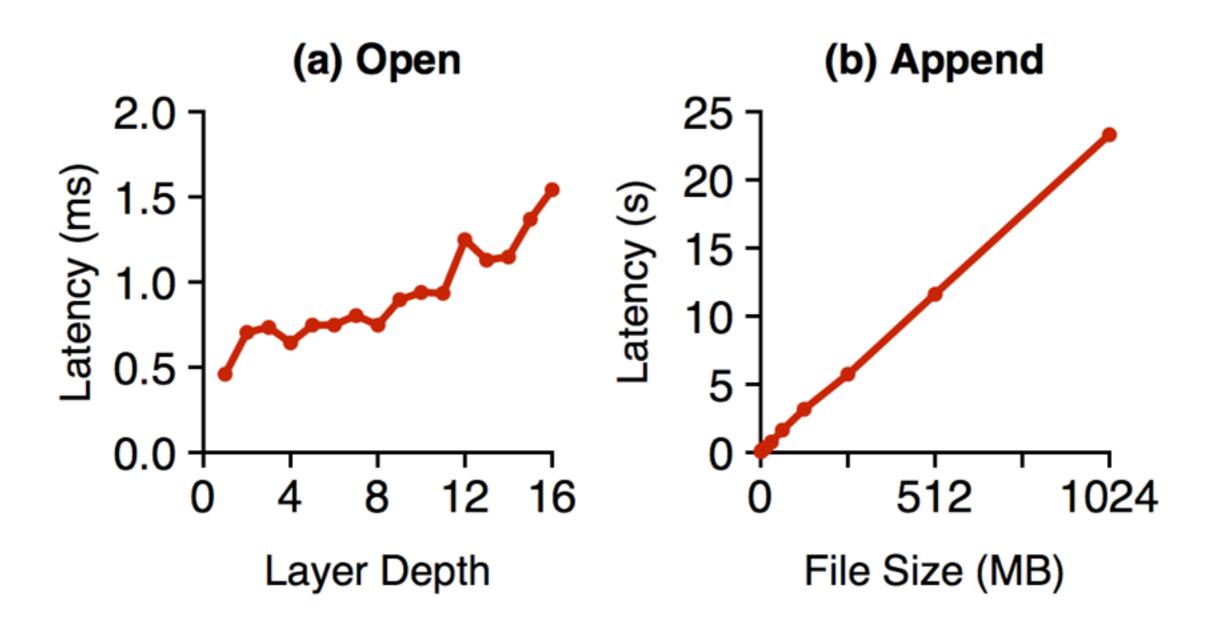
Evaluation

Conclusion

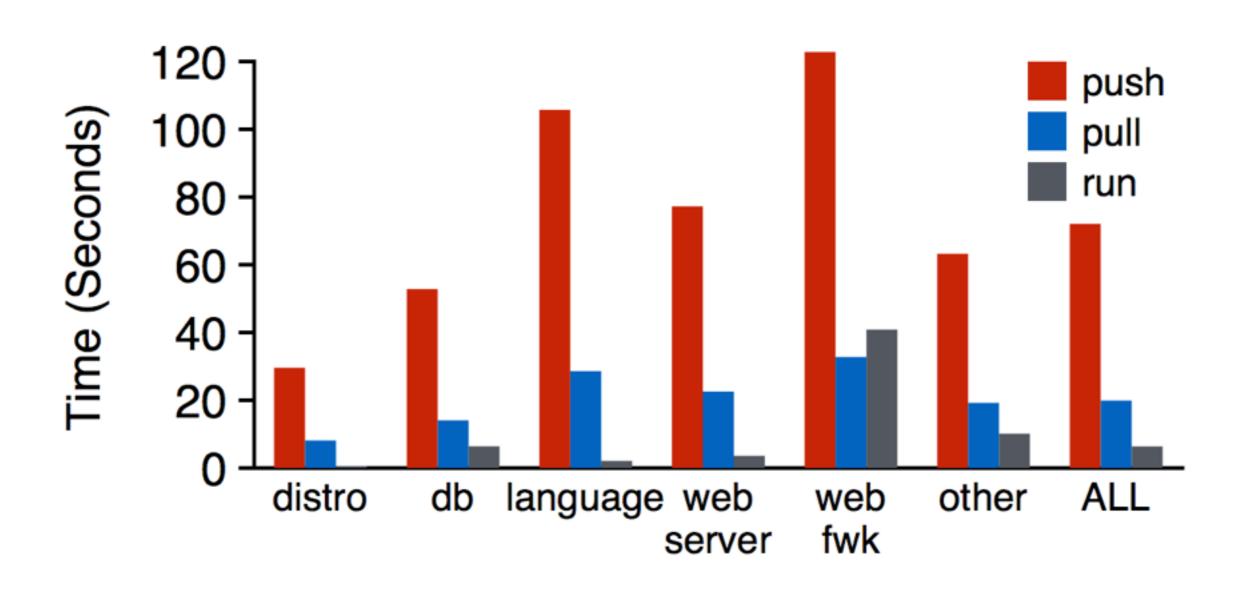
AUFS File System

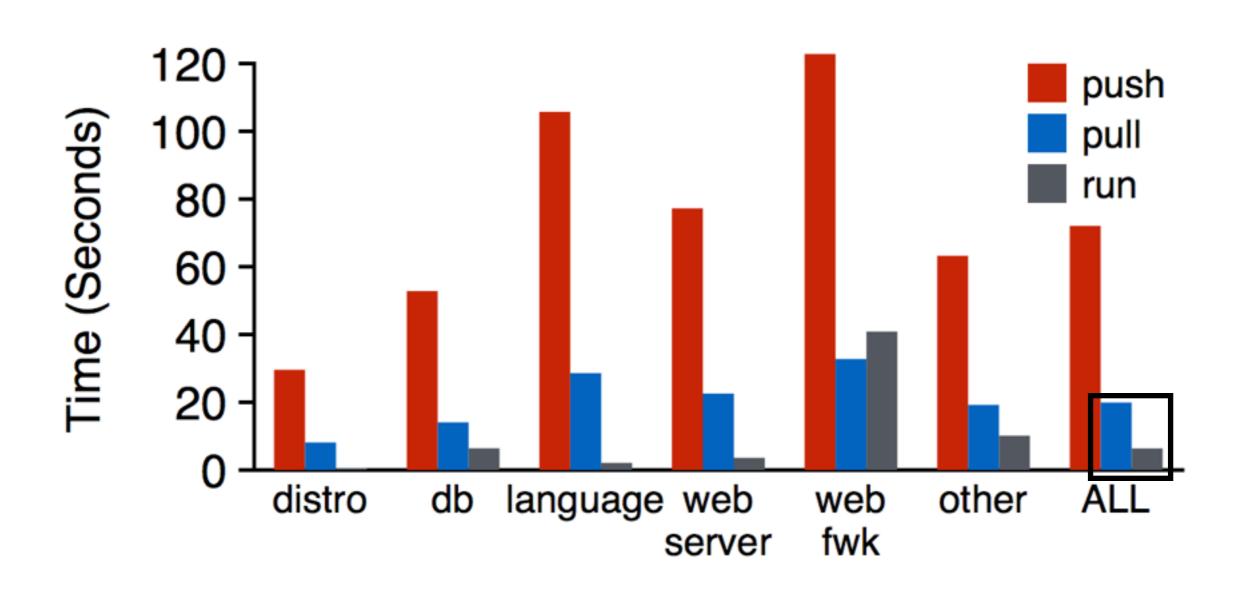


AUFS File System



Deep data is slow





76% of deployment cycle spent on pull

Slacker Outline

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Default Driver: AUFS

Our Driver: Slacker

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Conclusion

Slacker Driver

Goals

- make push+pull very fast
- utilize powerful primitives of a modern storage server (Tintri VMstore)
- create drop-in replacement; don't change Docker framework itself

