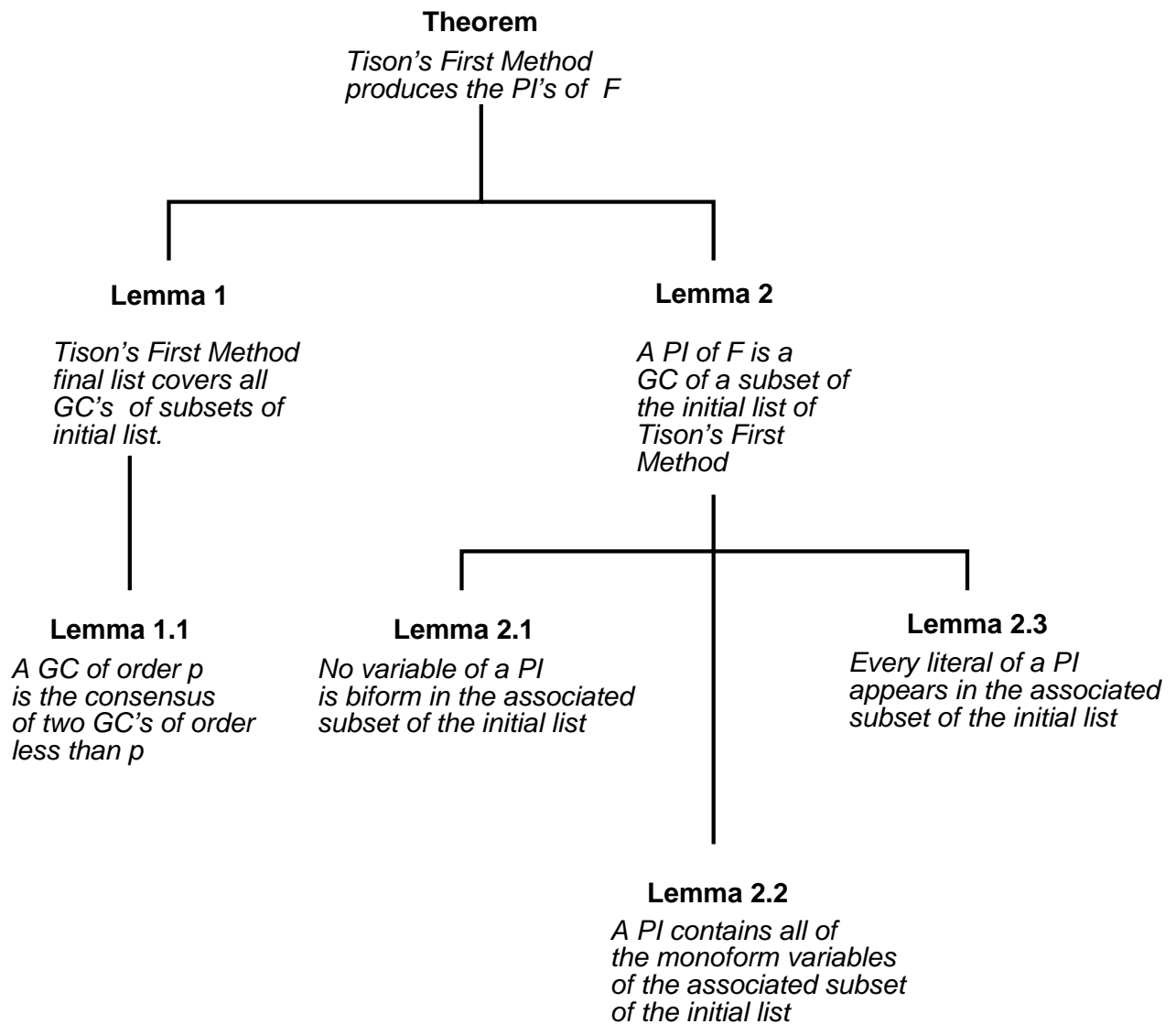


The next task is to prove the following theorem:

Theorem:(Correctness of Tison's First Method)

Let $\mathcal{F} = T_1 + T_2 + \dots + T_p$ be a sum-of-products expression for the switching function \mathcal{F} . If Tison's First Method is applied to the set of products $\mathcal{T} = (T_1, \dots, T_p)$ it produces a final set of products $L = (L_1, \dots, L_k)$ such that each L_i is a prime implicant of \mathcal{F} and any prime implicant of \mathcal{F} is equal to some L_j .

