Overview

Over the next four weeks, we will look at these topics:

- Building Blocks
- Advanced Authentication Issues
- **Security** ■



Overview

- Storage and its abstraction
- Virtualization and appliances
- Data Replication and Disaster Recovery (BCP/COOP)
- Physical Environments



At the system level: Unix/Linux Building blocks

I want to take some time to talk about one of the fundamental toolsets that most programs that system administrators work with are built over.



At the system level: Unix/Linux Building blocks

In the Unix/Linux world, the most important of these are the system calls. When we run a program like strace to see exactly what a Unix/Linux process is doing, we are watching this fundamental interaction between a program and its requests to the kernel, usually for access to resources controlled by the operating system.



Requests

A Unix/Linux system call is a direct request to the kernel regarding a system resource. It might be a request for a file descriptor to manipulate a file, it might be a request to write to a file descriptor, or any of hundreds of possible operations.



Requests

These are exactly the tools that every Unix/Linux program is built upon. Even the simplest of all programs, /bin/true and /bin/false, which do nothing except provide a process return code, must make at least one system call to exit(2).



File systems as a mainstay

In some sense, the mainstay operations are those on the file system.



File systems as a "permanent" resource

Unlike many other resources which are just artifacts of the operating system and disappear at each reboot, changing a file system generally is an operation that has some permanence. Filesystems have a naming system (we call these "pathnames" in the Unix/Linux world), and it is possible to explicitly operate on elements of a filesystem via these names. generally expect these files and directories to persist even over reboots of a system.



File systems as a "permanent" resource

Of course it is possible and even common to create "RAM" disk filesystems since they are quite fast and for items that are meant to be temporary, they are quite acceptable. As I mentioned with MailScanner setup, it is recommended to put /var/spool/incoming on a RAM disk since the information that it holds is inherently transient.



File systems as a "permanent" resource

But in the everyday world of system administration, the rule is that filesystems have some permanence, and it's the exception for it to be merely a RAM disk.



Filesystem operations versus file descriptor operations

While many of the most frequent operations happen on file descriptors, file descriptors are indirect references, and with the exception of operations such as dup (2), are not created from other file descriptors.



Filesystem operations

Unix/Linux offers a number of direct functions on files referred to by their names inside of a filesystem. These are important operations, such as creating a file descriptor associated with a file of a given name, where we want to be able to refer to a file or directory by name.



Fundamental filesystem operations

```
open() -- create a new file descriptor to access a file
close() -- deallocate a file descriptor

access() -- returns a value indicating if a file is accessible
chmod() -- changes the permissions on a file in a filesystem
chown() -- changes the ownership of a file in a filesystem
```



Important filesystem operations

```
link() -- create a hard link to a file
symlink() -- create a soft link to a file
```



Important filesystem operations

```
mkdir() -- create a new directory
rmdir() -- remove a directory
```



Important filesystem operations



Important file descriptor calls

A file descriptor is an int. It provides stateful access to an i/o resource such as a file on a filesystem, a pseudo-terminal, or a socket to a tcp session.

```
open() -- create a new file descriptor to access a file close() -- deallocate a file descriptor
```



```
dup() -- duplicate a file descriptor
dup2() -- improved way to duplicate a file descriptor
```





Important file descriptor calls

These two are "kitchen sink" calls: they do a wide miscellanea of operations.



Important file descriptor calls

flock() -- lock/unlock a file associated with a file descriptor



```
pipe() -- create a one-way association between two file
    -- descriptors so that output from
    -- one goes to the input of the other
    -- very important for simple ipc such as shells
    -- provide between processes
```



```
select() -- multiplex on pending i/o to or from a set of file
    -- descriptors
```



```
read() -- send data to a file descriptor
write() -- take data from a file descriptor
fsync() -- forces a flush for a file descriptor
```







Signals

Unix also supports signals, a very simple IPC mechanism. These signals, while they do not carry a payload of information, are distinguishable by their number. The most important of these are SIGKILL, SIGTERM, SIGHUP, and SIGALRM, but there are usually many more, even ones reserved to user processes such as SIGUSR1 and SIGUSR2.



Arbitrary signalling

```
    kill -- send an arbitrary signal to an arbitrary process
    killpg -- send an arbitrary signal to all processes in a
    process group
```



Signals



Signals

```
    alarm -- set an alarm clock for a SIGALRM to be sent to -- a process time measured in seconds
    setitimer -- set an alarm clock in fractions of a second to -- deliver one of SIGALRM, SIGVTALRM, or SIGPROF
```



Signals



Modifying the current process's state

Many system calls exist to change the current process's state.

```
chdir -- change the working directory for a process to dirname chroot -- change the root filesystem for a process
```



Modifying the current process's state

execve -- execute another binary in this current process fork -- create a new child process running the same binary clone -- allows sharing of execution context (unlike fork(2)) exit -- terminate the current process



Modifying the current process's state

```
getdtablesize -- report how many file descriptors this process
    -- can have active simultaneously (see select()
    -- for why this is useful)
```



Finding the current process's state

```
getgid -- return the group id of this process
getuid -- return the user id of this process
getpgid -- return process group id of this process
getpgrp -- return process group's group of this process
```



Finding the current process's state



Modifying the current process's state

```
nice() -- change the calling process's priority
setpriority() -- arbitrarily change any process's (or group or user)
-- priority
getpriority() -- get any process's priorities
setrusage -- set maxima for resource utilization by the
-- current process
```



Communications and Networking

```
bind -- bind a file descriptor to an address, such as a tcp port
listen -- specify willingness for some number of connections to be
-- blocked waiting on accept()
accept -- block until there is a new connection

-- actively connect to listen()ing socket
```



Communications and Networking



Communications and Networking



Others

```
brk
            -- allocate memory for the data segment for the
            -- current process. in Linux, the first call to brk
            -- actually creates the heap; the second and subsequent
            -- calls do allocation
gethostname -- gets a ''canonical hostname'' for the machine
gettimeofday -- gets the time of day for the whole machine
settimeofday -- sets the time of day for the whole machine
mount
             -- attaches a filesystem to a directory and makes
             -- it available
             -- flushes all filesystem buffers, forcing changed
sync
             -- blocks to ''drives'' and updates superblocks
```

futex

-- raw locking (lets a process block waiting on a change to a specific memory location)



Others

```
-- provides direct access from the kernel to:

load average

total ram for system

available ram

amount of shared memory existing

amount of memory used by buffers

total swap space

swap space available

number of processes currently in proctable
```



SYS V IPC

```
msgctl -- SYS V messaging control (uid, gid, perms, size)
msgget -- SYS V message queue creation/access
msgrcv -- receive a SYS V message
msgsnd -- send a SYS V message

shmat -- attach memory location to SYS V shared memory segment
shmctl -- SYS V shared memory contrl (uid, gid, perms, size, etc)
shmget -- SYS V shared memory creation/access
shmdt -- detach from SYS V shared memory segment
```

