

Rate Monotonic Analysis

Introduction

Periodic tasks

Extending basic theory

Synchronization and priority inversion

Aperiodic servers

Case study: BSY-1 Trainer



Purpose of Tutorial

- Introduce rate monotonic analysis
- Explain how to perform the analysis
- Give some examples of usage
- **Convince you it is useful**



Tutorial Format

Lecture

Group exercises

Case study

Questions welcome anytime



RMARTS Project

Originally called Real-Time Scheduling in Ada Project (RTSIA).

- focused on rate monotonic scheduling theory
- recognized strength of theory was in analysis

Rate Monotonic Analysis for Real-Time Systems (RMARTS)

- focused on analysis supported by (RMS) theory
- analysis of designs regardless of language or scheduling approach used

Project focused initially on uniprocessor systems.

Work continues in distributed processing systems.



Real-Time Systems

Timing requirements

meeting deadlines

Periodic and aperiodic tasks

Shared resources

Interrupts



What's Important in Real-Time

Criteria for real-time systems differ from that for timesharing systems.

	Time-Sharing Systems	Real-Time Systems
Capacity	High throughput	Schedulability
Responsiveness	Fast average response	Ensured worst- case latency
Overload	Fairness	Stability

- schedulability is the ability of tasks to meet all hard deadlines
- *latency* is the worst-case system response time to events
- stability in overload means the system meets critical deadlines even if all deadlines cannot be met



Scheduling Policies

CPU scheduling policy: a rule to select task to run next

- cyclic executive
- rate monotonic/deadline monotonic
- earliest deadline first
- least laxity first

Assume preemptive, priority scheduling of tasks

• analyze effects of non-preemption later



Rate Monotonic Scheduling (RMS)

Priorities of periodic tasks are based on their rates: highest rate gets highest priority.

Theoretical basis

- optimal fixed scheduling policy (when deadlines are at end of period)
- analytic formulas to check schedulability

Must distinguish between scheduling and analysis

- rate monotonic scheduling forms the basis for rate monotonic analysis
- however, we consider later how to analyze systems in which rate monotonic scheduling is not used
- any scheduling approach may be used, but all realtime systems should be analyzed for timing



Rate Monotonic Analysis (RMA)

Rate monotonic analysis is a method for analyzing sets of real-time tasks.

Basic theory applies only to independent, periodic tasks, but has been extended to address

- priority inversion
- task interactions
- aperiodic tasks

Focus is on RMA, not RMS.



Why Are Deadlines Missed?

For a given task, consider

- *preemption*: time waiting for higher priority tasks
- execution: time to do its own work
- *blocking*: time delayed by lower priority tasks

The task is *schedulable* if the sum of its preemption, execution, and blocking is less than its deadline.

Focus: identify the biggest hits among the three and reduce, as needed, to achieve schedulability



Rate Monotonic Theory - Experience

IBM Systems Integration Division delivered a "schedulable" real-time network.

Theory used successfully to improve performance of IBM BSY-1 Trainer.

Incorporated into IEEE FutureBus+ standard

Adopted by NASA Space Station Program

European Space Agency requires as baseline theory.

Supported in part by Ada vendors



Rate Monotonic Analysis - Products

Journal articles (e.g., IEEE Computer, Hot Topics)

Videotape from SEI

Courses from Telos and Tri-Pacific

A Practitioner's Handbook for Real-Time Analysis: Guide to Rate Monotonic Analysis for Real-Time Systems from Kluwer

CASE tools from Introspect and Tri-Pacific

Operating systems and runtimes from Alsys, DDC-I, Lynx, Sun, Verdix and Wind River

Standards: Futurebus+, POSIX, Ada 9X



Summary

Real-time goals are: fast response, guaranteed deadlines, and stability in overload.

Any scheduling approach may be used, but all real-time systems should be analyzed for timing.

Rate monotonic analysis

- based on rate monotonic scheduling theory
- analytic formulas to determine schedulability
- framework for reasoning about system timing behavior
- separation of timing and functional concerns

Provides an engineering basis for designing real-time systems



Plan for Tutorial

Present basic theory for periodic task sets

Extend basic theory to include

- context switch overhead
- preperiod deadlines
- interrupts

Consider task interactions:

- priority inversion
- synchronization protocols (time allowing)
- Extend theory to aperiodic tasks:
 - sporadic servers (time allowing)
- **Present BSY-1 Trainer case study**



A Sample Problem



 τ_2 's deadline is 20 msec before the end of each period