In order to do this we need a progression of 6 related lemmas. The relationships are shown in the following diagram.



Lemma 1.1:

Let Q_i $i = 1, \dots, p$ be products of literals and $X = GC(Q_1, \dots, Q_p)$ be the generalized consensus of order p of those products. There exists a variable x and two products R and S such that

$$\begin{aligned} R &= GC(Q_{i_1}, \dots, Q_{i_r}) \\ S &= GC(Q_{j_1}, \dots, Q_{j_s}) \\ r &\leq p \quad \text{and} \quad s \leq p \\ X &= cons(R, S, x) \end{aligned}$$

In fact, the proof of this lemma shows that such a consensus can be formed for each biform variable in (Q_1, \dots, Q_p) . Note that this is one part of the basic step of both of Tison's methods.

Proof of Lemma 1.1:

See Handout on Tison's Method

Lemma 1.1 allows us to prove the following result that forms a large part of the proof of the correctness theorem.

Lemma 1:

Let $\mathcal{T} = (T_1, \dots, T_n)$ be a set of products and let $X = GC(T_{i_1}, \dots, T_{i_k})$, $k \leq n$ be a generalized consensus of a subset of \mathcal{T} . If Tison's First method is applied to \mathcal{T} then there exists a product L_j such that

$$\begin{array}{l} L_j \in L = (L_1, \dots, L_m) \\ X \rightarrow L_j \end{array}$$

where L is the final list produced by Tison's First method.

In other words all generalized consensi of subsets of \mathcal{T} are covered by elements of L .

Proof of Lemma 1:

Without loss of generality, assume that $T_{i_j} = T_j, j = 1, \dots, k$, i.e., given X and i_1, \dots, i_k , renumber the products so that those involved in the generalized consensus appear first. Let x_1, x_2, \dots, x_b be the biform variables in order of use in generating consensus and let Stage i refer to computing all consensus w.r.t. x_i . Let t be the smallest integer such that all biform variables of (T_1, \dots, T_k) are contained in x_1, \dots, x_t . (Note that this implies x_t is biform in (T_1, \dots, T_k) .) We will use induction on t to show that at the end of stage t L has a term which covers X .

Assume $t = 1$. Then there is a single biform literal, x , in (T_1, \dots, T_k) . Therefore, B_i , the biform variable part of T_i must be either x or x' . (We will assume the literal x is also an uncomplemented switching variable. If it is not the following reasoning is easily modified substituting x' for x .) Since for $X = GC(T_1, \dots, T_k)$ to exist we must have $B_1 + \dots + B_k = 1$ irredundantly we must have $k = 2$. If $k = 1$ the sum is not always 1 and if $k \geq 3$ it cannot be irredundant since there are only two possible literals. So $k = 2$ and $X = RS$ where $T_1 = xR$ and $T_2 = x'S$.

On the first step of Tison's method, the consensus of $T_1 = xR$ and $T_2 = x'S$ w.r.t. x will be evaluated. Since it is known that R and S contain no other biform variables the definition of consensus is satisfied and RS is added to the list L . Since $RS = GC(T_1, \dots, T_k)$ this completes the case for $t = 1$.

Now suppose $t > 1$. Since x_t is biform in (T_1, \dots, T_k) by Lemma 1.1, we know that $X = \text{cons}(Y, Z, x_t)$, where Y is a generalized consensus of some subset, \mathcal{Y} , of terms in (T_1, \dots, T_k) and Z is a generalized consensus of another subset \mathcal{Z} .

