

CIS 4360, SPRING 2026

DIGITAL SIGNATURE

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The slides are loosely based on those of
Prof. Mihir Bellare, UC San Diego.

Agenda

1. High-level Overview

2. Building Signature Scheme

3. Application: DNSSEC

The Need For Signing Is Ubiquitous

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CHECK N° 0007
DATE: Aug. 11, 2019

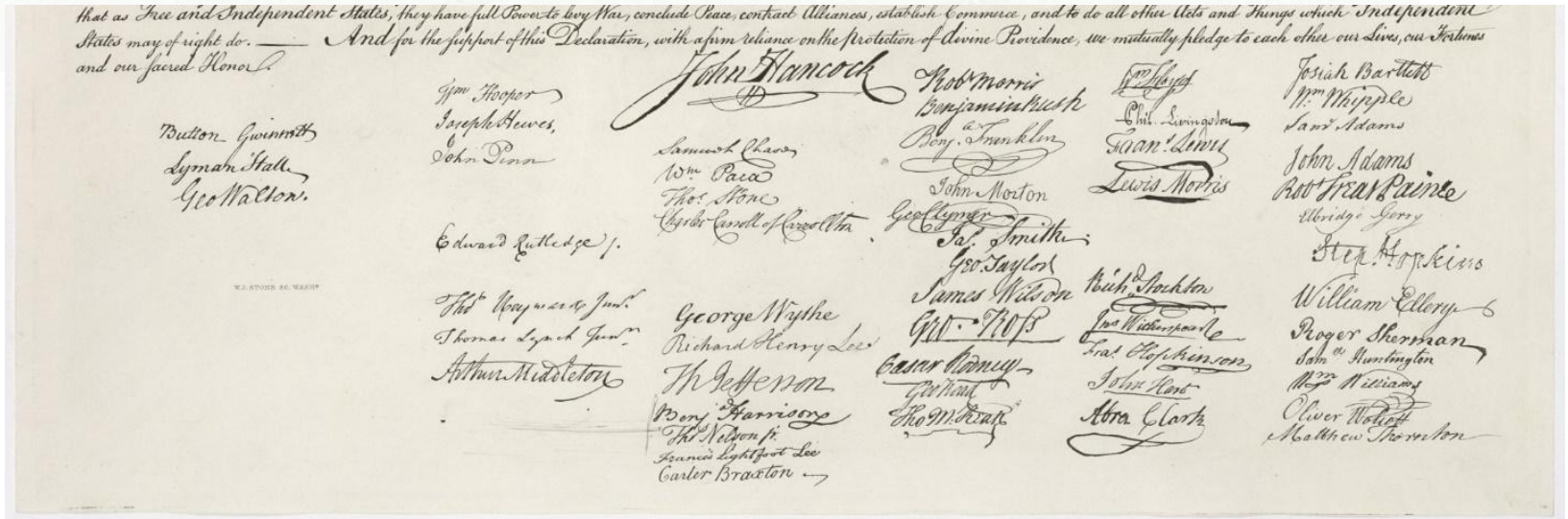
PAY TO THE ORDER OF: Mary Johnson \$ 715,39
Seven hundred fifteen and ³⁹/₁₀₀ DOLLARS