Design

- lazy pull
- layer flattening
- cache sharing

Slacker Driver

Goals

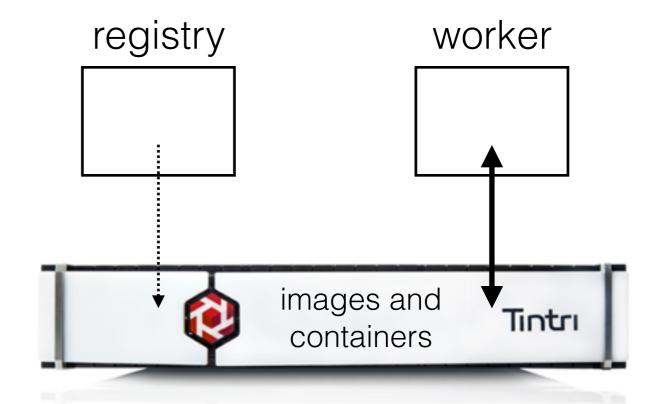
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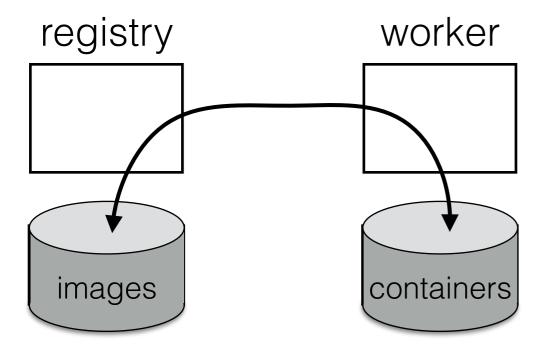
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AUFS

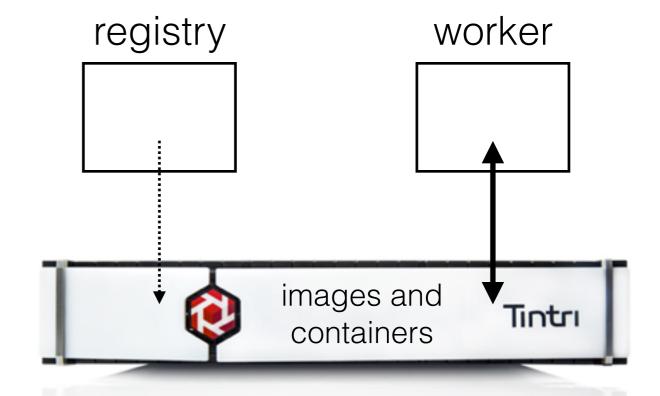
registry worker images containers



AUFS



Slacker



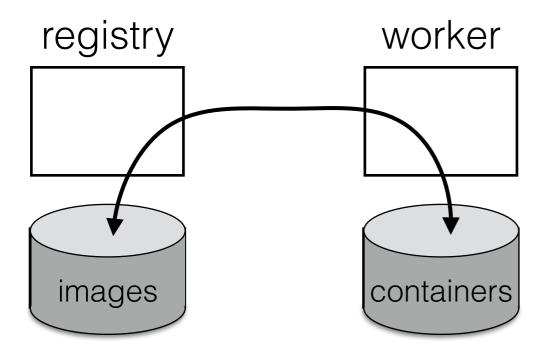
significant copying

- over network
- to/from disk

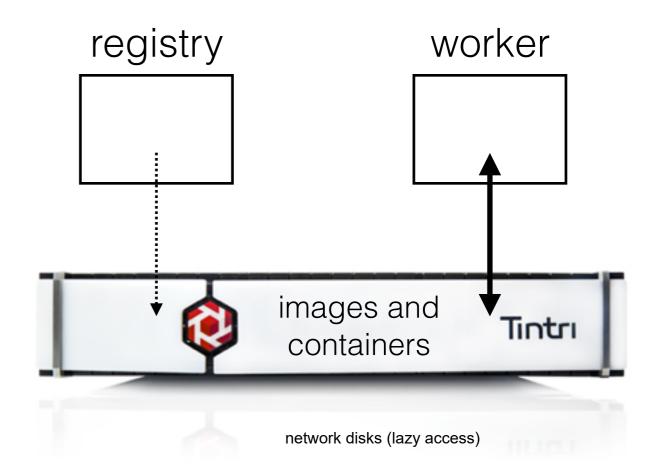
centralized storage

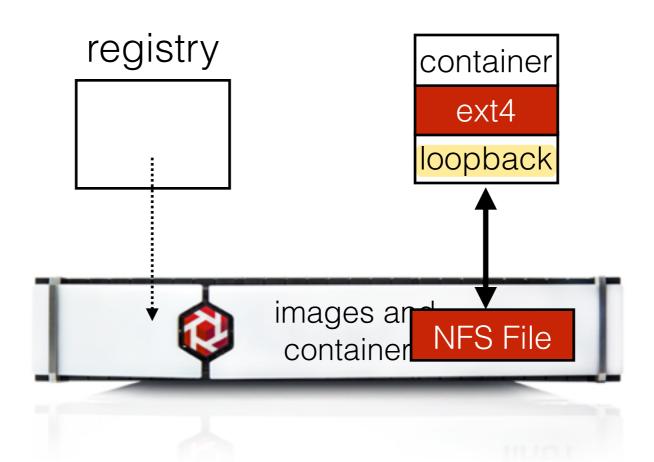
easy sharing

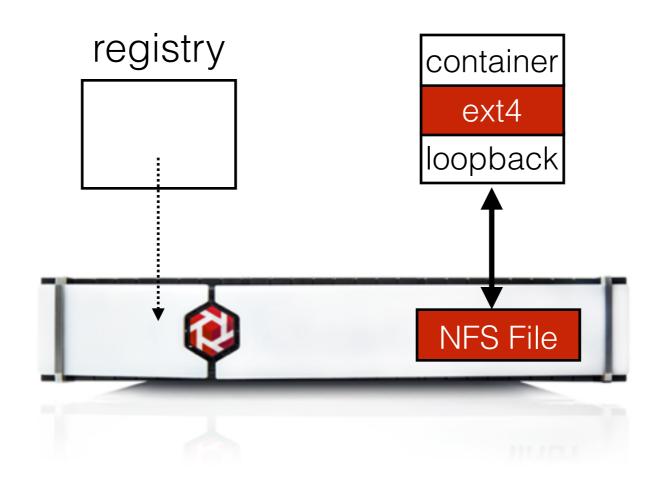
AUFS



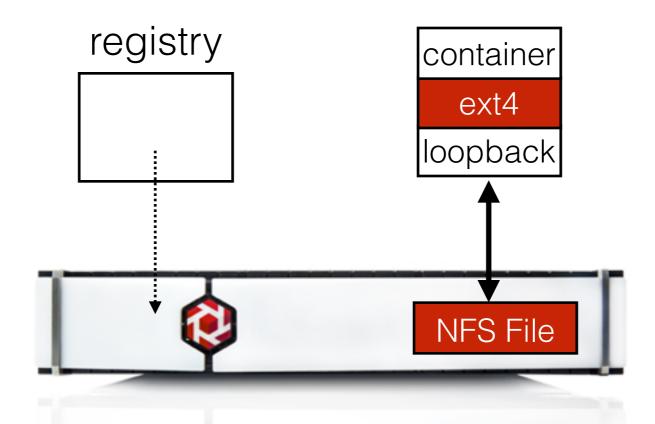
local disks (need to copy content to local disks)







Slacker



VMstore abstractions...