Note, by definition, Y and Z are products of variables that are monofrom in \mathcal{Y} and \mathcal{Z} respectively. x_t must be monofrom in \mathcal{Y} and \mathcal{Z} and must appear as x_t in one (assume it is \mathcal{Y}) and x'_t in the other since X is formed by consensus w.r.t. x_t . So $X = Y_0 Z_0$ where $Y = x_t Y_0$ and $Z = x'_t Z_0$.

Since by definition of t , x_t is the last biform variable in (T_1, \dots, T_k) , all biform variables in \mathcal{Y} and \mathcal{Z} must be among x_1, \dots, x_{t-1} . By the induction hypothesis that all GC of sets whose biform variables are included in x_1, \dots, x_{t-1} have terms in L that cover them, two terms Y^* and Z^* exist in L such that

$$\begin{aligned} Y &\leq Y^* \\ Z &\leq Z^* \end{aligned}$$

There are only 3 possible cases:

If x_t not in Y^* , then Y^* contains only literals from Y_0 and therefore X is covered:

$$X = Y_0 Z_0 \leq Y_0 \leq Y^*$$

If x'_t not in Z^* , then Z^* contains only literals from Z_0 and therefore X is covered:

$$X = Y_0 Z_0 \leq Z_0 \leq Z^*$$

If x'_t in Z^* and x_t in Y^* then we have $Y = x_t Y_0$ and $Y^* = x_t Y_0^*$. Since $Y \leq Y^*$ it follows that $Y_0 \leq Y_0^*$. Similarly, $Z^* = x'_t Z_0^*$ and $Z_0 \leq Z_0^*$. Therefore,

$$X = Y_0 Z_0 \leq Y_0^* Z_0^* = \text{cons}(Y^*, Z^*, x_t) = X^*$$

But, X^* is produced by Tison's method when taking consensi w.r.t. x_t .

Once any product that covers X is placed in L it will not be removed except by another product that covers X . \square

We need three fairly simple lemmas to prove our second significant contribution to the proof of the correctness theorem. Note that \leq is used to denote the implication coverage relationship, i.e., $x \rightarrow y$ or y covers x is denoted $x \leq y$.

Lemma 2.1:

Suppose P is a prime implicant of a switching function \mathcal{F} and (Q_1, \dots, Q_k) is a set of products that are implicants of \mathcal{F} such that

$$P \leq Q_1 + Q_2 + \dots + Q_k \leq \mathcal{F}$$

irredundantly (no Q_i can be removed without destroying the coverage).

If u is a literal in P then u' is not in (Q_1, \dots, Q_k) , i.e., no variable in P is biform in (Q_1, \dots, Q_k) .

Proof of Lemma 2.1:

Suppose that $P = uR$ and, without loss of generality, $Q_1 = u'S$. Since $P \leq Q_1 + \dots + Q_k$, for any assignment such that $P = 1$ we have $P(Q_1 + \dots + Q_k)' = 0$. If $P = 0$ the product must also be 0 so we have for all assignments to the switching variables

$$\begin{aligned} 0 &= P(Q_1 + \dots + Q_k)' \\ &= PQ'_1 \dots Q'_k \\ &= uR(u + S')Q'_2 \dots Q'_k \\ &= (uR + uRS')Q'_2 \dots Q'_k \\ &= uRQ'_2 \dots Q'_k \\ &= PQ'_2 \dots Q'_k \end{aligned}$$

This says that $P \leq Q_2 + \dots + Q_k$ which contradicts the assumption of irredundancy. \square

Lemma 2.2:

Suppose P is a prime implicant of a switching function \mathcal{F} and (Q_1, \dots, Q_k) is a set of products that are implicants of \mathcal{F} such that

$$P \leq Q_1 + Q_2 + \dots + Q_k \leq \mathcal{F}$$

irredundantly (no Q_i can be removed without destroying the coverage).