PAYABLE AT
ALL LOREM BANK BRANCHES IN USA
ACCOUNT N° 001234567

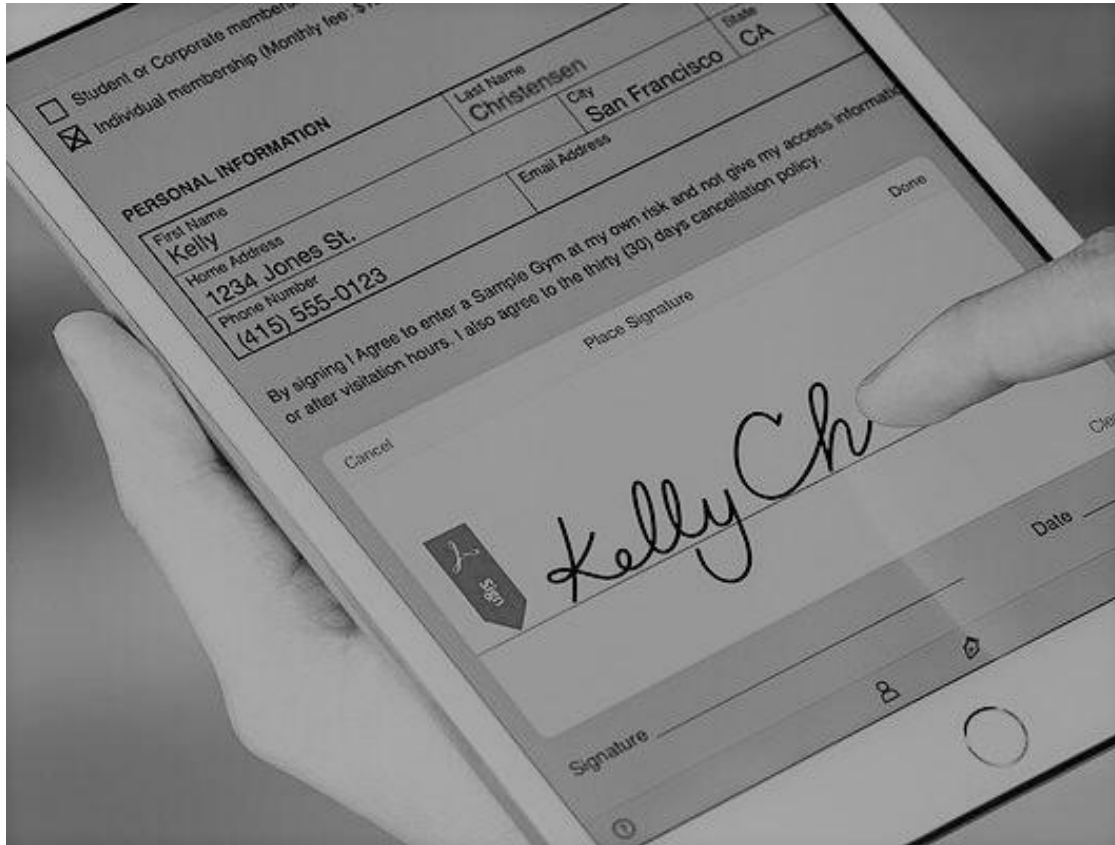
MEMO Monthly rent

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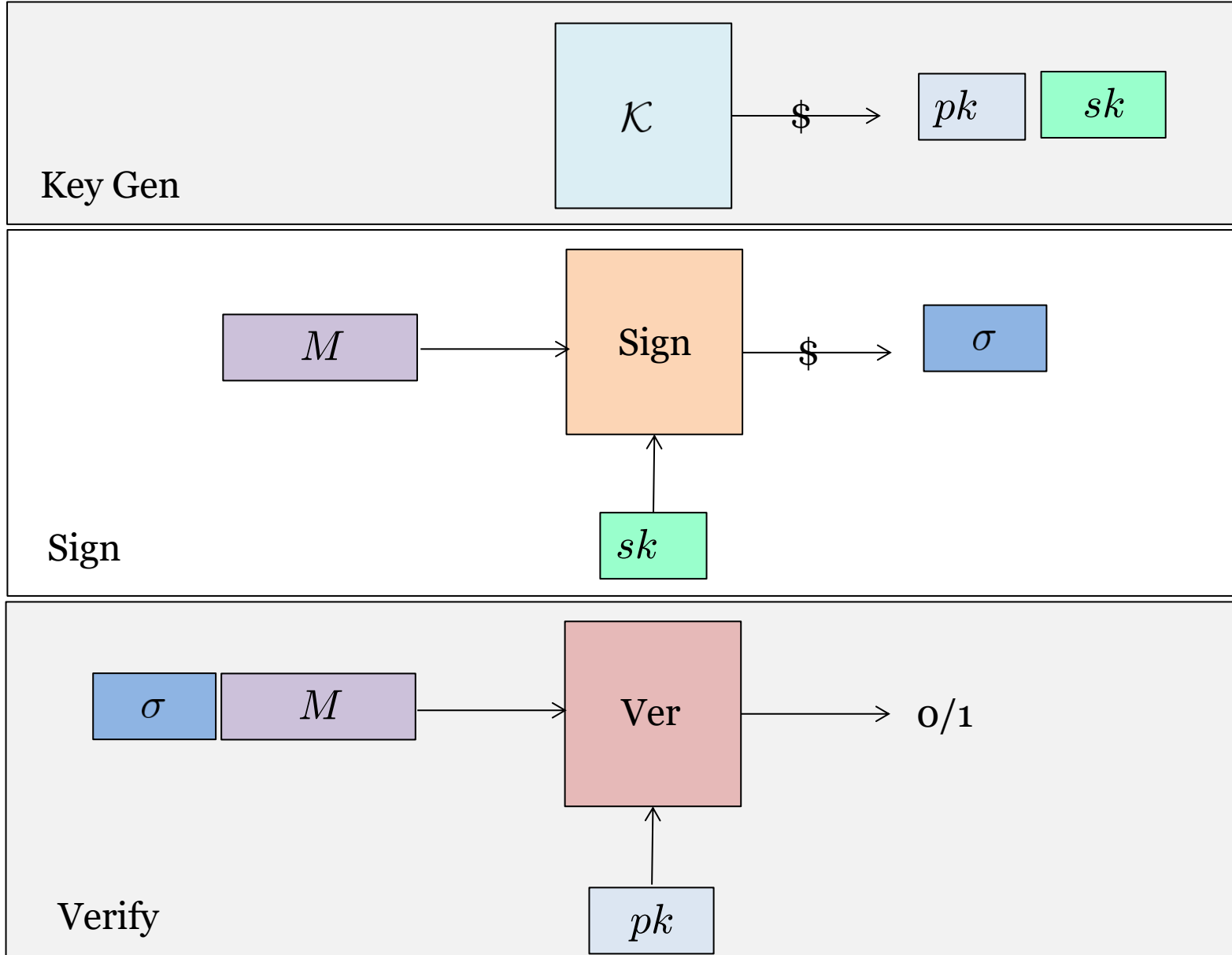
How To Sign Electronically?



Lots of apps to digitize
