Hash Functions

Viet Tung Hoang

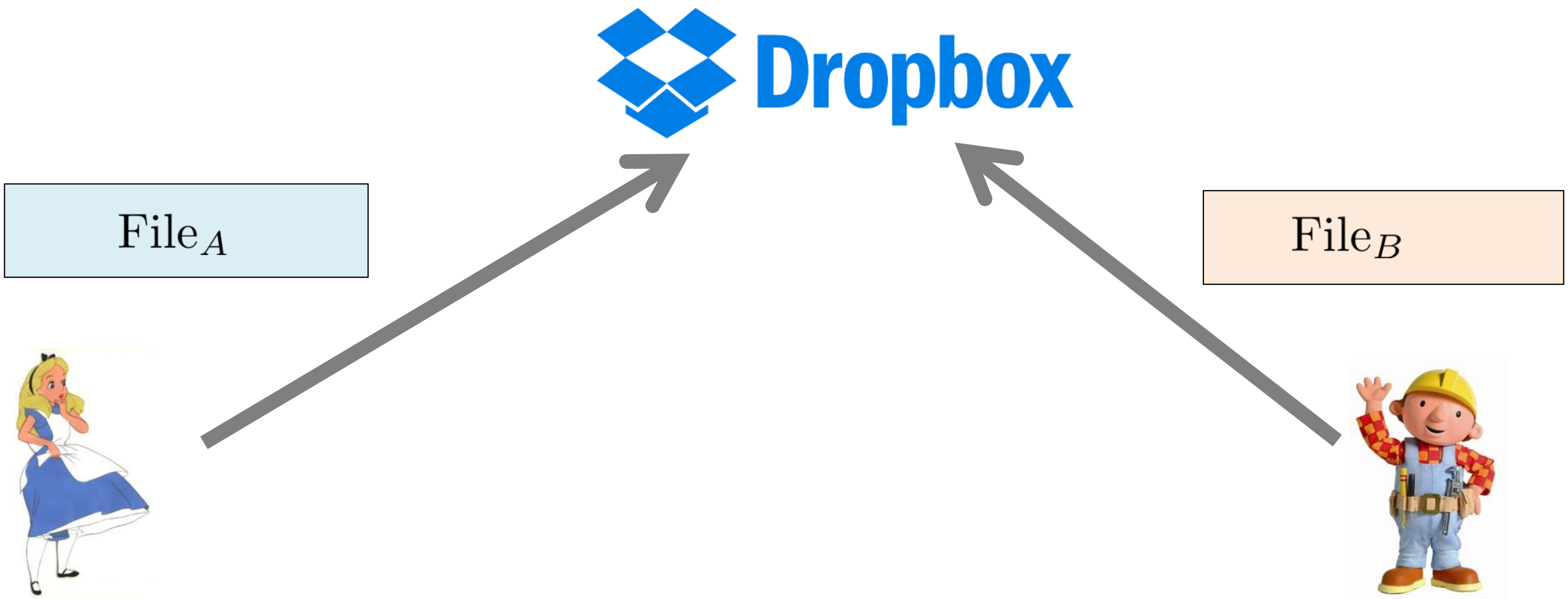
Agenda

1. Security Modeling for Hash Functions

2. Building Hash Function: MD Transform

3. Application: Password Storage

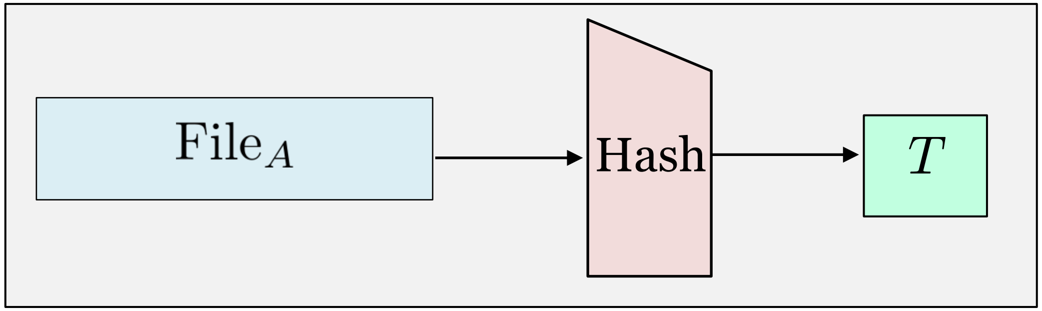
Motivating Application: Data Deduplication



Dropbox's goals:

- If many users store the same file, keep only a **single** copy
- Minimize bandwidth usage