If u is a literal that is monofom in (Q_1, \dots, Q_k) then u appears in P .

Proof of Lemma 2.2:

Suppose a monofom u appears in Q_1, \dots, Q_m but not in P, Q_{m+1}, \dots, Q_k . As above we have

$$P(Q_1 + \dots + Q_m + Q_{m+1} + \dots + Q_k)' = 0$$

for all assignments to the variables. Now consider all assignments such that $u = 0$. $Q_1 = \dots = Q_m = 0$ and P, Q_{m+1}, \dots, Q_k are unaffected, i.e., they take on the same set of values that they achieve for ALL ASSIGNMENTS. We therefore have for ALL ASSIGNMENTS

$$P(Q_{m+1} + \dots + Q_k)' = 0$$

which says that $P \leq Q_{m+1} + \dots + Q_k$. This contradicts the irredundancy assumption. \square

Lemma 2.3:

Suppose P is a prime implicant of a switching function \mathcal{F} and (Q_1, \dots, Q_k) is a set of products that are implicants of \mathcal{F} such that

$$P \leq Q_1 + Q_2 + \dots + Q_k \leq \mathcal{F}$$

irredundantly (no Q_i can be removed without destroying the coverage).

If u is a literal in P then u occurs in (Q_1, \dots, Q_k) .

Proof of Lemma 2.3: Suppose $P = uR$ and u is not in (Q_1, \dots, Q_k) . As above we have

$$P(Q_1 + \dots + Q_k)' = 0.$$

Set $u = 1$. Q_1, \dots, Q_k are unaffected so

$$R(Q_1 + \dots + Q_k)' = 0$$

for all assignments to the variables. It follows that

$$R \leq Q_1 + \dots + Q_k \leq \mathcal{F}$$

which contradicts the primeness of P . \square

Corollary 2.4:

Suppose P is a prime implicant of a switching function \mathcal{F} and (Q_1, \dots, Q_k) is a set of products that are implicants of \mathcal{F} such that

$$P \leq Q_1 + Q_2 + \dots + Q_k \leq \mathcal{F}$$

irredundantly (no Q_i can be removed without destroying the coverage).

Let S_Q, S_m , and S_b be the sets of all literals, all monoform literals, and all biform literals in (Q_1, \dots, Q_k) , respectively. If S_P is the set of literals in the product P , then $S_m = S_P$.

Proof: By Lemma 2.3, $S_P \subseteq S_Q = S_b \cup S_m$. By Lemma 2.1, $S_P \cap S_b = \emptyset$. Therefore, $S_P \subseteq S_m$. By Lemma 2.2, $S_m \subseteq S_P$. Therefore, $S_P = S_m$. \square

We can now state the following which associates a prime implicant with a generalized consensus.

Lemma 2:

Suppose P is a prime implicant of a switching function \mathcal{F} and $\mathcal{Q} = (Q_1, \dots, Q_n)$ is a set of products that are implicants of \mathcal{F} such that

$$\mathcal{F} = Q_1 + Q_2 + \dots + Q_n.$$

Then there exists a subset of \mathcal{Q} with k implicants and $k \leq n$, such that

$$P = GC(Q_{i_1}, \dots, Q_{i_k})$$

Proof of Lemma 2: By the assumptions of the lemma,

$$P \leq Q_1 + Q_2 + \dots + Q_n = \mathcal{F}.$$

Since $\mathcal{F} = 1$ does not necessarily mean that $P = 1$ there may be redundant implicants Q_j on the right side of the inequality. These can be removed while preserving the coverage of P by the resulting sum. Since P is prime none of its literals can be removed without violating the implication of \mathcal{F} . So after removing the Q_j which are redundant with respect to coverage of P we have

$$P \leq Q_{i_1} + \dots + Q_{i_k} \leq \mathcal{F}.$$

By the summary of Lemmas 2.1, 2.2, and 2.3 expressed in Corollary 2.4, the literals of P are the monofom literals of $(Q_{i_1}, \dots, Q_{i_k})$. Write $Q_{i_j} = B_{i_j} X_{i_j}$ by splitting the literals into biform and monofom respectively. Therefore, $P = X_{i_1} \dots X_{i_k}$ and we must consider $B_{i_1} + \dots + B_{i_k}$.