VMstore Abstractions

Copy-on-Write

- VMstore provides snapshot() and clone()
- block granularity avoids AUFS's problems with file granularity

snapshot(nfs_path)

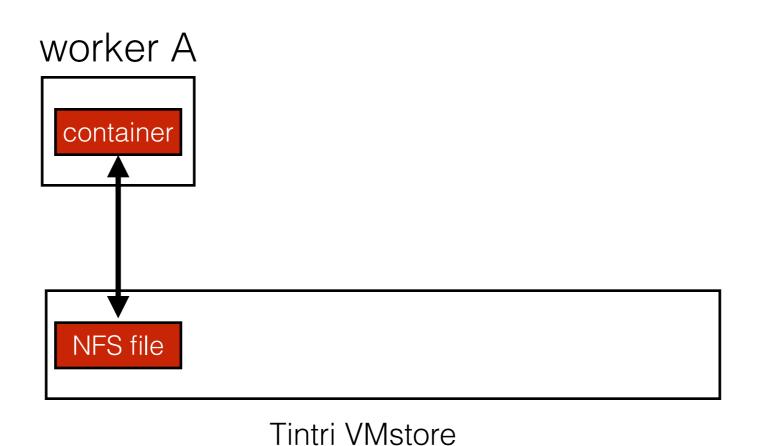
- create read-only copy of NFS file
- return snapshot ID

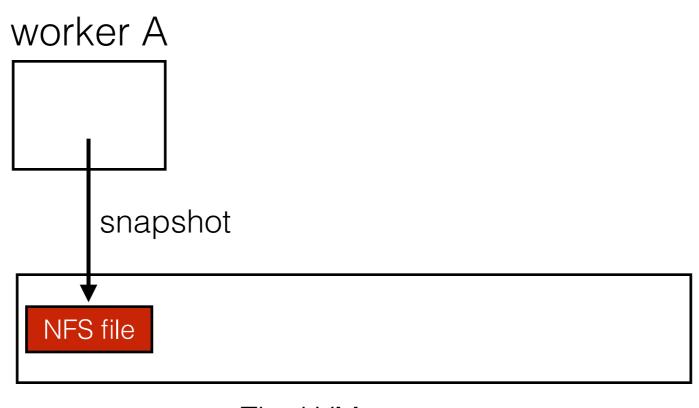
clone(snapshot_id)

create r/w NFS file from snapshot

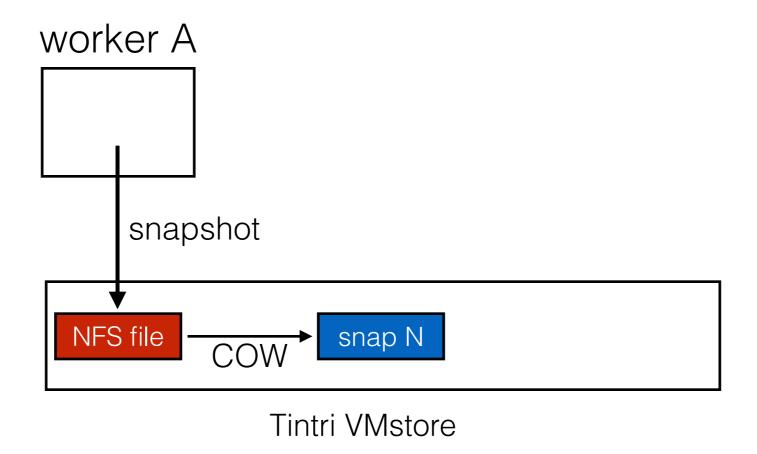
Slacker Usage

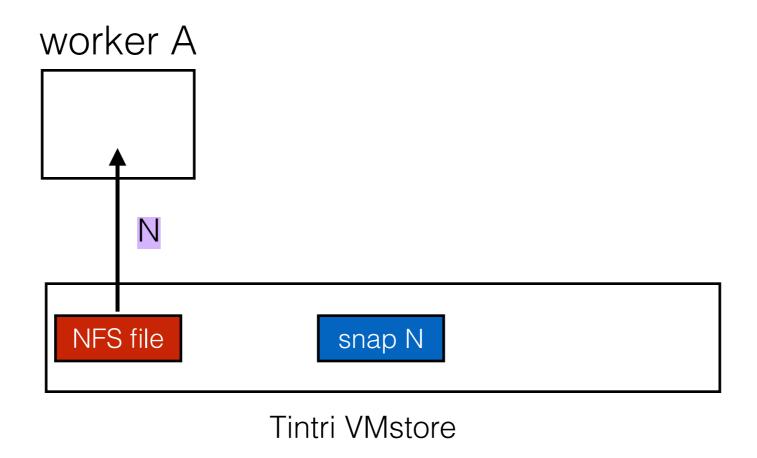
- NFS files ⇒ container storage
- snapshots ⇒ image storage
- clone () ⇒ provision container from image
- snapshot () ⇒ create image from container

