

# Computer Security

CIS 5370

## Security Policies

## Take-Grant Protection Model

- A specific (not generic) system
  - Set of rules for state transitions
- Safety decidable, and in time linear with the size of the system
- Goal: find conditions under which rights can be transferred from one entity to another in the system

## Take-Grant Notation

○: objects (files, ...)  
 ●: subjects (users, processes, ...)  
 ⊗: either a subject or an object

$G \xrightarrow{x} G'$ : apply a rewriting rule  $x$  to  $G$  to get  $G'$

$G \xrightarrow{*} G'$ : apply a sequence of rewriting rules to  $G$  to get  $G'$

$R = \{ t, g, r, w, \dots \}$ : set of rights

## Four Rules in Take-Grant

Take, grant, create, remove

## More Rules

create      ●      | -      ● —  $\alpha$  — ⊗

remove      ● —  $\alpha$  — ⊗      | - ④      ● —  $\alpha - \beta$  — ⊗

These four rules are called the *de jure* rules

Rule sequences are called *witnesses*

## Subjects and Objects

- Both subjects and objects can hold rights to objects
- Not sure what it means for an object to hold rights to another object (or to a subject)
- Only subjects can take or grant rights
  - See definition of take and grant
  - Important in later results

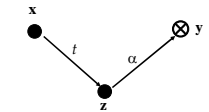
## Transferring rights

- Two connected subject vertices may transfer rights in 4 ways:
  1. Take
  2. Grant
  3. Symmetric Take
  4. Symmetric Grant

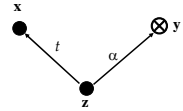
7

## Symmetry

We know that if node  $z$  has  $\alpha$  rights to  $y$  and  $x$  has take ( $t$ ) rights to  $z$ , then  $x$  can gain  $\alpha$  rights to  $y$ . This is the take rule.



What if  $x$  has no rights to  $z$ , and  $z$  has  $\alpha$  to  $y$ , but has only take ( $t$ ) rights to  $x$ ? Can  $x$  gain  $\alpha$  rights to object  $y$ ?



Yes, but we will need a *witness*!

8

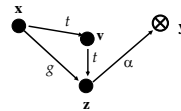
## can•share Predicate

Definition:

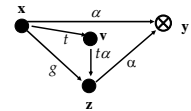
- $can\bullet share(r, \mathbf{x}, \mathbf{y}, G_0)$  if, and only if, there is a sequence of protection graphs  $G_0, \dots, G_n$  such that  $G_0 \vdash^* G_n$  using only *de jure* rules and in  $G_n$  there is an edge from  $\mathbf{x}$  to  $\mathbf{y}$  labeled  $r$ .

9

## Can-share

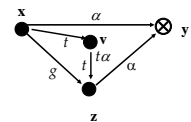


$\vdash ?$



Can subject  $x$  gain  $\alpha$  rights to object  $y$ ?

1.  $v$  takes ( $\alpha$  to  $y$ ) from  $z$
2.  $x$  takes ( $\alpha$  to  $y$ ) from  $v$



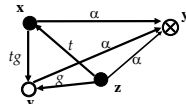
10

## Symmetry

Can subject  $x$  gain  $\alpha$  rights to object  $y$ ?



$\vdash \textcircled{4}$



1.  $x$  creates ( $tg$  to new)  $v$
2.  $z$  takes ( $g$  to  $v$ ) from  $x$
3.  $z$  grants ( $\alpha$  to  $y$ ) to  $v$
4.  $x$  takes ( $\alpha$  to  $y$ ) from  $v$

11

## Definitions

- **tg-path**: what you would think, path where each pair of connected vertices share at least one  $t$  or  $g$  edge
- **tg-connected**: two vertices are **tg-connected** if there is a **tg-path** between them
- **island**: subject-only subset of vertices in a protection graph that are **ALL** **tg-connected**

12

## Exercise

- Prove that any right possessed by any vertex in an island can be shared with any other vertex in the island

13

## Take – Grant Properties

- **Safety question is decidable in linear time on the complexity of the protection system**
- **In any Take-Grant protection system, there is an efficient algorithm to decide if a specific subject can gain a right to specific object**

14

## Key Question

- Characterize class of models for which safety is decidable
  - Existence: Take-Grant Protection Model is a member of such a class
  - Universality: In general, question undecidable, so for some models it is not decidable
- What is the dividing line?

15