Choose an assignment of the variables so that all monomial literals evaluate to 1. Since $P \leq Q_{i_1} + \dots + Q_{i_k}$ it follows that FOR THIS ASSIGNMENT

$$\begin{aligned} 1 &= Q_{i_1} + \dots + Q_{i_k} \\ &= B_{i_1} X_{i_1} + \dots + B_{i_k} X_{i_k} \\ &= B_{i_1} + \dots + B_{i_k} \end{aligned}$$

Recall, however, that by the definitions of B_j and X_j there are no literals in common between P and any B_j so setting the assignment above does not affect the values of literals in any B_j . Therefore,

$$B_{i_1} + \dots + B_{i_k} = 1$$

FOR ALL ASSIGNMENTS of variables in $(Q_{i_1}, \dots, Q_{i_k})$.

The irredundancy of the sum follows from the irredundancy of $P \leq Q_{i_1} + \dots + Q_{i_k}$. If, say, B_{i_1} was a redundant term then when $P = 1$

$$\begin{aligned} 1 &= B_{i_2} + \dots + B_{i_k} \\ &= B_{i_2}X_{i_2} + \dots + B_{i_k}X_{i_k} \\ &= Q_{i_2} + \dots + Q_{i_k} \end{aligned}$$

which is a contradiction that $P \leq Q_{i_1} + \dots + Q_{i_k}$ irredundantly (Q_{i_1} could have been removed).

Therefore,

$$P = GC(Q_{i_1}, \dots, Q_{i_k})$$

□

We now have all of the theoretical groundwork to state and prove the correctness theorem.

Theorem:(Correctness of Tison's First Method)

Let $\mathcal{F} = T_1 + T_2 + \dots + T_p$ be a sum-of-products expression for the switching function \mathcal{F} . If Tison's First Method is applied to the set of products $\mathcal{T} = (T_1, \dots, T_p)$ it produces a final set of products $L = (L_1, \dots, L_k)$ such that each L_i is a prime implicant of \mathcal{F} and any prime implicant of \mathcal{F} is equal to some L_j .

Proof of Correctness Theorem: The proof consists of two parts:

1. Tison's First Method generates all prime implicants.
2. All implicants in the final list L are prime implicants.

The first part results directly from Lemma 2. For any prime implicant P we have

$$P \leq T_1 + \dots + T_p = \mathcal{F}.$$

By Lemma 2, there must be a subset of terms of this sum such that

$$P \leq T_{i_1} + \dots + T_{i_k} \leq \mathcal{F}$$

with $P = GC(T_{i_1}, \dots, T_{i_k})$.

By Lemma 1, there must exist an $L_j \in L$ such that $P \leq L_j$. Since L_j is built by a series of consensus operations from an initial set of implicants of \mathcal{F} it must also be an implicant. Since P is prime, it follows that $P = L_j$, i.e., P cannot be covered by another implicant that is not prime.

The second part follows from the first part. Suppose L contains all of the prime implicants and one implicant Q which is not prime. Since Q is not prime some of its literals can be removed while preserving the implication property. Denote the product of the redundant literals by R . It follows that $Q = PR$ where

$$P \leq PR = Q \leq \mathcal{F}$$

Since no more literals can be removed P must be a prime implicant and by the first part of the proof $P = L_j$ for some j . Therefore, Q would be removed from L by the fact that it is covered by L_j . So only prime implicants appear in the final list L . \square