## **Class Participation Quiz**

## Complete and submit this within 15 minutes.

## **General Information**

Last Name \_\_\_\_\_ First Name \_\_\_\_\_

ID Number \_\_\_\_\_ Email \_\_\_\_\_

## Fill in the blanks (or select T/F):

 $\underline{\qquad} T(n) = T(\sqrt{n}) + O(1)$ 

 $\underline{\qquad} T(n) = 2T(\sqrt{n}) + O(1)$ 

\_\_\_\_\_ IOs are necessary and sufficient to sort n comparable objects of O(1) size in the IO Efficient model.

\_\_\_\_\_ is an upper bound on the query time of a point in kd-tree in fixed dimensions.

\_\_\_\_\_ is an upper bound on the sorting time in the PRAM model of computation.

Show proof on the back of this page.

 ${\bf T} \quad {\bf F} \mbox{ For all weighted input graphs (with positive capacities), The Preflow push algorithm can be sped up by using capacity scaling. }$ 

 ${\bf T} - {\bf F}$  Unit weight bipartite matching can be done using max flow algorithm.

 $\mathbf{T}$  **F** In every directed graph with nodes s and t (and all edges having unit weight), the maximum number of edge-disjoint s-t paths is equal to the minimum number of edges whose removal separates s from t.

**T F** If a bipartite graph  $G = (L \cup R, E)$  has a perfect matching, then |N(S)| < |S| for all subsets  $S \subseteq L$ .

**T F** Factor is not in NP $\cap$ co-NP.

 ${\bf T} - {\bf F}$  In fixed dimensions, LP can be solved in linear time.

**T F** LRU is 2k-competitive.

Show in a two line proof, LP Feasibility and LP optimization are equivalent. (Use the following space on this page)