CNT4406/5412 Network Security PKI (Public Key Infrastructure)

Zhi Wang

Florida State University

Fall 2014

Zhi Wang (FSU)

CNT4406/5412 Network Security

Fall 2014 1 / 25

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- a repository for retrieving certificates
- a method of revoking certificates
- a method of evaluating a chain of certificates from trust anchors to ۰ the target name

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- A certificate is a signed message vouching a name with a public key
 [Bob's public key is 829248]_{Alice}
 - Alice is the issuer, and Bob is the subject

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- A trust anchor is a public key that is trusted to sign a certificate

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X.509 Certificate

X.509 Certificate: Example

A certificate for Bank of America signed by VeriSign



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X.509 Certificate: Example

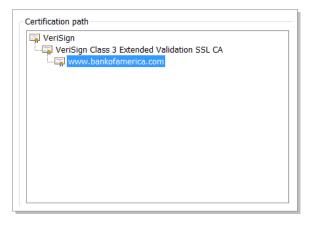
Details for the BOA certificate

Field	Value	-
Version	V3	
📑 Serial number	77 24 50 6d 4f 9a 87 9d	Ξ
📴 Signature algorithm	sha1RSA	
📴 Signature hash algorithm	sha1	
📑 Issuer	VeriSign Class 3 Extende	
🗒 Valid from	Tuesday, February 28, 20	
🖾 Valid to	Thursday, February 28, 2	
🗐 Subject	www.bankofamerica.com	-

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X.509 Certificate: Example

Verification path for the BOA certificate



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X.509 Certificate: Example

A certificate from VeriSign signed by VeriSign (why?)



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X.509 Certificate Format

• Basic fields:

Field	Description	Example
version	X.509 certificate version	V3
serial number	unique id for the certificate in the CA	77 24 50 6d
signature	algorithm used to signed the certificate	sha1 RSA
issuer	X.500 name of the issuer	VeriSign Class 3
validity from	starting date of the certificate	02/28/2012 8:00:00 PM
validity to	end date of the certificate	02/28/2013 8:00:00 PM
subject	X.500 name of the subject	www.bankofamerica.com
public key	public key of the subject	RSA 30 82 01 0a

*CA may have multiple keys

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• Optional fields: basic constrains, key usage, CRL distribution points, authority key identifier*, subject key identifier...

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PKI Trust Models

PKI trust models defines where to get trust anchors, and how to evaluate a chain of certificates from trust anchors to the target

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- Monopoly model
- Monopoly + RAs
- Oelegated CAs
- Oligarchy model
- Anarchy model
- Top-down with name constraints
- Ø Bottom-up with name constraints

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Monopoly Model

One CA is universally trusted by the world
 we everyone must get certificates from the (only) CA

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Monopoly Model

- One CA is universally trusted by the world
 we everyone must get certificates from the (only) CA
- The CA's public key is embedded in all software and hardware
 the only trust anchor

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• There is NO universally trusted organization

Image: A matched and A matc

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- \bullet Monopoly control ${\rightarrow}\mathsf{CA}$ can charge whatever it wants for a certificate
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 infeasible to change the CA's key everywhere if compromised
- It is expensive and insecure for the CA to remotely certify a public key

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Delegated CAs

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 Firefox (ver 15.0.1): hundreds of root certificates from 86 organizations around the world, none revoked

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- Some software (browsers) accept expired certificates
 Google Chrome includes a certificate from Microsoft which has expired since 12/30/1999

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- The database would be unworkably large if deployed on Internet scale
- It's hard to verify the trustworthiness of every certificate on the chain
 mo problem for a small community where everyone is trustworthy
 mo can we trust a certificate if there are multiple chains?

Models with Name Constraints

- A CA is only trusted for certifying users in his domains
 FSU CA certifies FSU students
 - a certificate for Amazon.com from Nigeria?

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- A CA is only trusted for certifying users in his domains
 FSU CA certifies FSU students
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- Name constraints defines they way to verify (search) a target
 im top-down or bottom-up

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 imitiar to the monopoly model with delegate CAs

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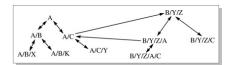
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- It has the other problems of the monopoly model

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Bottom-up with Name Constraints

• The model assumes a hierarchical namespace (like DNS)

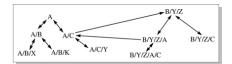


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Fall 2014 17 / 25

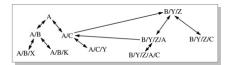
Bottom-up with Name Constraints

- The model assumes a hierarchical namespace (like DNS)
- Each organization can create its own PKI and link to others
 a parent certifies its children (down-link)
 - **••** a child also certifies its parent (**up-link**)
 - **a cross-link** connects two nodes neither is an ancestor of the other



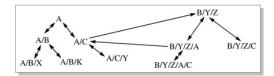
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 - **a cross-link** connects two nodes neither is an ancestor of the other
- The namespace can be traversed starting from any node
 follow up-links and/or one cross-link to an ancestor of the target
 - follow donw-links (only) from there



Bottom-up with Name Constraints...

• How can A/C/Y verify the certificate of B/Y/Z/C?



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 Fall 2014
 18 / 25

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Bottom-up with Name Constraints...

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- No monopoly is possible

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 PKI can be deployed in any organization independent of others
 - Authentication between users in the same organization stays within
- It is easy to find out if a path exists
- No monopoly is possible
- Replacing any key is reasonably easy

Certificate Revocation

• Certificates might need to be revoked

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- Certificates might need to be revoked
 - employment is terminated, private key is stolen...
- Certificates normally have an expiration time
 - validity time is months, lots of damage can be done in between
 - browsers normally do not enforce expiration time!

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• CA maintains and periodically releases a CRL

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- CA maintains and periodically releases a CRL
- Each transaction is checked against the CRL

Certificate Revocation List

• Why are CRL issued periodically even if no certificates are revoked?

Certificate Revocation List

- Why are CRL issued periodically even if no certificates are revoked?
- How frequent should CRL be issued?

Certificate Revocation List

- Why are CRL issued periodically even if no certificates are revoked?
- How frequent should CRL be issued?
- If a CRL is maintained, why set an expiration time for certificates?

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Delta CRL

• CRL may be very large and needs to be downloaded frequently

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Delta CRL

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- A delta CRL lists changes from the last complete CRL
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- Delta CRLs can be issued very frequently
 im full CRLs are issued less frequently

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- OLRS is a online system to query the revocation status of certificates
 OLRS maints the full CRL list
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- OLRS can be a performance bottleneck
 - Alice can obtain a (timestamped) certificate from OLRS "Alice's certificate was not revoked as of ..."
 - Bob can query OLRS in advance and cache the result

Image: A Image: A

• CRL is a **black-list** of revoked certificates

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- CRL is a black-list of revoked certificates
- How about maintaining a list of valid certificates?
 Is it more secure than bad-lists?

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 solution: use hashes of the certificates

Summary

- X.509 certificate
- PKI trust models
- Certificate revocation
- Next lecture: IPSec

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