## COP5621 Exam 2-Spring 2005

Name: $\qquad$ (Please print)
Put the answers on these sheets. Use additional sheets when necessary. Show how you derived your answer when applicable (this is required for full cred it and helpful for partial credit). You can collect 100 points in total for this exam. A bonus question for an additional 15 points is included. If the total number of points for this exam exceed 100, the excess points are carried over to the next exams.

1. Match the terms below with the given sentences so as to best complete each sentence. Use no term more than once. Some terms will go unused. (10 points)

| (1) | context-free | $(7)$ | leftmost |
| :--- | :--- | ---: | :--- |
| $(2)$ | terminal | $(8)$ | rightmost |
| $(3)$ | LL(1) | $(9)$ | LR(1) |
| $(4)$ | SLR | $(10)$ | LALR(1) |
| $(5)$ | lexeme | $(11)$ | string |
| $(6)$ | nonterminal | $(12)$ | context-sensitive |

(a) In a $\qquad$ derivation, the rightmost nonterminal in a sentential form is replaced in each step.
(b) Recursive-descent parsing can be used for $\qquad$ grammars (list all that apply).
(c) A shift-reduce conflict may arise in constructing a $\qquad$ parsing table (list all that apply).
(d) In a context-free grammar, each production has a single $\qquad$ on the left-hand side.
(e) The yield of a parse tree is a sequence of $\qquad$ _s.
2. Name the four error recovery strategies. (10 points)
3. Show that the following grammar is ambiguous by constructing two distinct leftmost derivations for the string abab. Also show the two distinct parse trees. (15 points)

$$
S \rightarrow \mathbf{a} S \mathbf{b} S|\mathbf{b} S \mathbf{a} S| \varepsilon
$$

4. Consider the following grammar:

$$
\begin{aligned}
& \text { start } \rightarrow \text { label stmt } \\
& \text { label } \rightarrow \text { num } \\
& \mid \varepsilon \\
& \text { stmt } \rightarrow \text { id }:=\text { expr } \\
& \mid \text { goto num } \\
& \text { expr } \rightarrow \text { id } \\
& \mid \\
& \text { num }
\end{aligned}
$$

For each production $A \rightarrow \alpha$, determine $\operatorname{FIRST}(\alpha)$ and $\operatorname{FOLLOW}(A)$. (15 points)
5. Disprove that the following grammar is LL(1). (15 points).

$$
\begin{aligned}
& S \rightarrow A \mathbf{a} \\
& A \rightarrow \mathbf{b} \mid B \mathbf{b} \mathbf{c} \\
& B \rightarrow \mathbf{c} \mid \varepsilon
\end{aligned}
$$

6. Eliminate left recursion from the following grammar. Use the algorithm described in the book and illustrated in class. (15 points)

$$
\begin{array}{lll}
A & \rightarrow & B A \mid \mathbf{a} \\
B & \rightarrow & A B \mid \mathbf{b}
\end{array}
$$

7. Given below is the collection of $\mathrm{LR}(0)$ items for the grammar

$$
\left.\begin{array}{ll}
\text { 1. } & L \\
\text { 2. } & L
\end{array}\right] L \| L
$$

Construct an SLR parsing table. There are shift-reduce conflicts. You must resolve the shiftreduce conflicts by choosing either a shift or reduce depending on the following assumptions:

- The operators $\|, \& \&,!$ are listed in increasing order of precedence,
- The operators $\|$ and $\& \&$ are left associative, i.e. $\mathrm{b} \& \& \mathrm{~b} \& \& \mathrm{~b}=(\mathrm{b} \& \& \mathrm{~b}) \& \& \mathrm{~b}$.

Indicate your choice for each conflict by circling the shift or reduce. (20 points)

8. (bonus question). Given the following grammar, construct the $\mathrm{LR}(1)$ collection of sets of items with goto transitions. (15 points)

$$
\begin{aligned}
& S \quad \rightarrow \quad X Y \\
& X \rightarrow Y \mathbf{b} \\
& X \rightarrow \mathbf{b} \\
& Y \rightarrow \mathbf{c} \\
& Y \rightarrow \epsilon
\end{aligned}
$$

To get started, the first set of items $I_{0}$ is

$$
\begin{array}{lll}
I_{0}: & S^{\prime} \rightarrow \cdot S, \$ & \operatorname{goto}(S)=I_{1} \\
S \rightarrow \cdot X Y, \$ & \operatorname{goto}(X)=I_{2} \\
X & \rightarrow \cdot Y \mathbf{b}, \$ / \mathbf{c} & \operatorname{goto}(Y)=I_{3} \\
X & \rightarrow \cdot \mathbf{b}, \$ / \mathbf{c} & \operatorname{goto}(\mathbf{b})=I_{4} \\
X & \rightarrow \mathbf{c}, \mathbf{b} & \operatorname{goto}(\mathbf{c})=I_{5} \\
Y & \rightarrow, \mathbf{b} &
\end{array}
$$