signatures

Problem: A digitized signature is easily copied → forgery

Digital Signature Scheme: Syntax



Digital Signature versus MAC

MAC

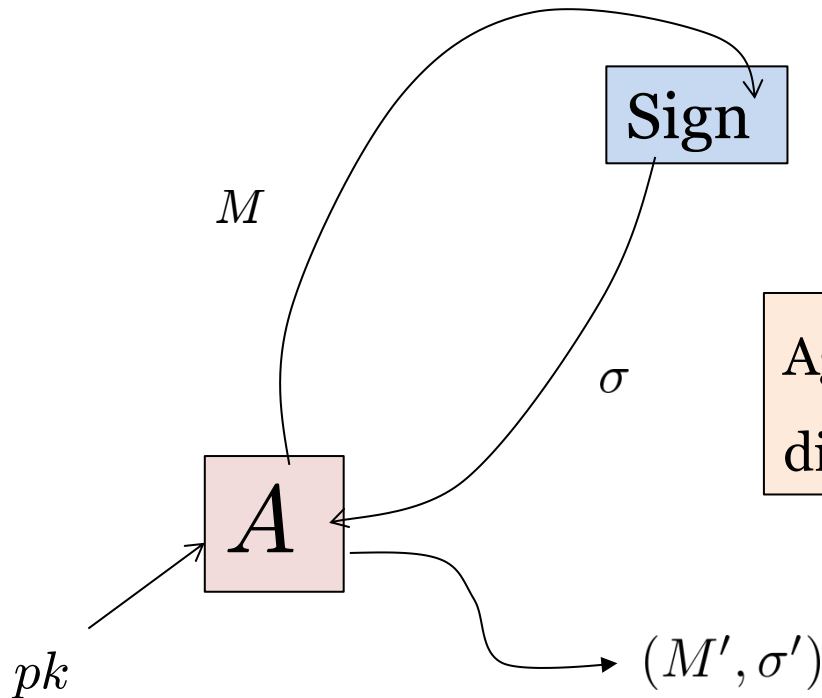
- Verifier needs to share a secret key with signer
- Verifier can impersonate signer