Motivating Application: Data Deduplication



Dropbox

T

Check tags of existing files for duplication

Send / Not Send

File_A

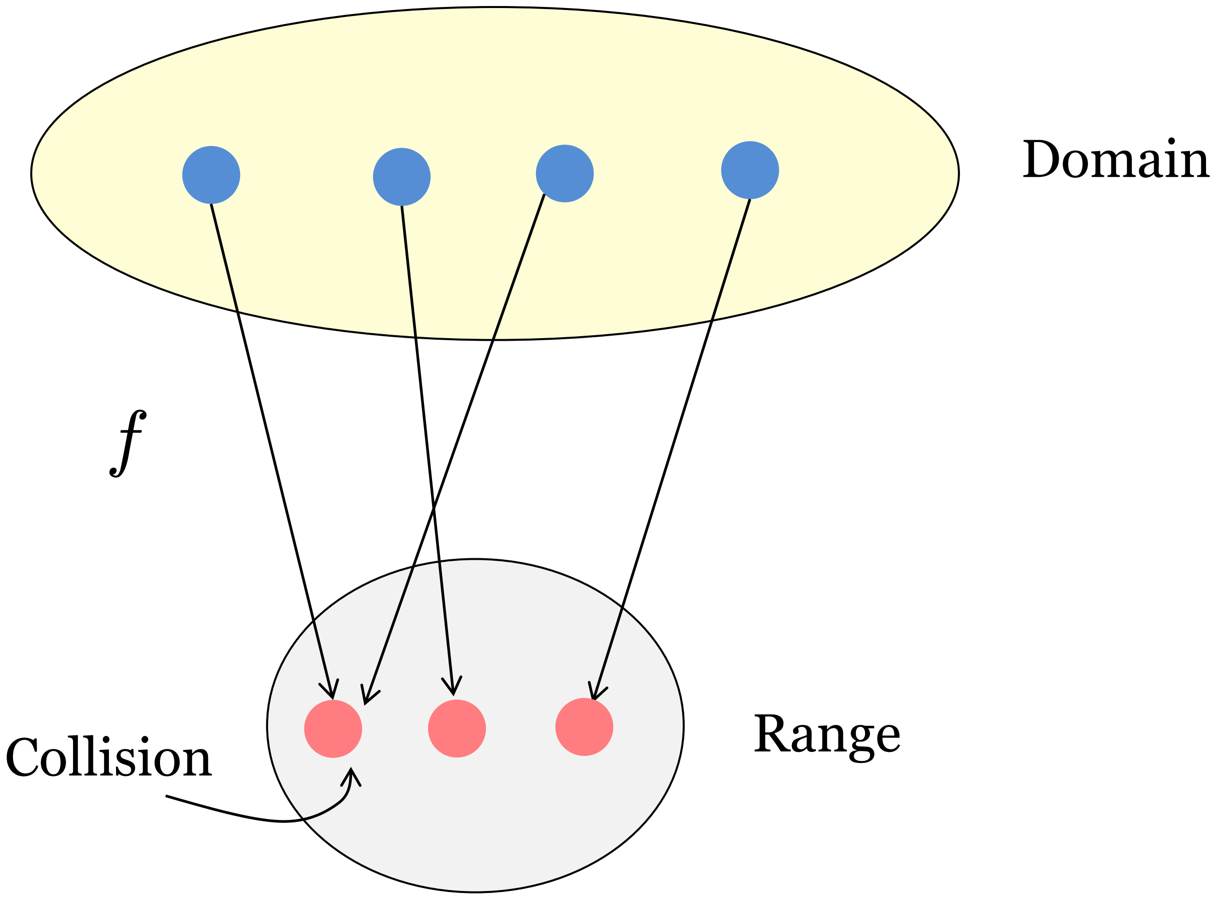
Send only when requested

What property

do we need for the hash?

Collision-Resistance

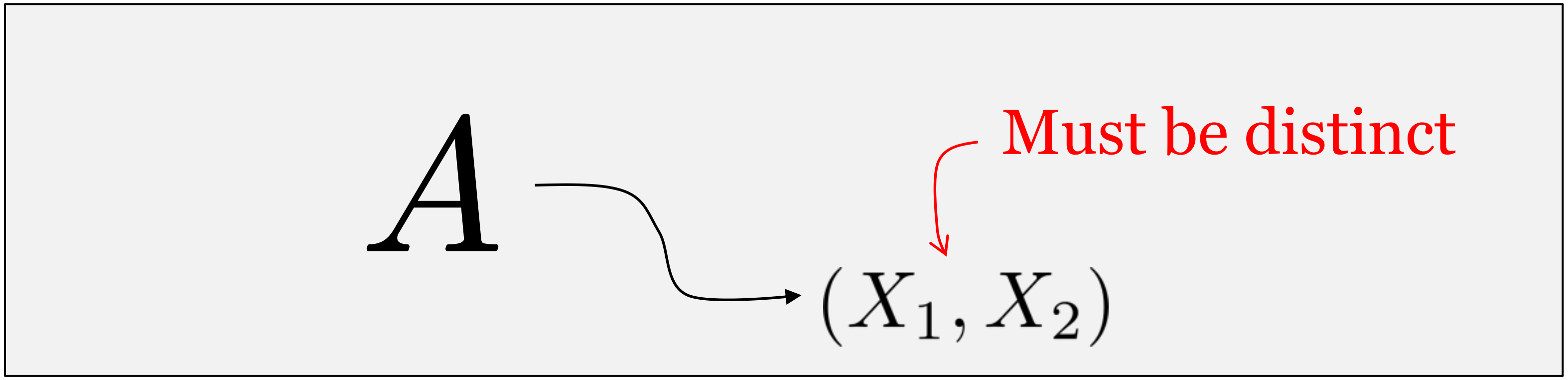
$f : \text{Domain} \rightarrow \text{Range}$