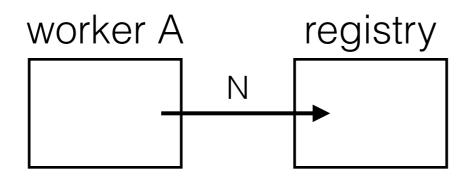


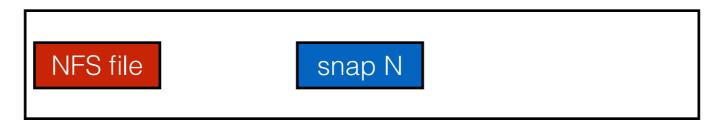


Tintri VMstore

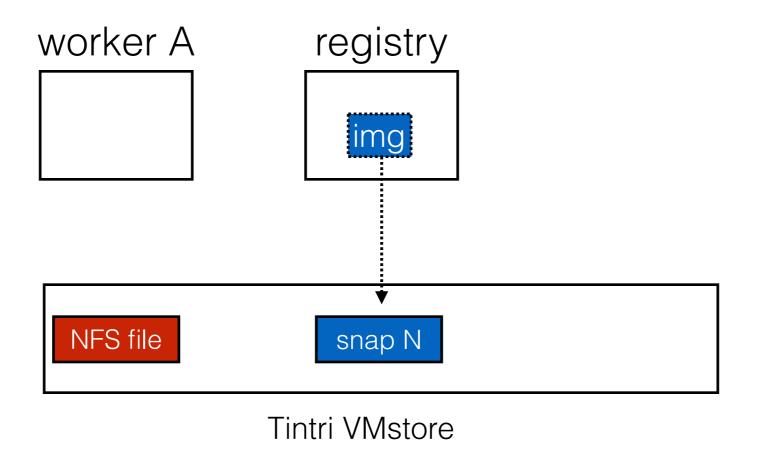






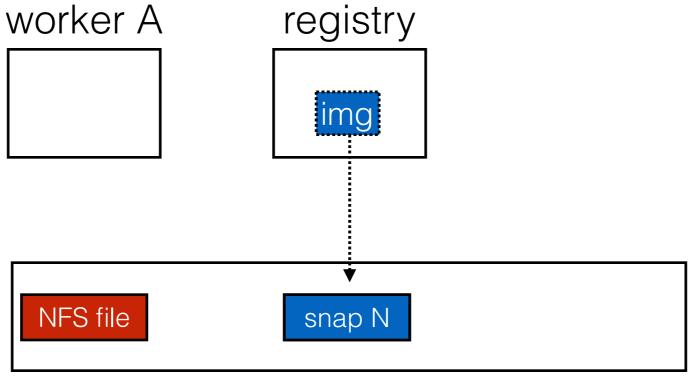


Tintri VMstore

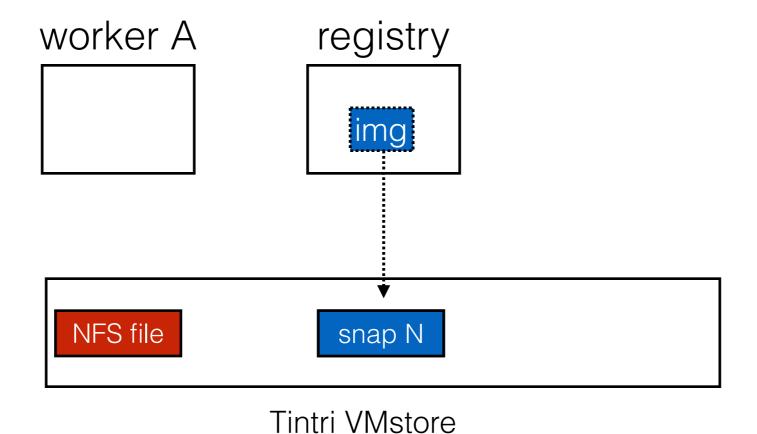


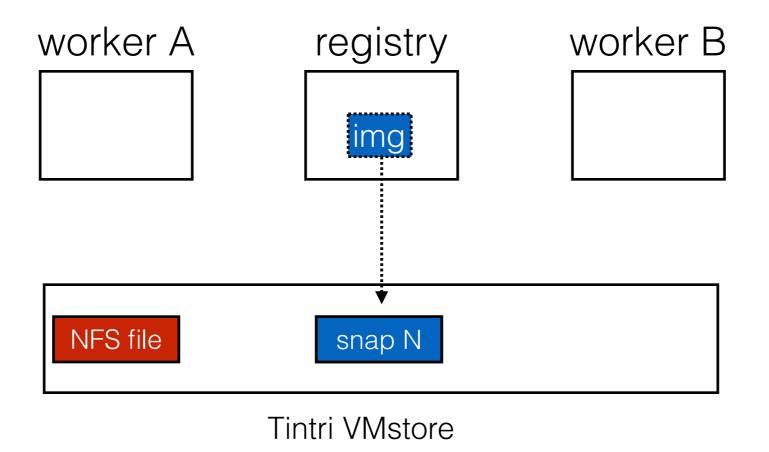
Note: registry is only a name server.