Digital Signature

- Verifier needs no secret
- Verifier cannot impersonate signer

Digital Signature: Unforgeability Security

- Similar to MAC security
- **Difference:** The adversary is given the public key



Again, digital signature doesn't directly thwart replay attack.

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A Bad Scheme: Plain RSA Signature

Key generation: Like RSA encryption

Sign:

- To sign a message, “decrypt” it:

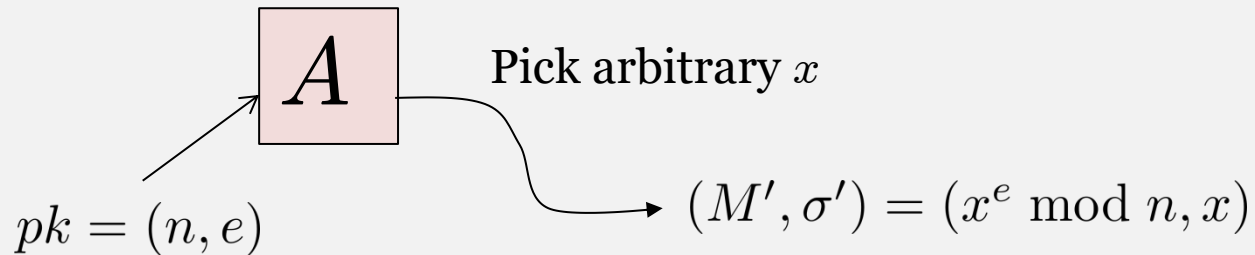
Verify:

- To verify a signature, “encrypt” it and compare with the message

Issues with Plain RSA Signature

- **Feasibility:** Can sign only short messages
- **Security:** Can easily break unforgeability security

No sign query needed!

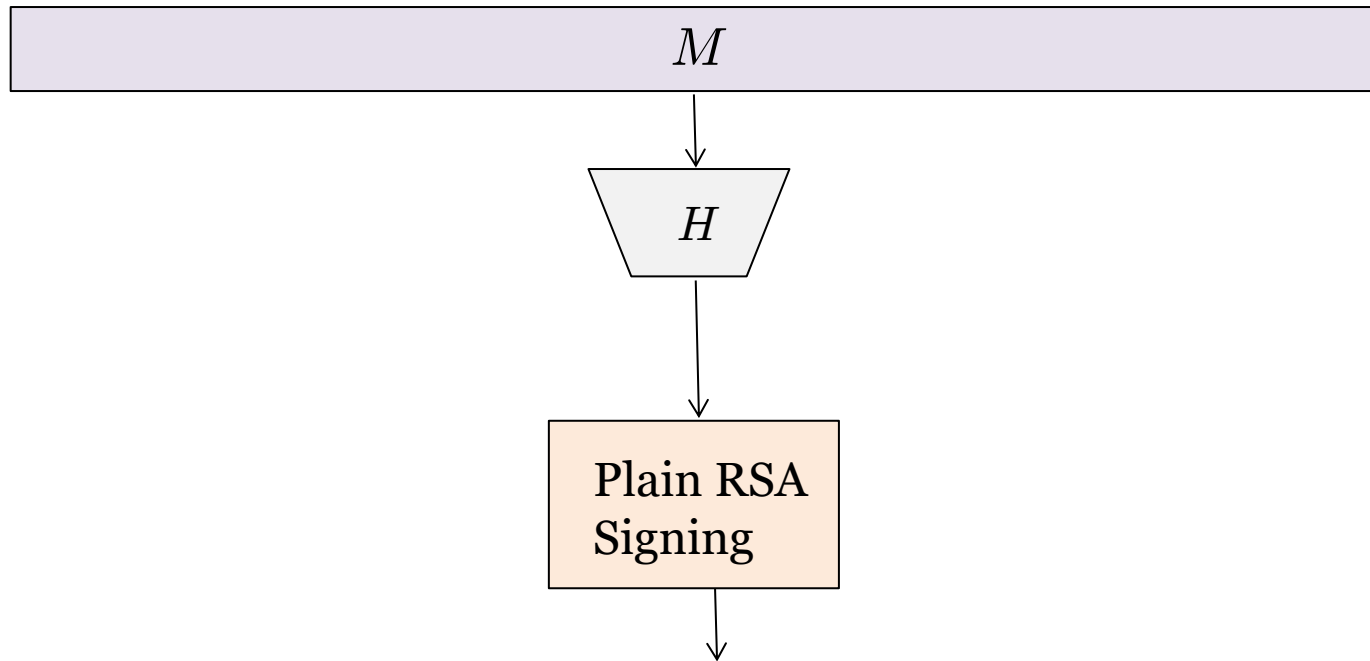


Hash-then-Sign Paradigm

Plain RSA Signature → Full Domain Hash (FDH)

Key generation: Like Plain RSA

Sign: To sign message M



Question: How to verify?