By Pigeonhole Principle, if $|\text{Domain}| > |\text{Range}|$ then collision exists

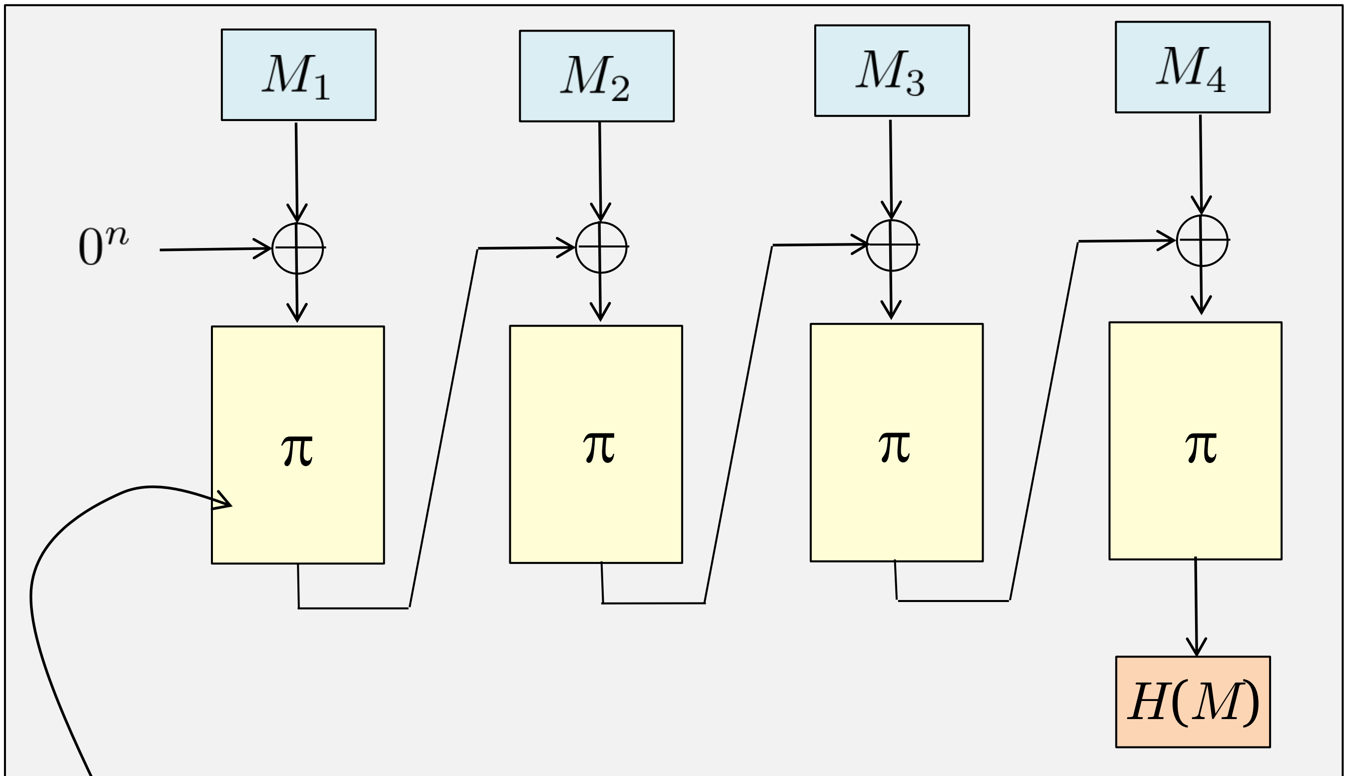
Want: collisions are hard to find, although they exist