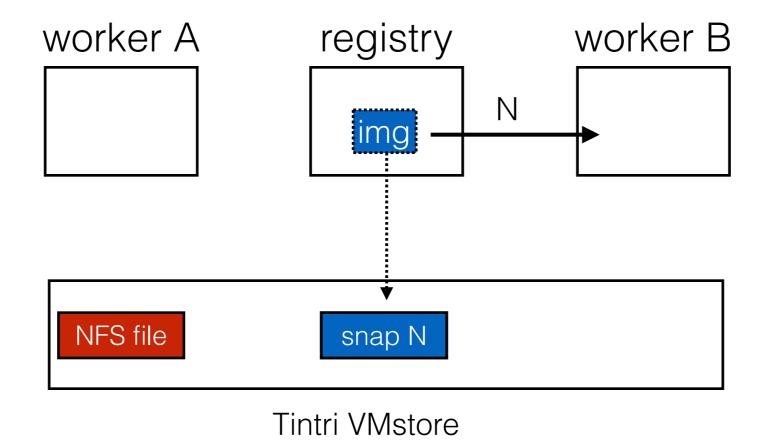
Maps layer metadata ⇒ snapshot ID

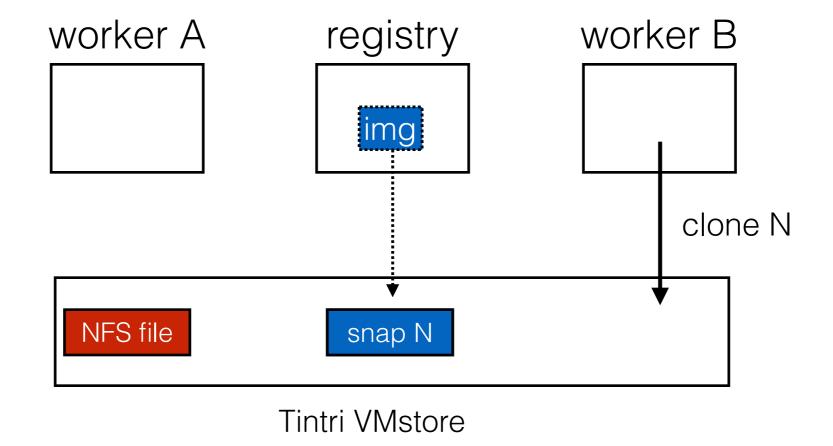


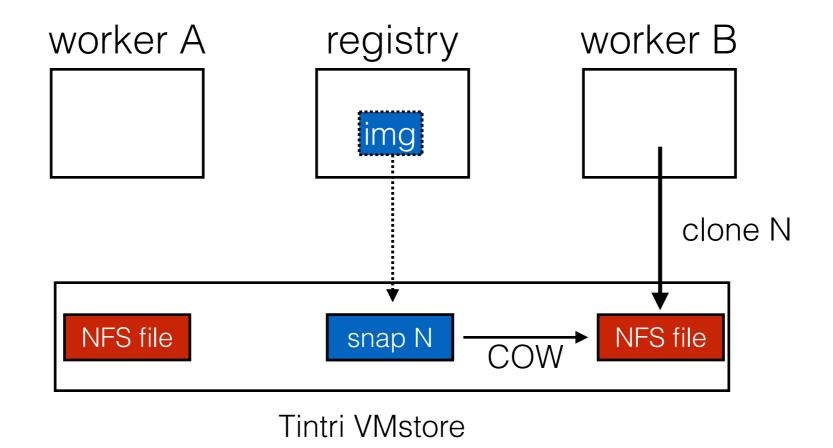
Tintri VMstore

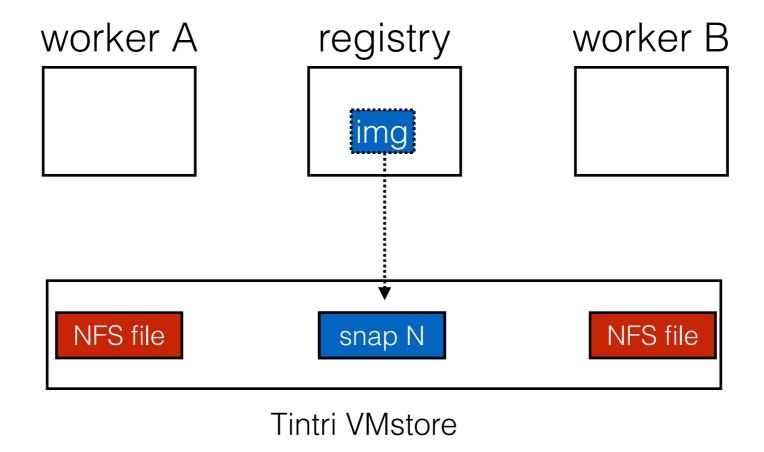


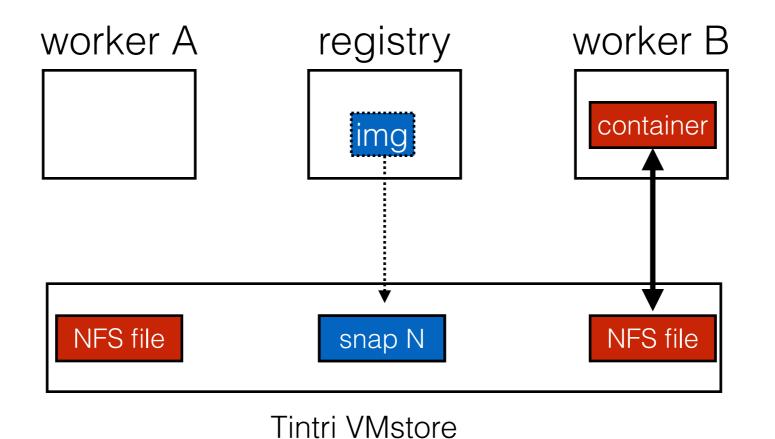












Slacker Driver

Goals

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Design

- lazy pull
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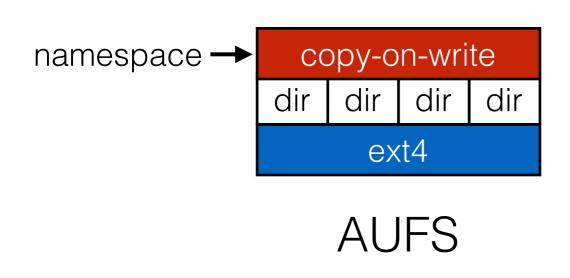
Slacker Flattening

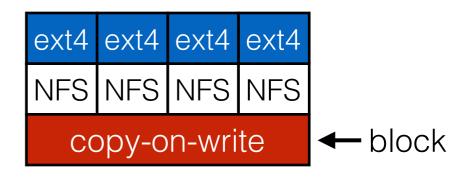
File Namespace Level

- flatten layers
- if B is child of A, then "copy" A to B to start. Don't make B empty

Block Level

do COW+dedup beneath NFS files, inside VMstore





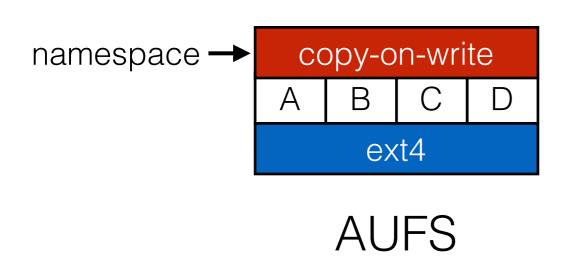
Slacker Flattening

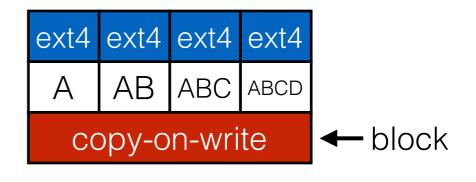
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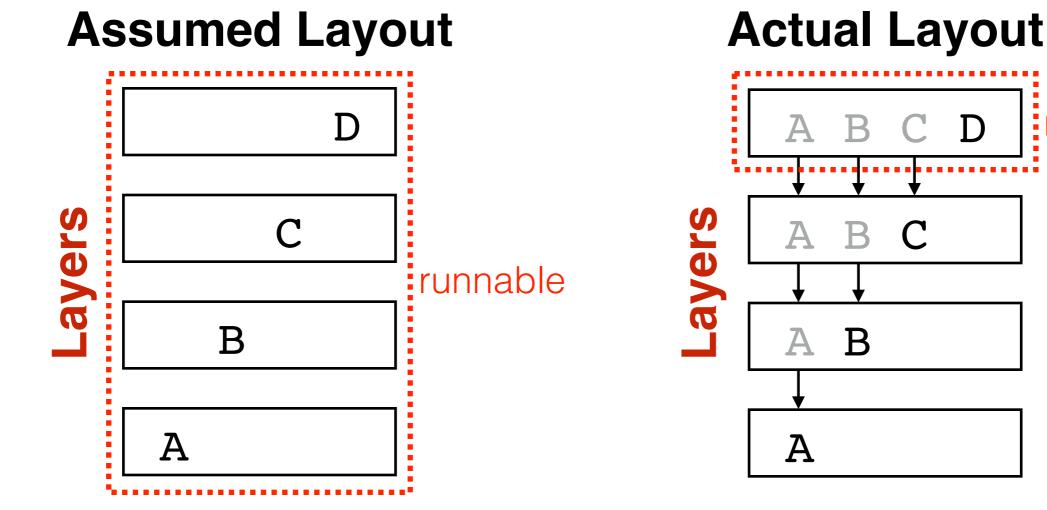
do COW+dedup beneath NFS files, inside VMstore





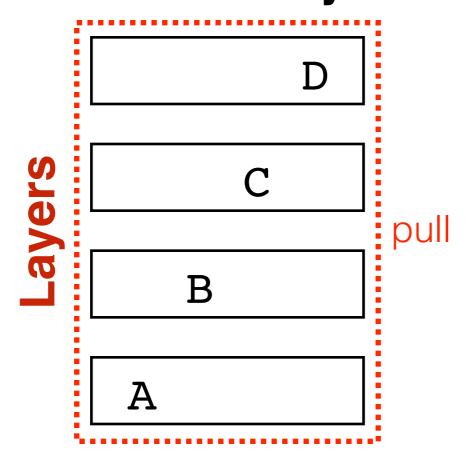
Challenge: Framework Assumptions

runnable

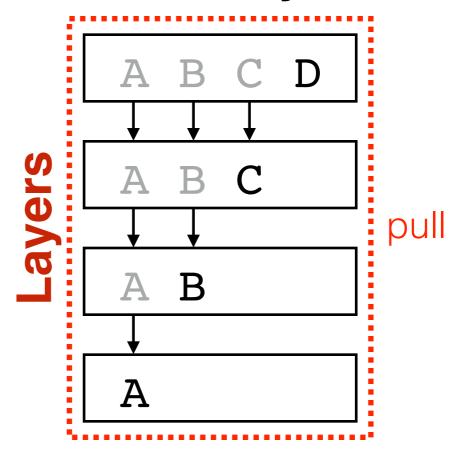


Challenge: Framework Assumptions

Assumed Layout



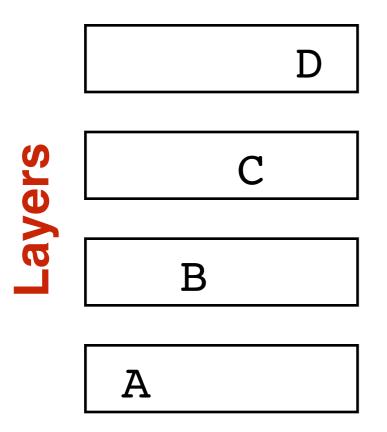
Actual Layout



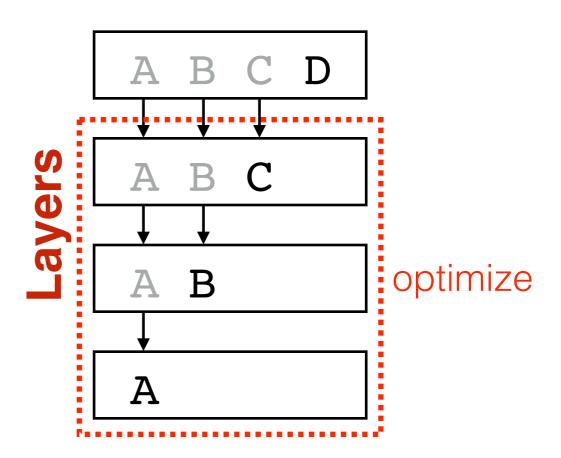
Challenge: Framework Assumptions

Strategy: **lazy cloning**. Don't clone non-top layers until Docker tries to mount them.