Security Requirement for Hash Function

What intuition suggests: Hash must be collision-resistant

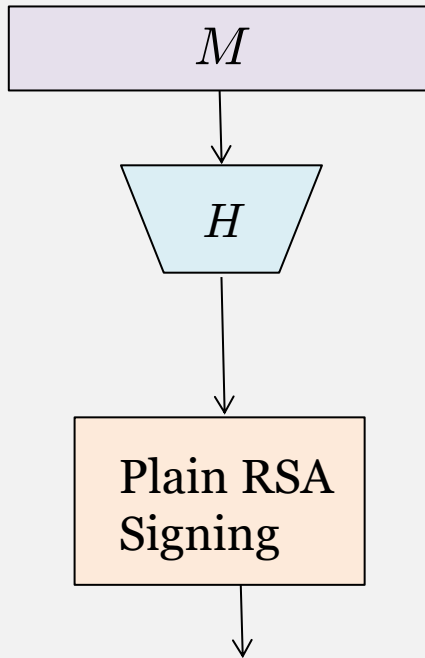
If $H(M) = H(M')$ then M and M' have the same signature

What proof requires: Hash is modeled as a random oracle

A Gap of **Demand** and **Supply**

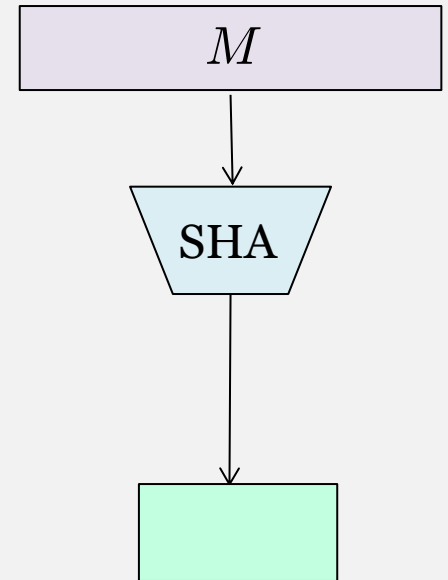
2048

bits of output

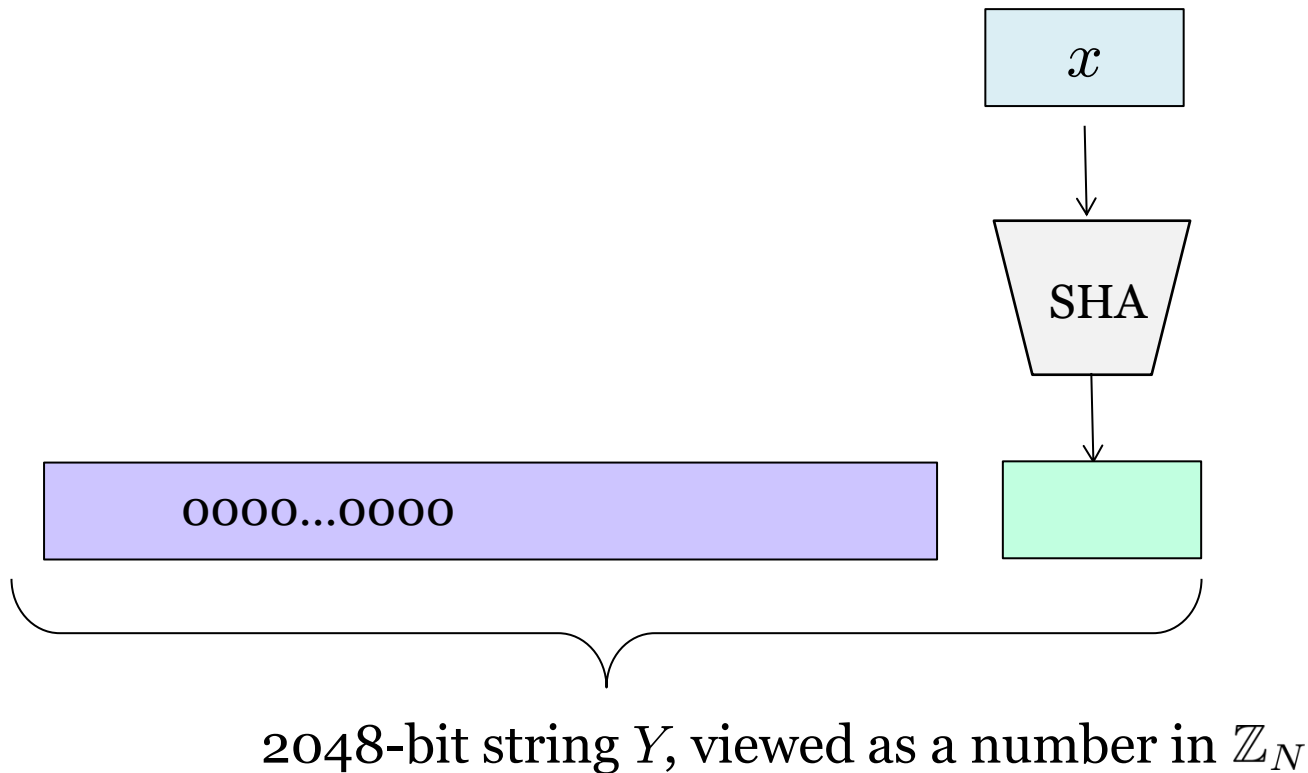


512

bits of output

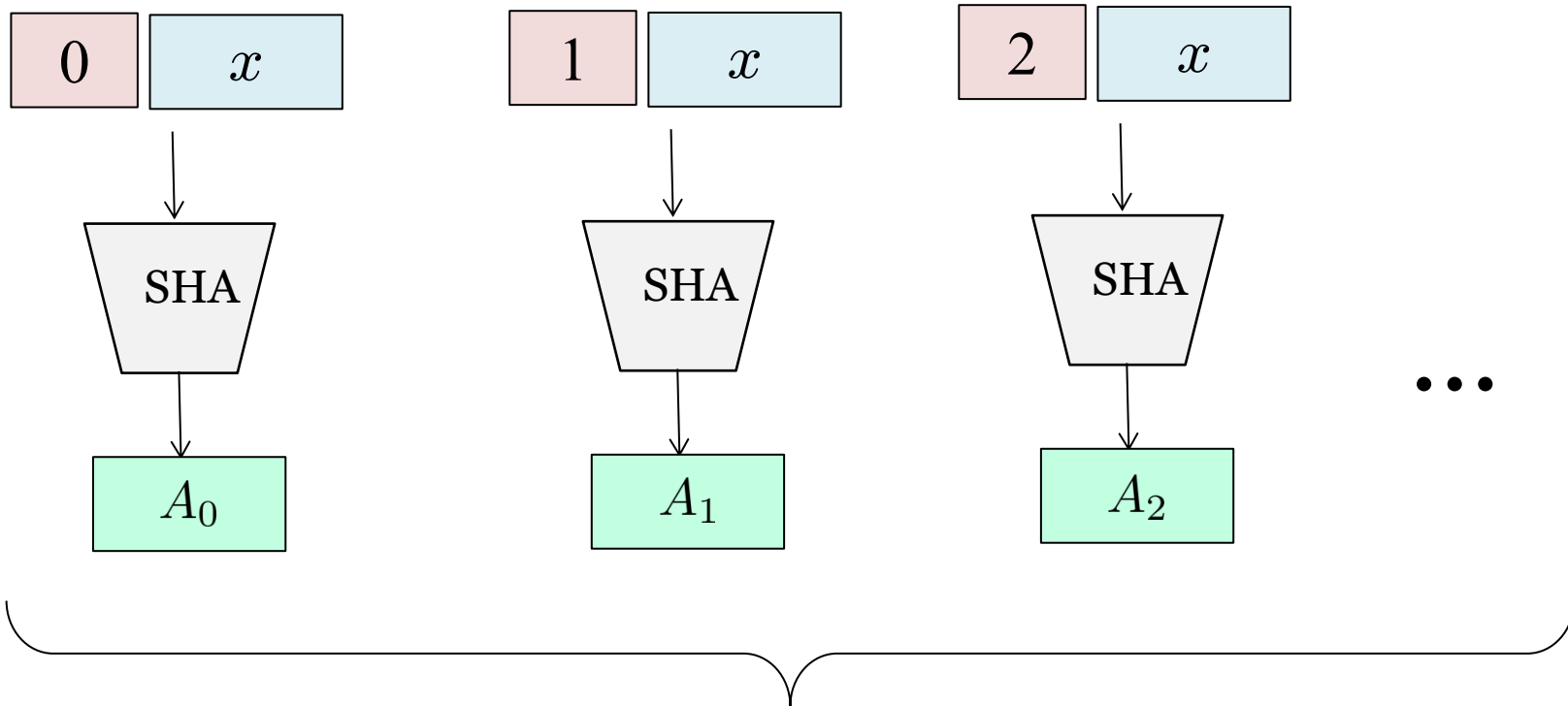


A Common Wrong Way to Hash



Broken by Desmedt and Odlyzko in 1985

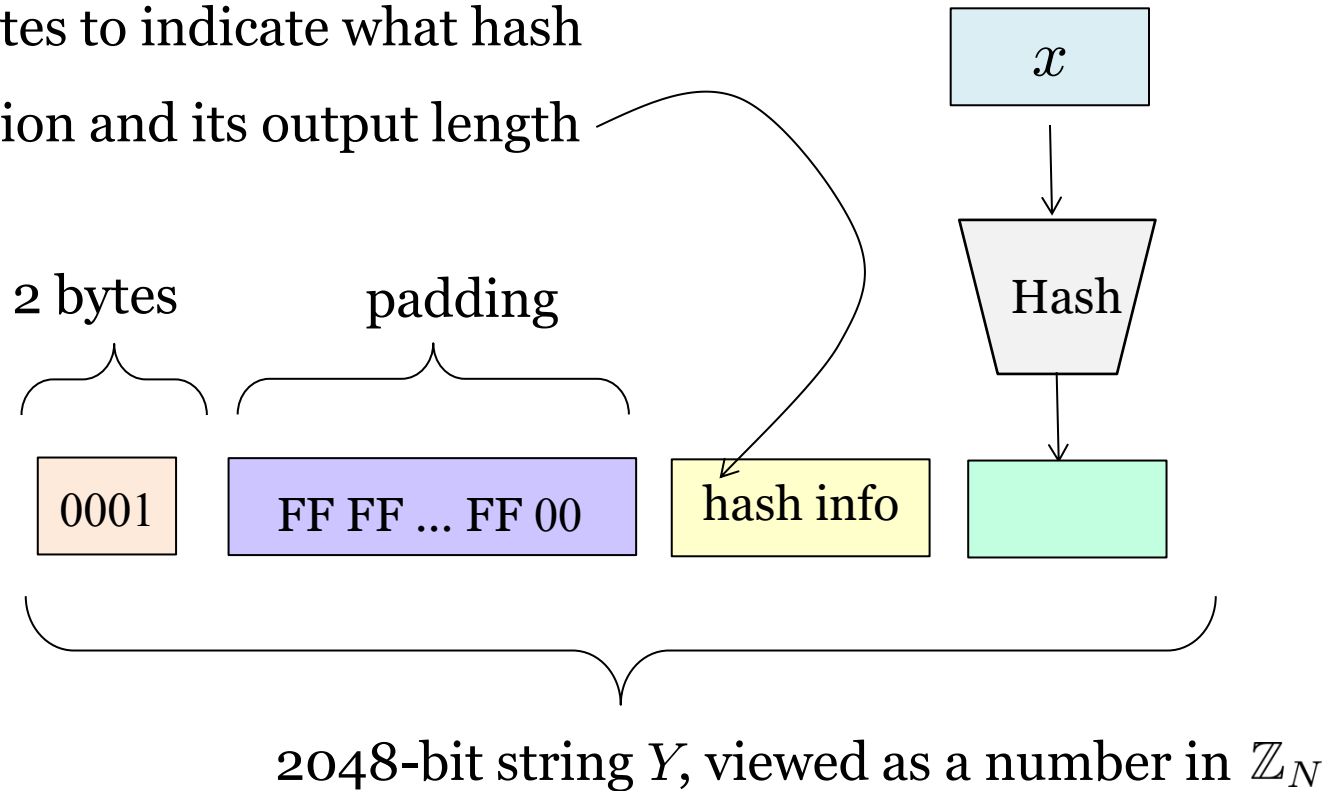
How to Hash Properly