Defining Collision-Resistance



$$\mathbf{Adv}_H^{\mathrm{cr}}(A) = \Pr[H(X_1) = H(X_2)]$$

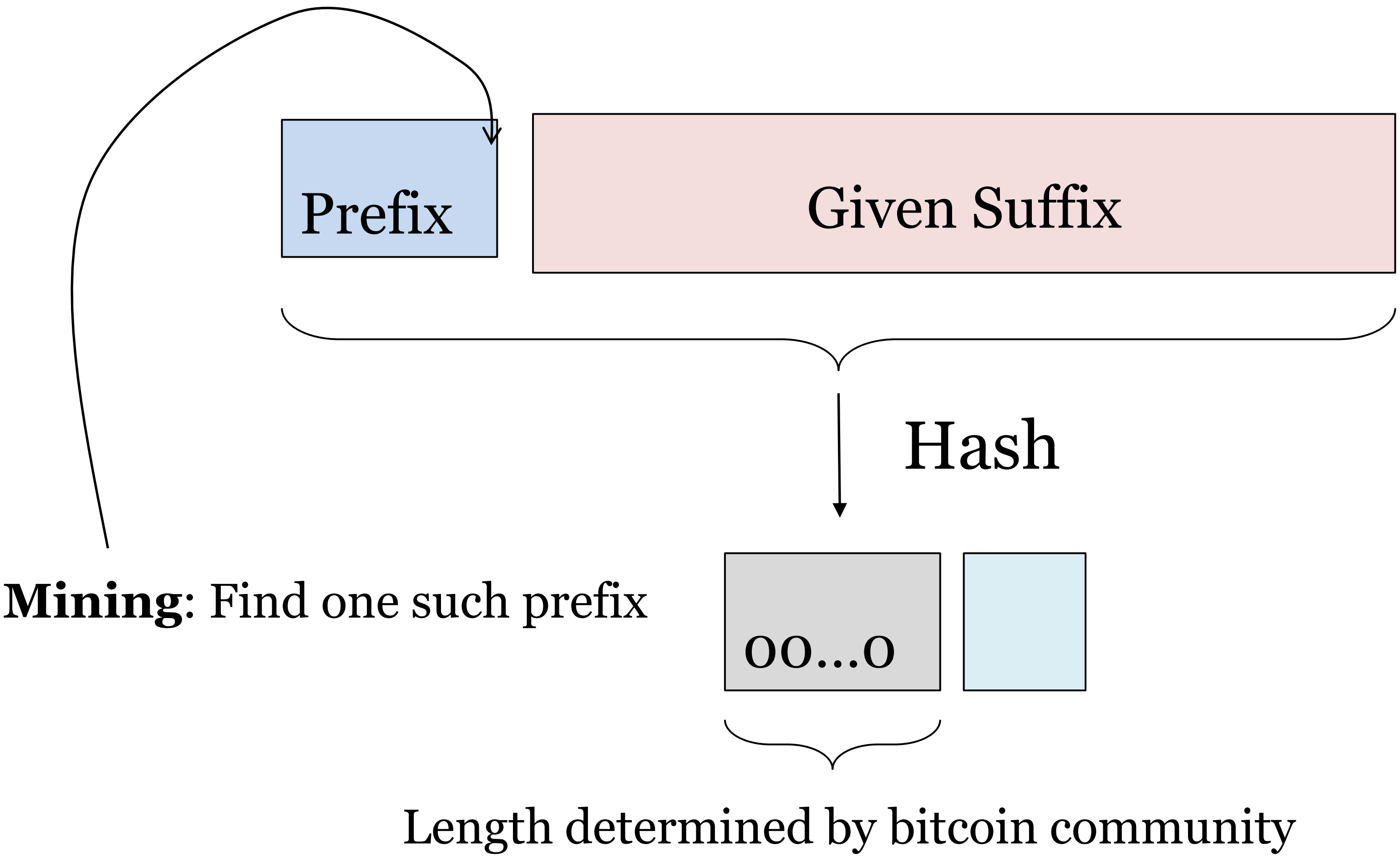
Exercise: Break Collision Resistance



Public permutation

π, π^{-1} are public

CR Is Not Enough: Bitcoin Mining



Want: Can't mine faster than brute-force

Modeling Security of Hash Functions