Assumed Layout



Actual Layout



Slacker Driver

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Slacker Driver

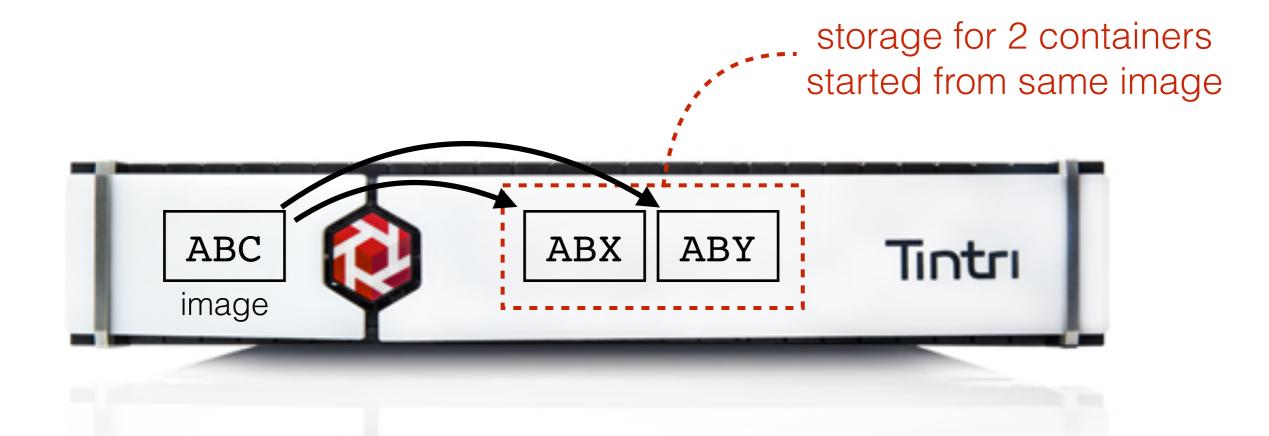
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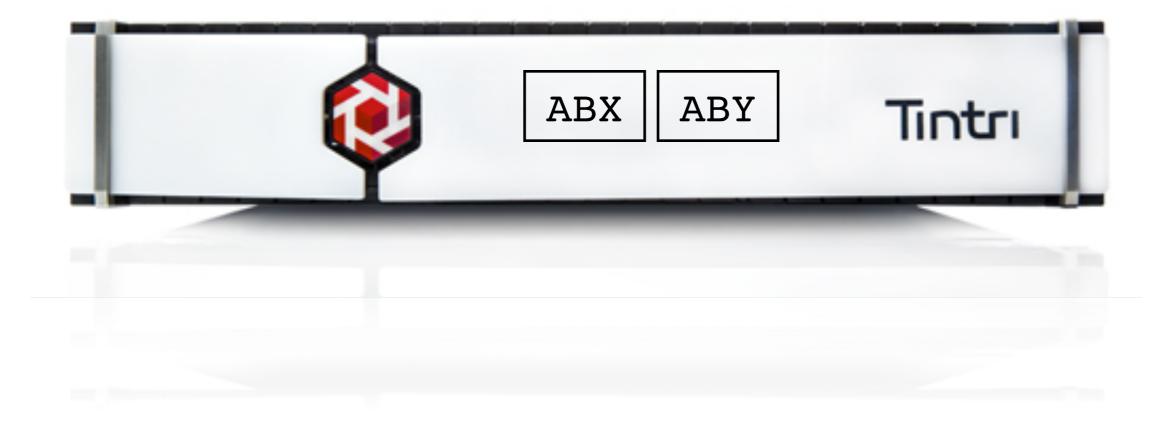
Design

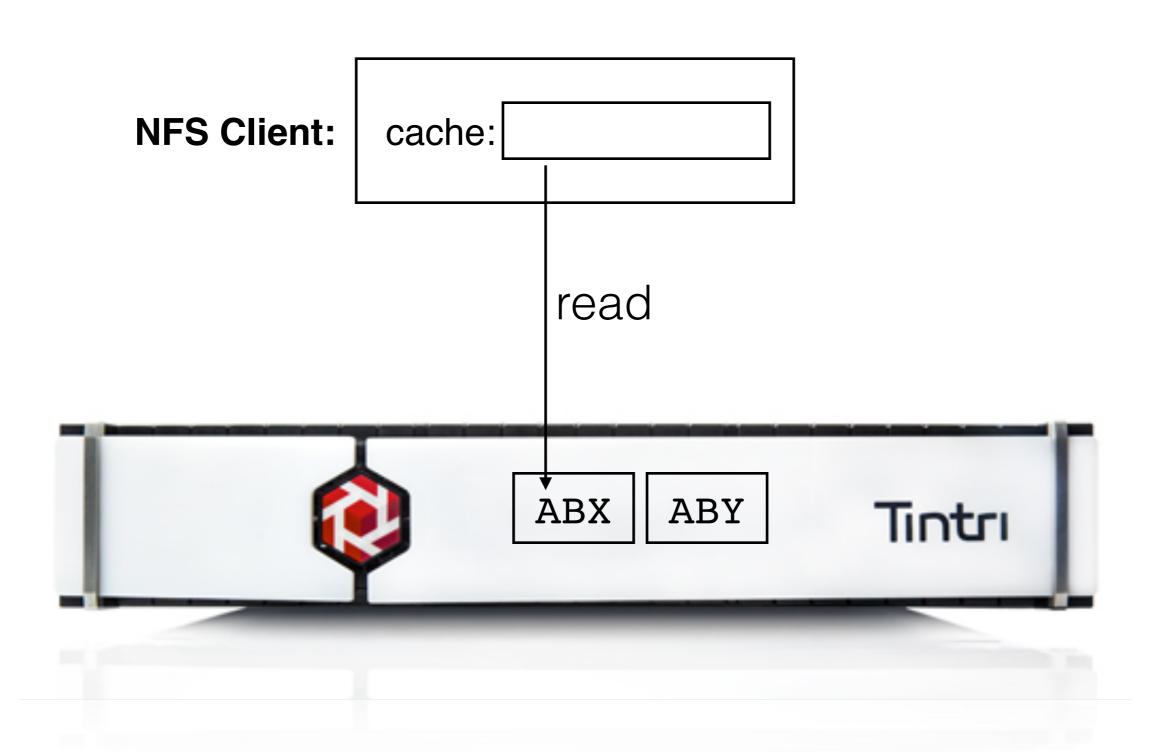
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NFS Client: cache:

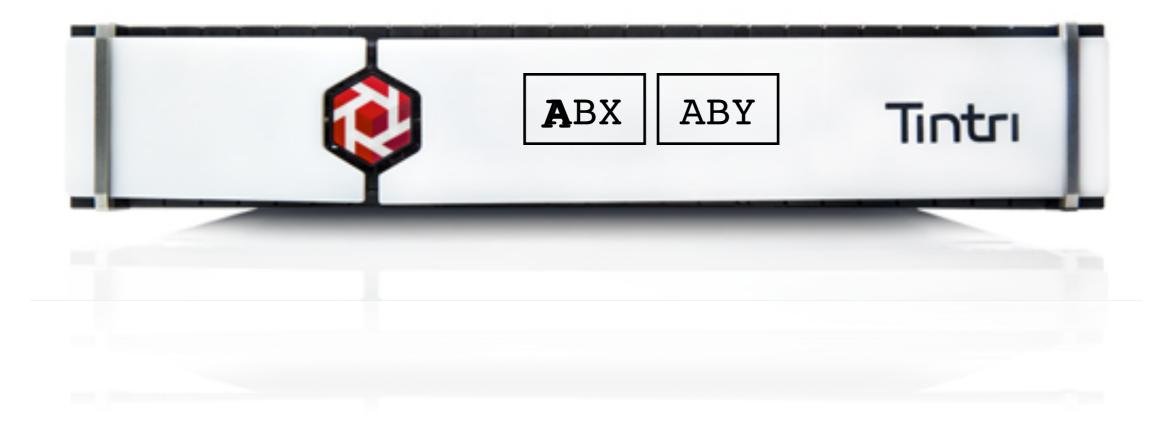


NFS Client: cache:

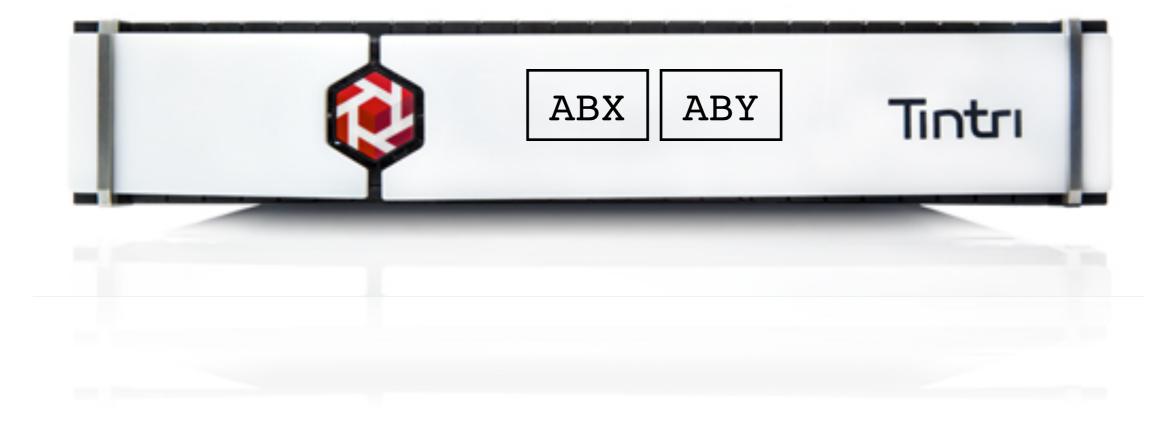


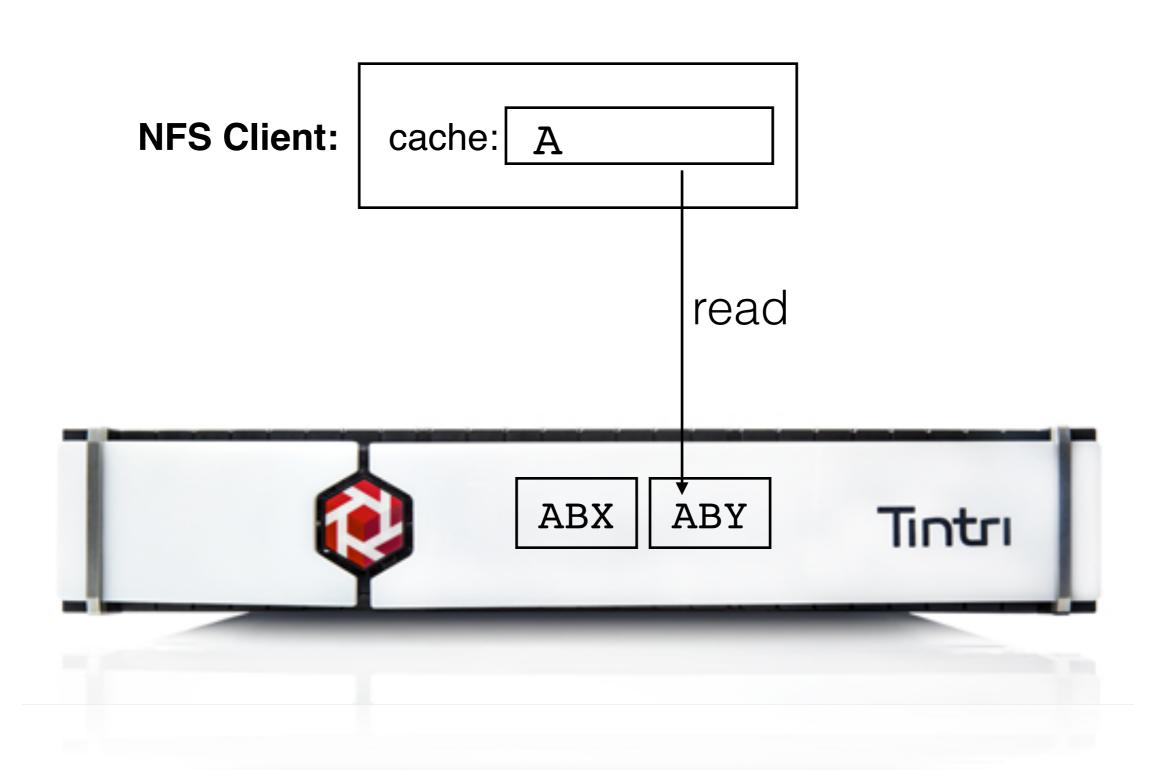


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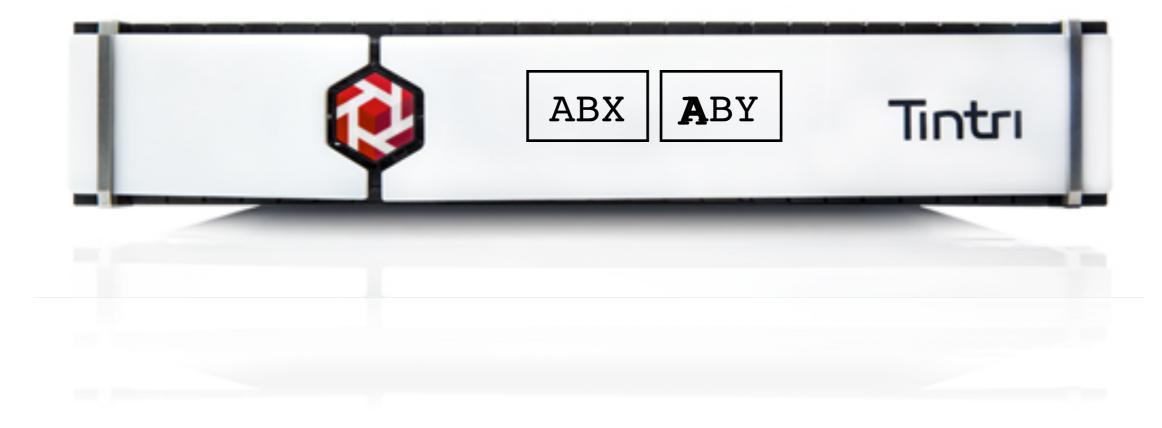