Use the first $m = \lceil \log_2(N) \rceil$ bits and take mod N

Hashing in PKCS#1

19 bytes to indicate what hash function and its output length



Agenda

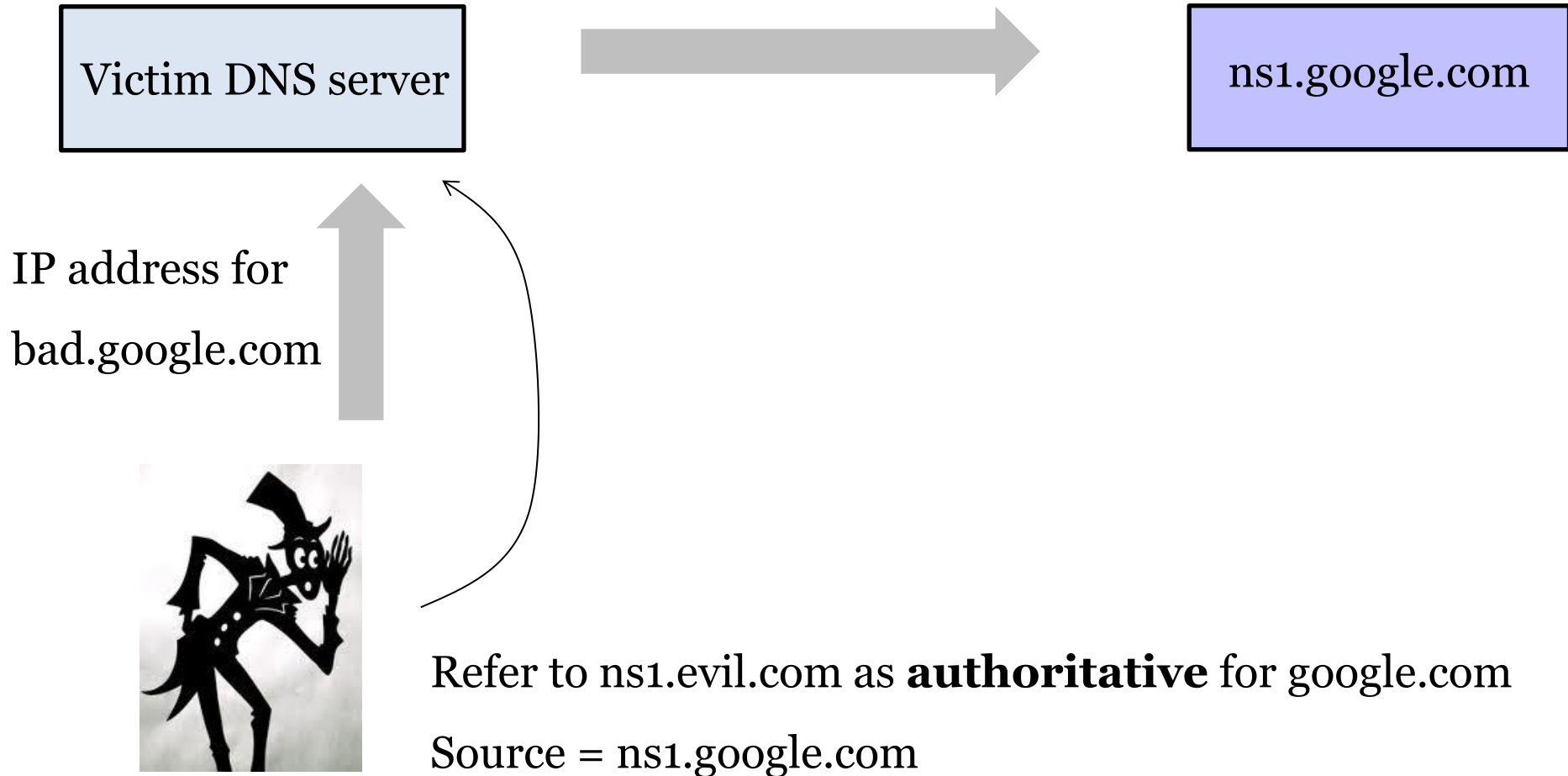
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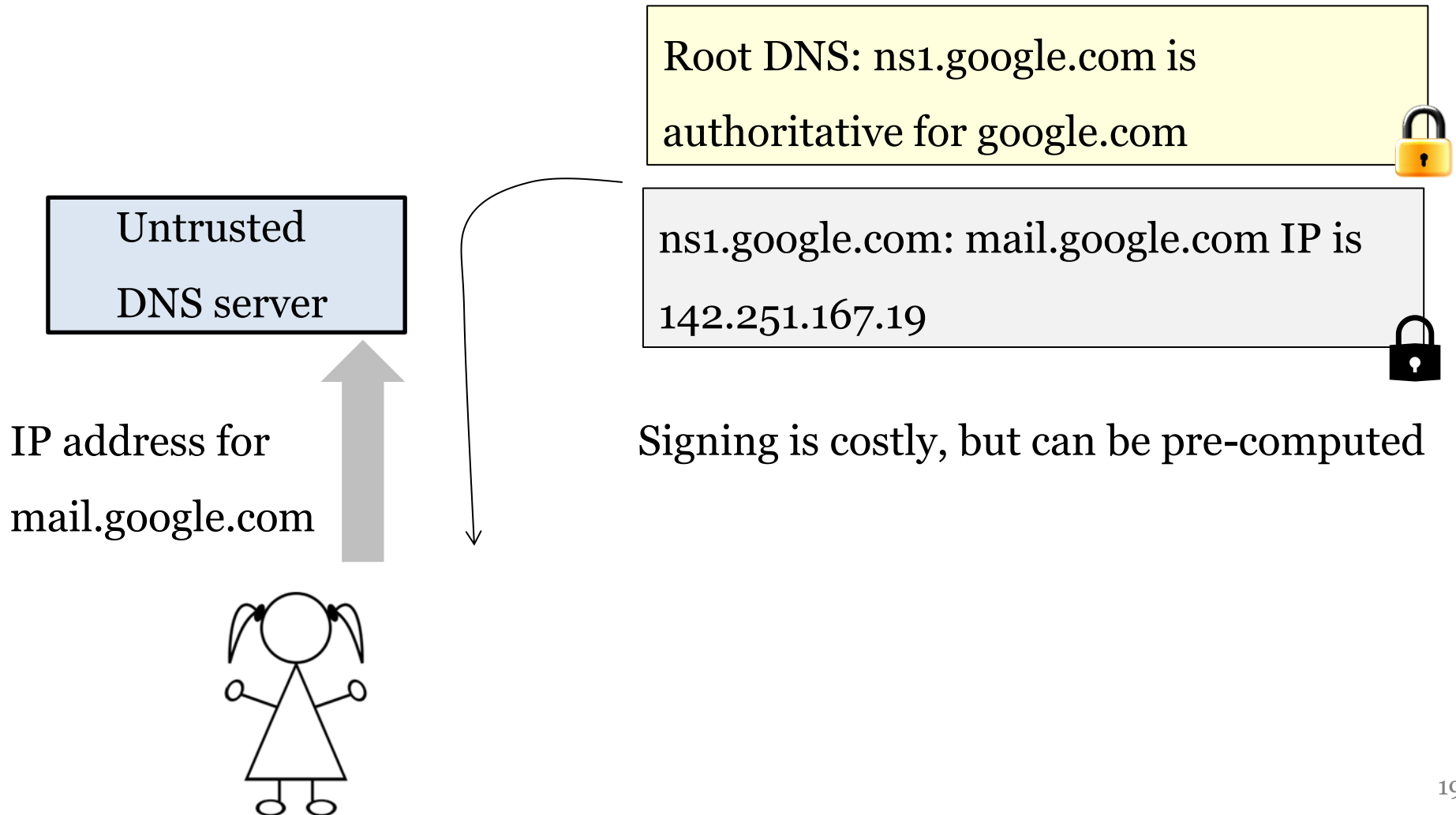
Recap: DNS Cache Poisoning Attack