The Random Oracle Model

RO

Maintain a secret random $f : \{0,1\}^n \to \{0,1\}^m$

Everybody, including the adversary, has access to RO

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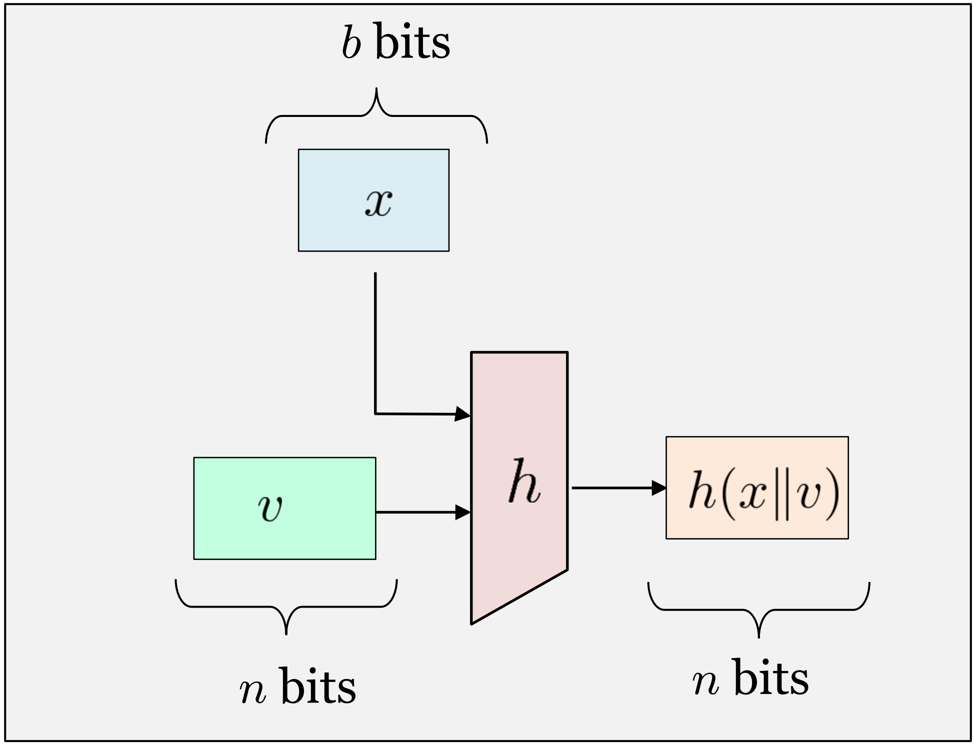
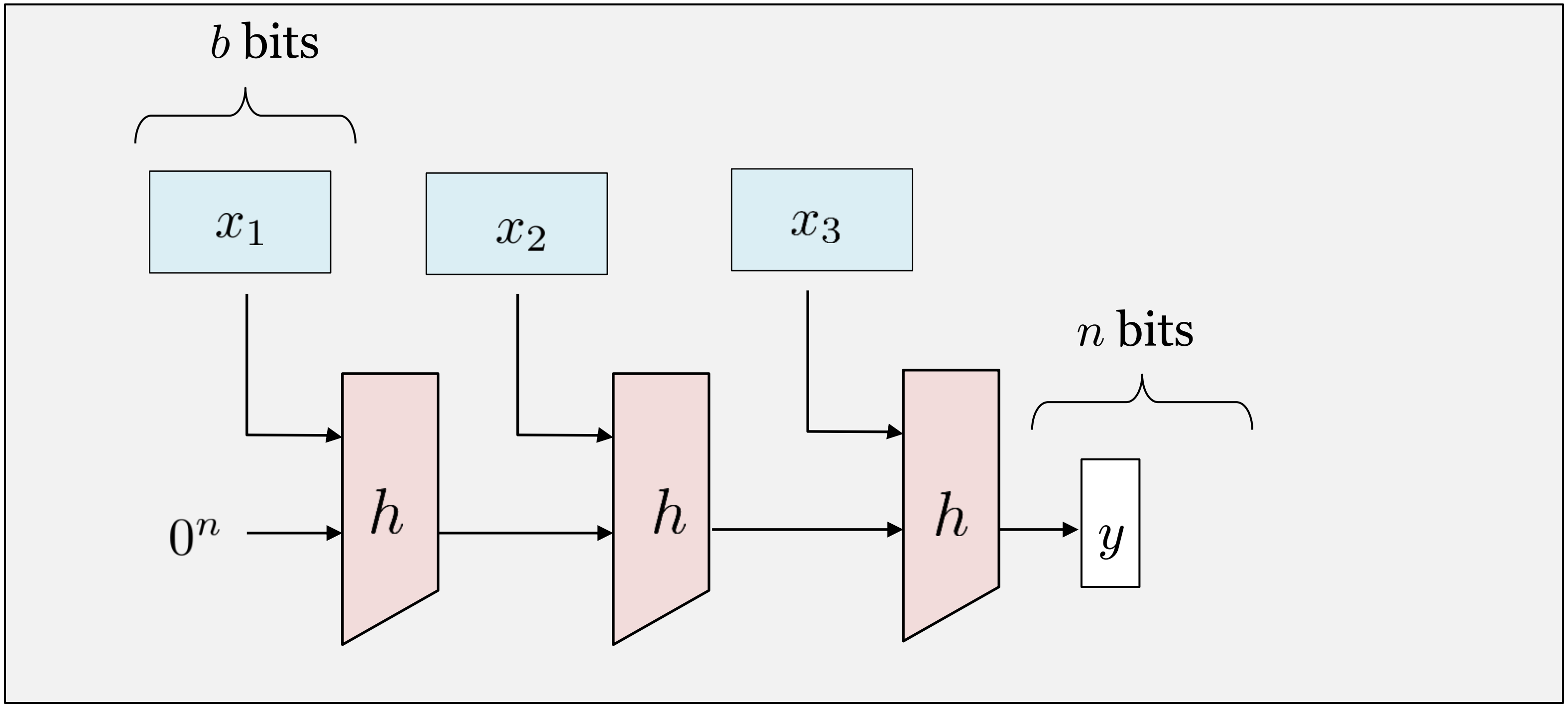
Compression Functions