NFS Client: cache: A

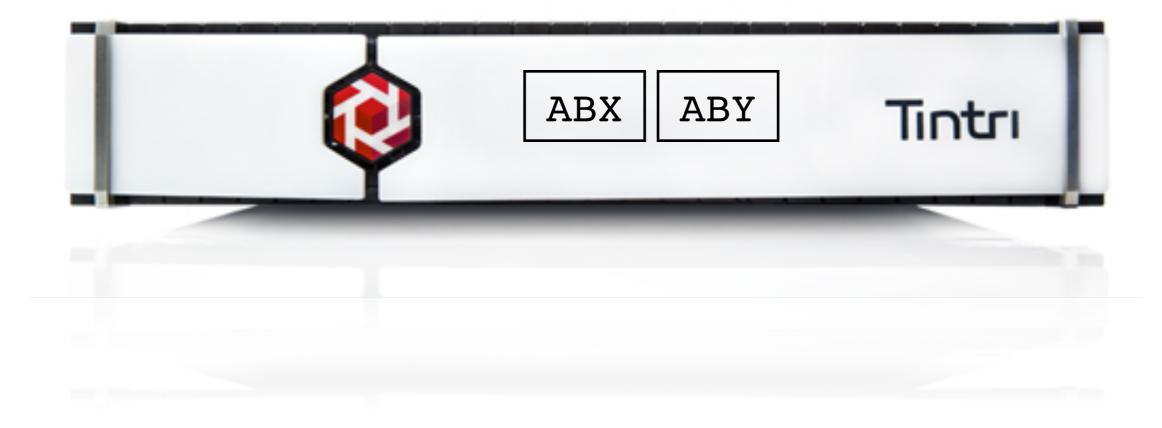




NFS Client: cache: A



NFS Client: cache: A A

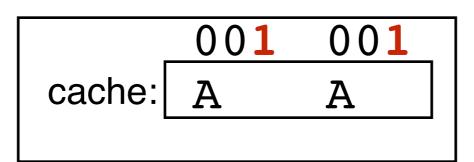


NFS Client: cache: A A

Challenge: how to avoid space and I/O waste?



NFS Client:



Strategy: track differences and deduplicate I/O (more in paper)



Slacker Outline

Background

Container Workloads

Default Driver: AUFS

Our Driver: Slacker

Evaluation

Conclusion

Questions

What are deployment and development speedups?

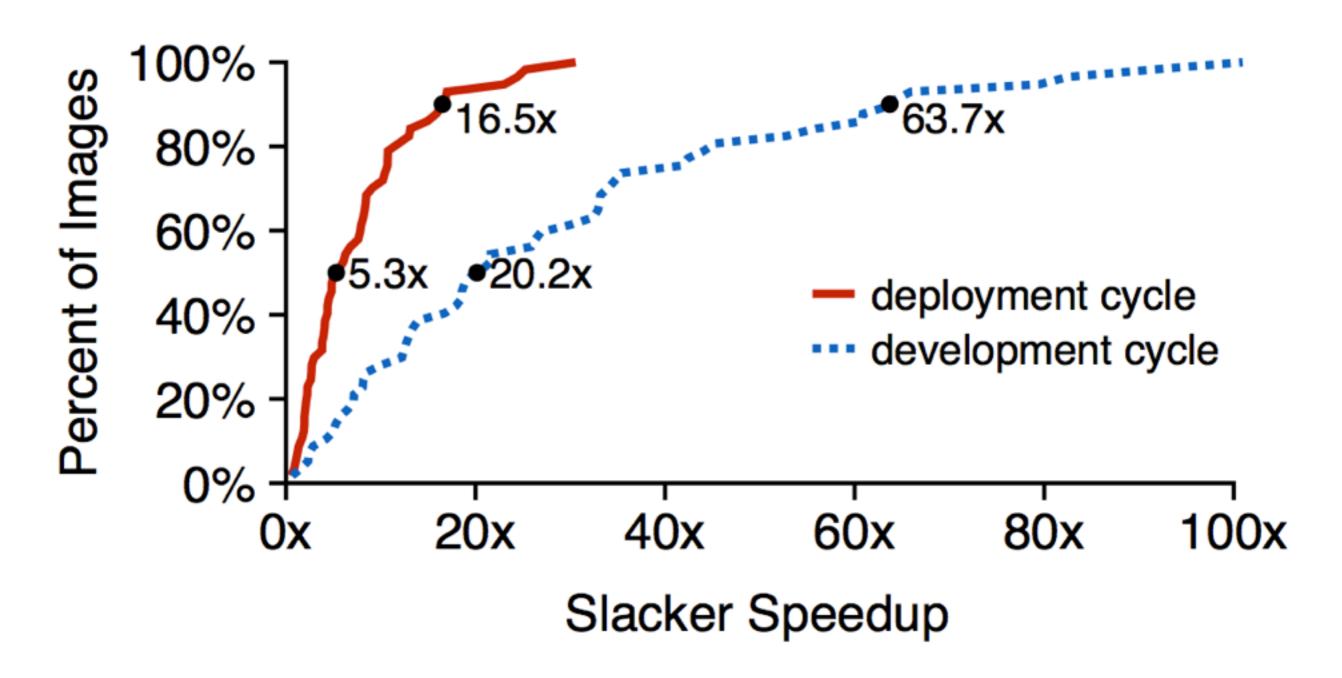
How is long-term performance?

Questions

What are deployment and development speedups?

How is long-term performance?

HelloBench Performance



deployment: pull+run

development: push+pull+run

Questions

What are deployment and development speedups?

• 5x and 20x faster respectively (median speedup)

How is long-term performance?

Questions

What are deployment and development speedups?

• 5x and 20x faster respectively (median speedup)

How is long-term performance?

Server Benchmarks

Databases and Web Servers

- PostgreSQL
- Redis
- Apache web server (static)
- io.js Javascript server (dynamic)

Experiment

- measure throughput (after startup)
- run 5 minutes

Server Benchmarks

Databases and Web Servers

- PostgreSQL
- Redis
- Apache web server (static)
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Experiment

- measure throughput (after startup)
- run 5 minutes

Result: Slacker is always at least as fast as AUFS

Questions

What are deployment and development speedups?

5x and 20x faster respectively (median speedup)

How is long-term performance?

there is no long-term penalty for being lazy

Slacker Outline

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Default Driver: AUFS

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Strengths:

- 1, very nice design of HelloBench, Slacker and MultiMaker
- 2. Good Iverage of VMStore based on their prior work for block ID based snapshot and clone
- 3. Good optimization with lazy fetch and ID-based cache; Modified kernel driver further enables client side caching.

Weaknesses:

- 1. The presentation has too many animation and failed to convey the beauty of the design and implementation.
- 2. Could have discussed the implication of Slacker to large parallel file or storage systems because it is unlikely for big data centers to use NSF-based backup.

Conclusion

Containers are inherently lightweight

but existing frameworks are not

COW between workers is necessary for fast startup

- use shared storage
- utilize VMstore snapshot and clone

Slacker driver

- 5x deployment speedup
- 20x development speedup

HelloBench: https://github.com/Tintri/hello-bench