Kaminsky, 2008

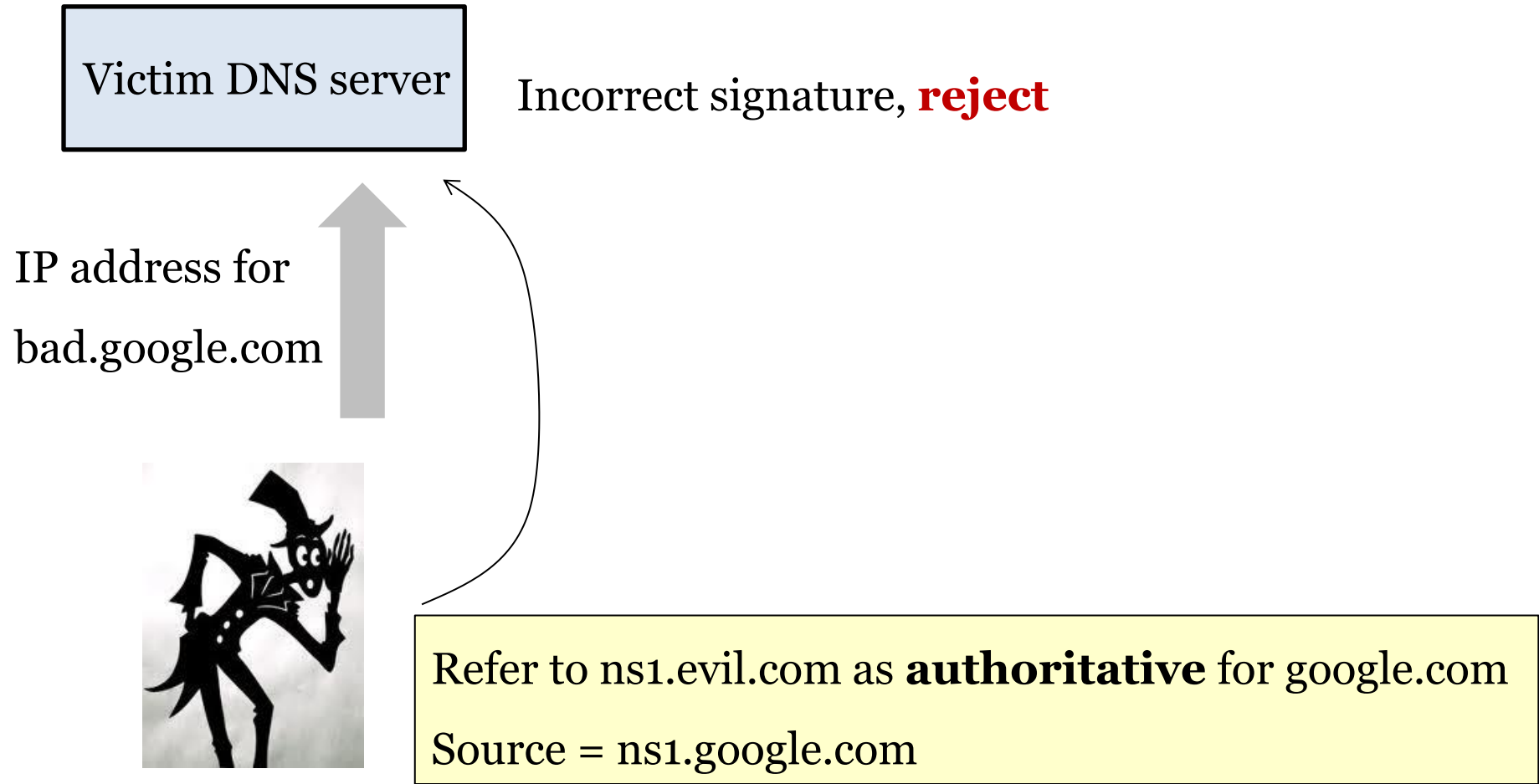


The Fix: DNSSEC

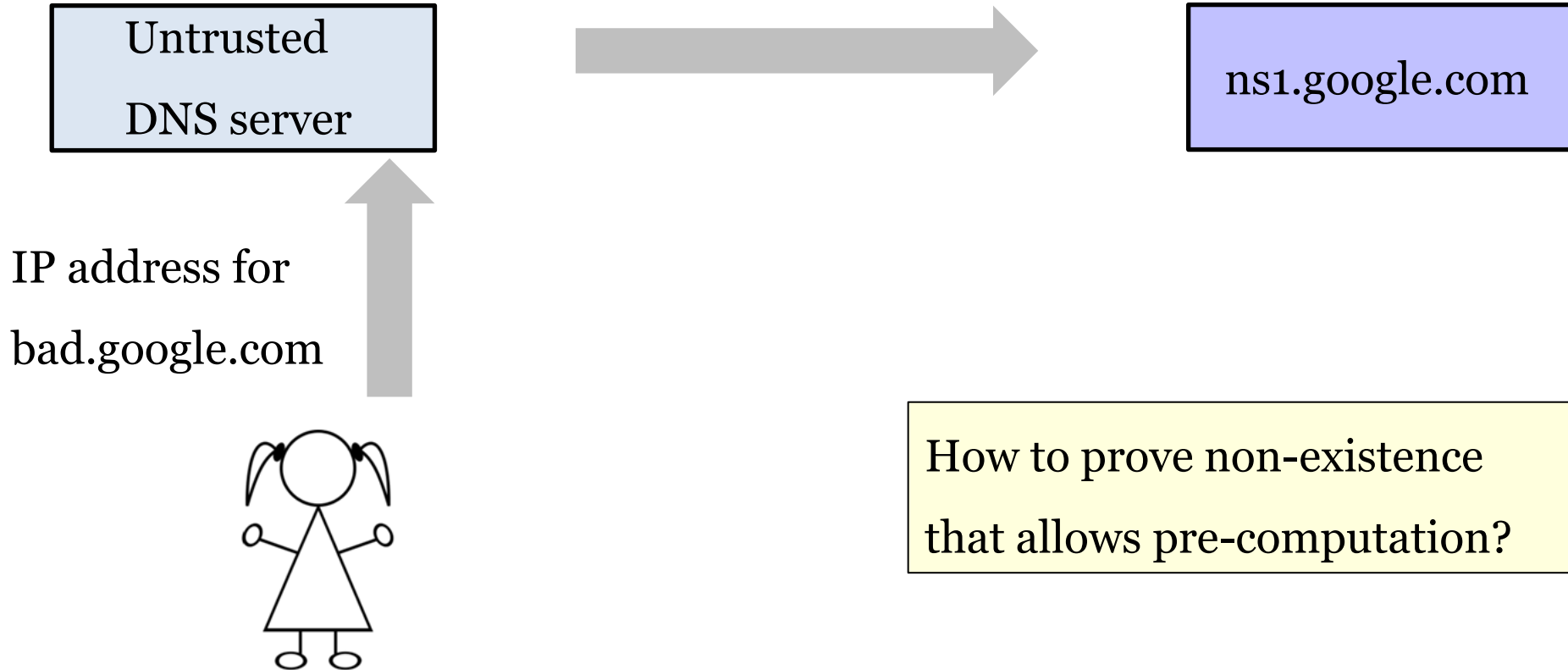
DNS replies need to be signed by authority



Thwarting Cache Poisoning Attack

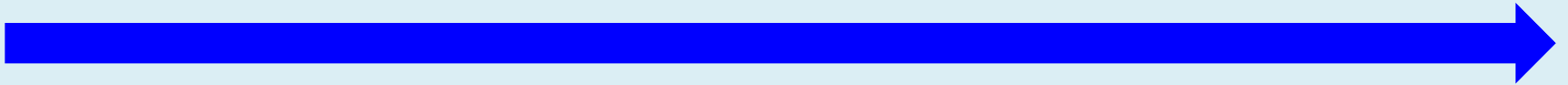


Issue: Prove Non-Existence



Proving Non-Existence: Precomputation


Google sorts its subdomain names alphabetically



account chrome mail policies site

Sign every consecutive pairs

google.com	account .google.com	
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account.google.com	chrome .google.com	
--------------------	--------------------	---

chrome.google.com	mail .google.com	
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Proving Non-Existence: Respond to Query

Unsuccessful Binary Search