$$h : \{0,1\}^{b+n} \rightarrow \{0,1\}^n$$



For SHA-2, b = 512 and n = 256

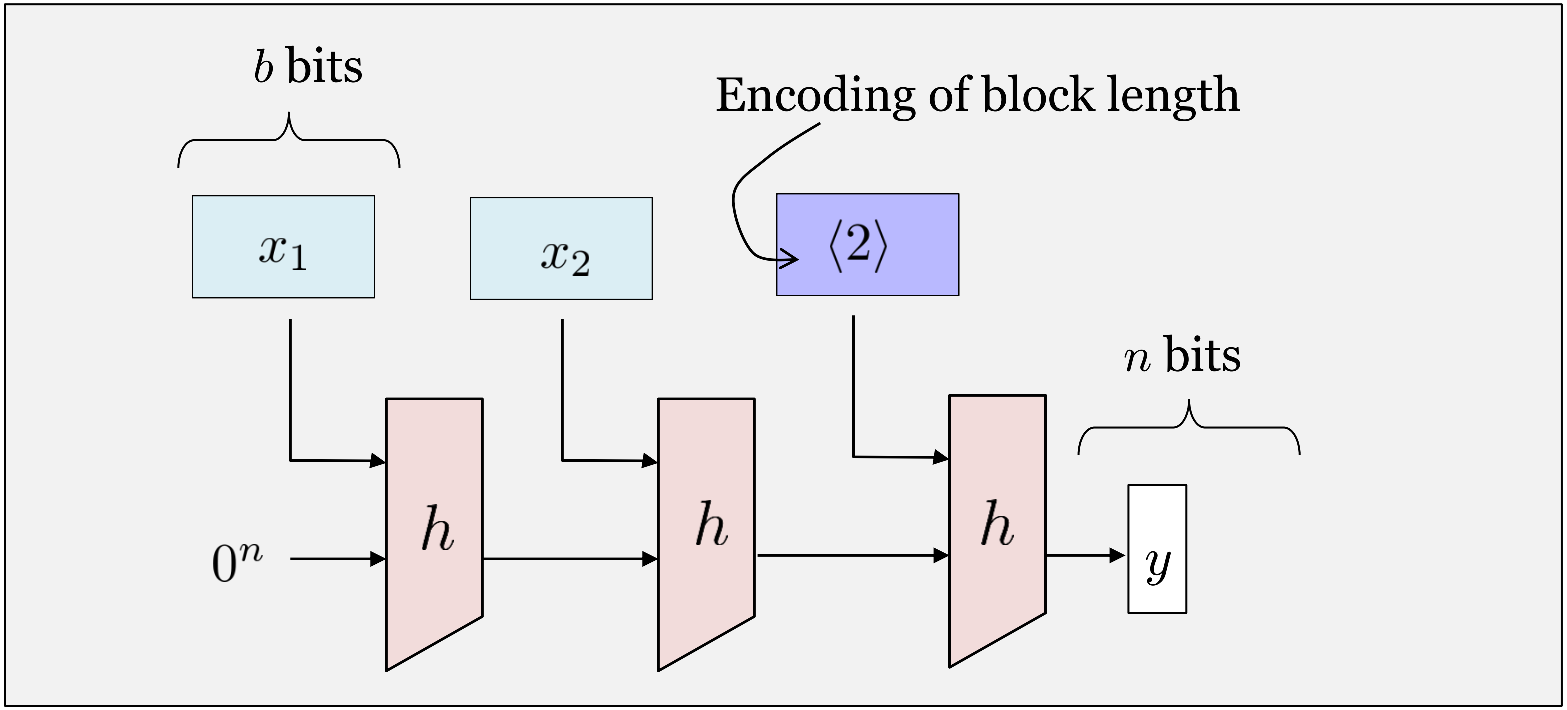
First Attempt



Question: Suppose that $h(0^b \| 0^n) = 0^n$

Break the collision resistance of H

Second Attempt: Plain Merkle-Damgard



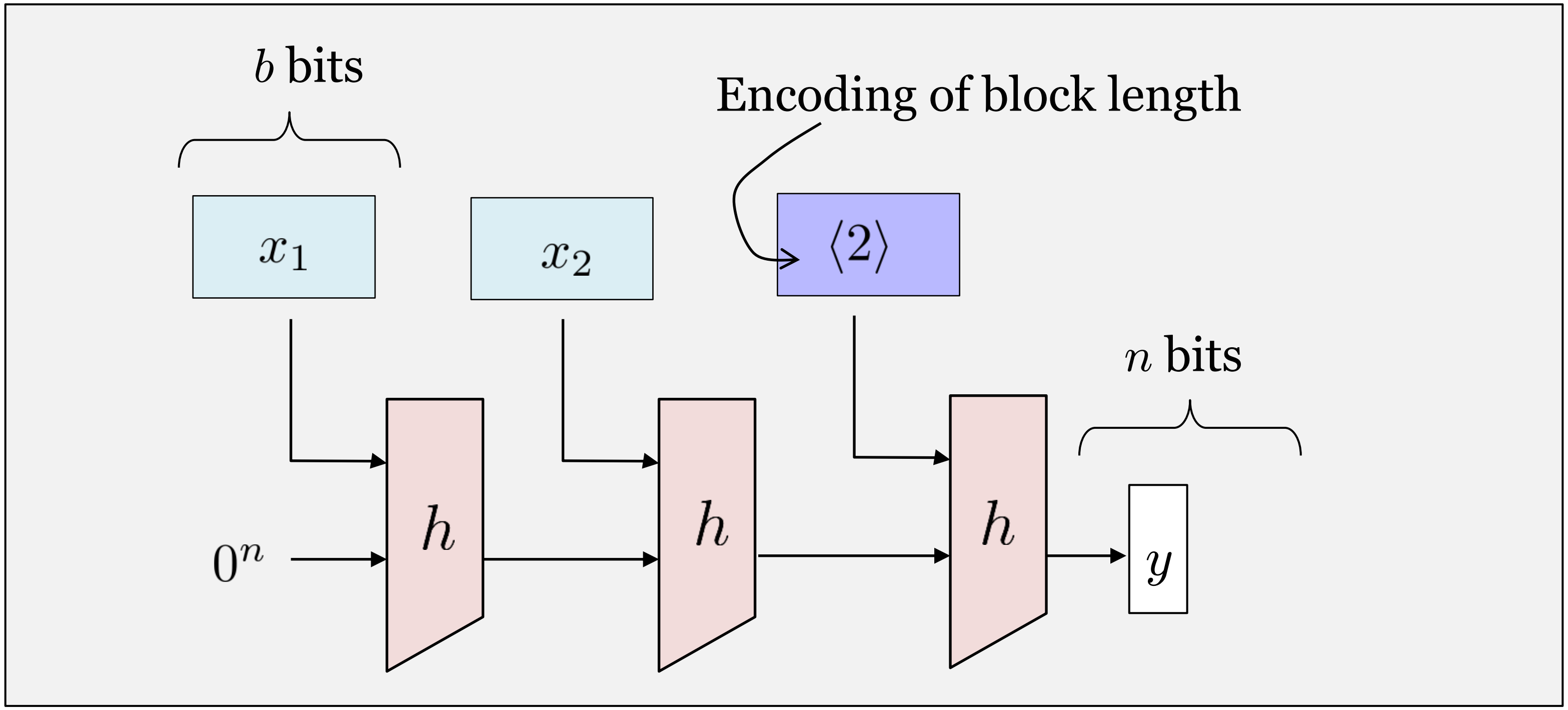
This is the structure of SHA-256

Theorem: If h is CR then H = **MD**(h) is also CR

Can't attack H if h has no weakness

Plain MD Is Not Enough for All Applications

Length-Extension Attack



Question: Consider the following MAC F

$$F_K(x) = H(K \| x)$$

Break the MAC security of F using a single Tag query

The Damage of Length Extension Attack

Hacking Trick: Bypass Authentication

K

Bank of America

bank.com

K

bank.com/api?token=ad6613c382&user=alice&cmd=NoOp

$H(K \| \text{"user=alice\&cmd=NoOp"})$

Adversary tricks Alice to perform a harmless command to learn an authentication token

The Damage of Length Extension Attack

Hacking Trick: Bypass Authentication

K

bank.com

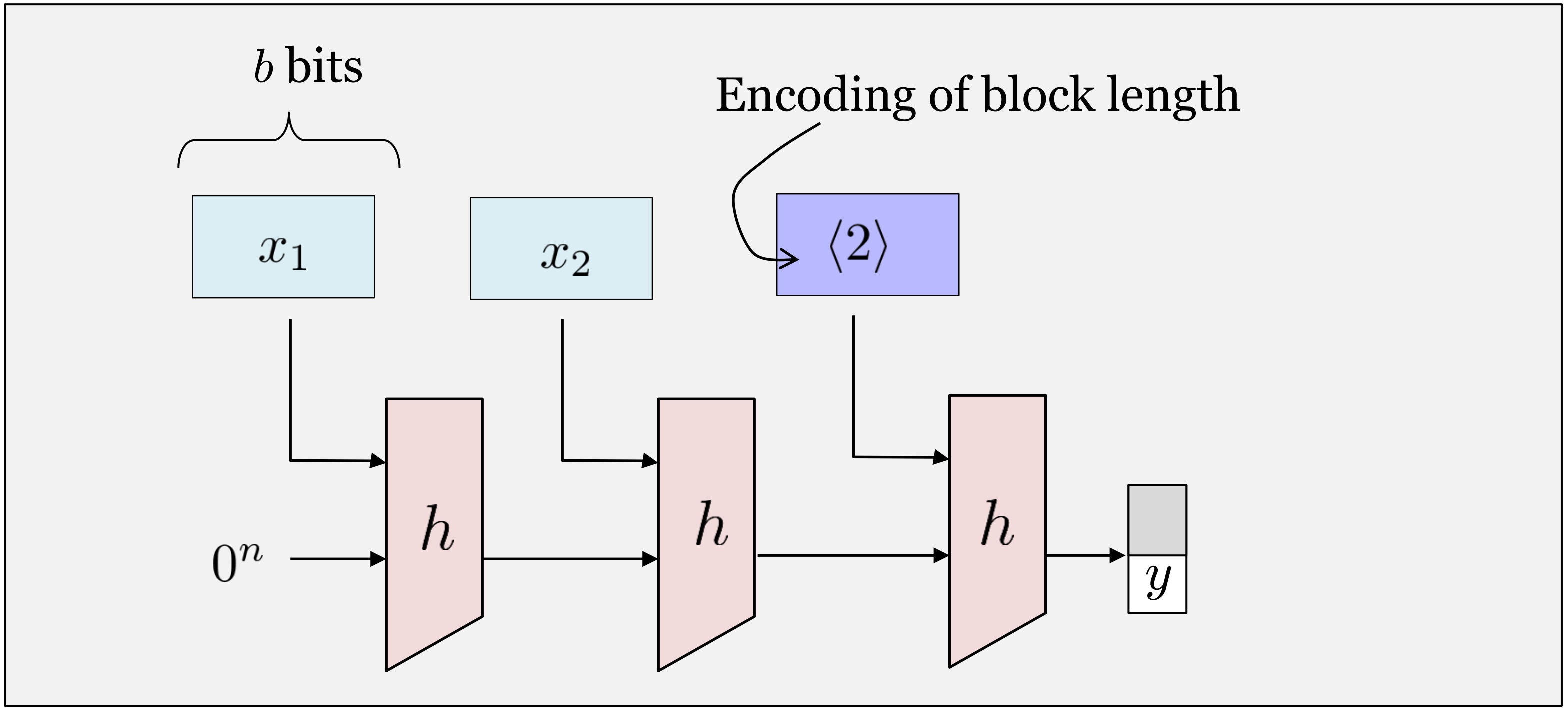
K

bank.com/api?token=dbb78b593f&user=alice&cmd=NoOp&cmd=OpenSafe

$$H(K \| \text{“user=alice\&cmd=NoOp\&cmd=OpenSafe”})$$

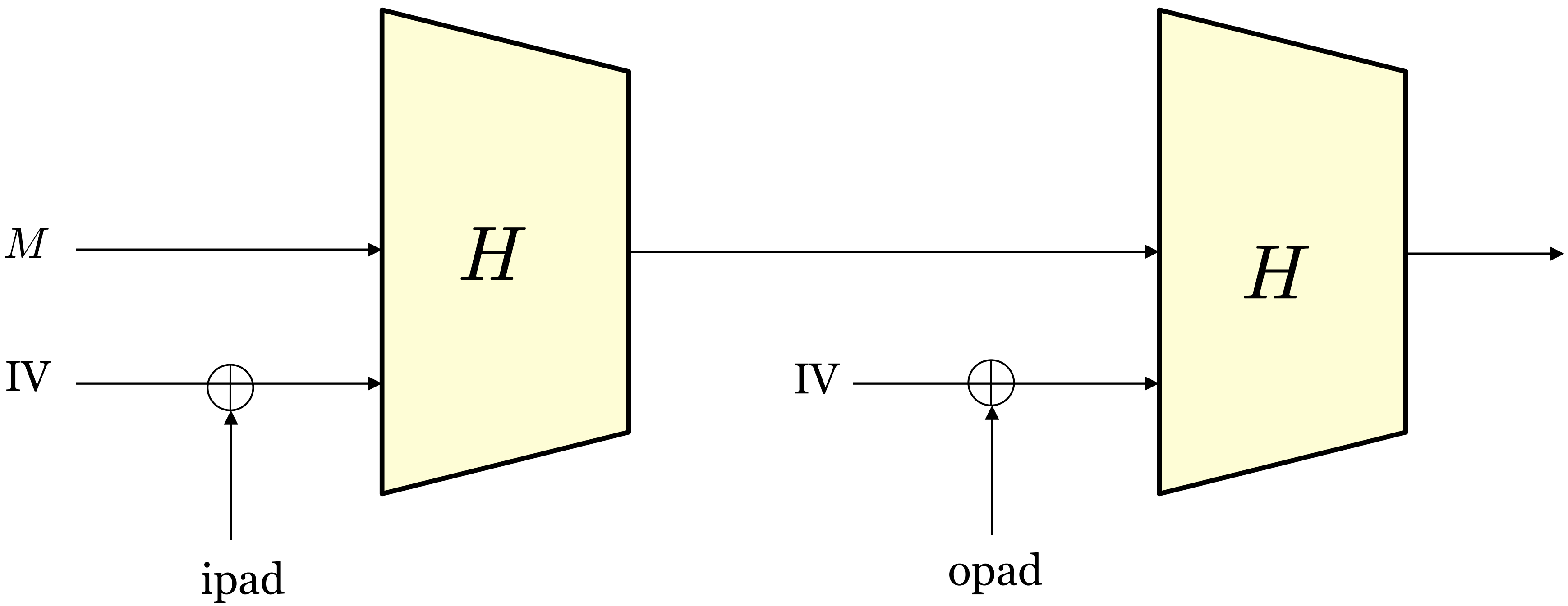
Adversary can compute the authentication token for a damaging command

The (Strengthened) MD Transform



The output needs to be truncated

How To Have Large Output: HMAC



On large input, HMAC is only a bit more expensive than SHA-256

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3. **Application: Password Storage**

Password Storage

MOTHERBOARD
TECH BY VICE

T-Mobile Stores Part of Customers' Passwords In Plaintext, Says It Has 'Amazingly Good' Security

A T-Mobile Austria customer represe

admission in a Twitter thread.

ars TECHNICA

BIZ & IT —

How an epic blunder by Adobe could strengthen hand of password crackers

Engineers flout universal taboo by *encrypting* 130 million pilfered passwords.

NEWS

Hackers crack more than 60% of breached LinkedIn passwords

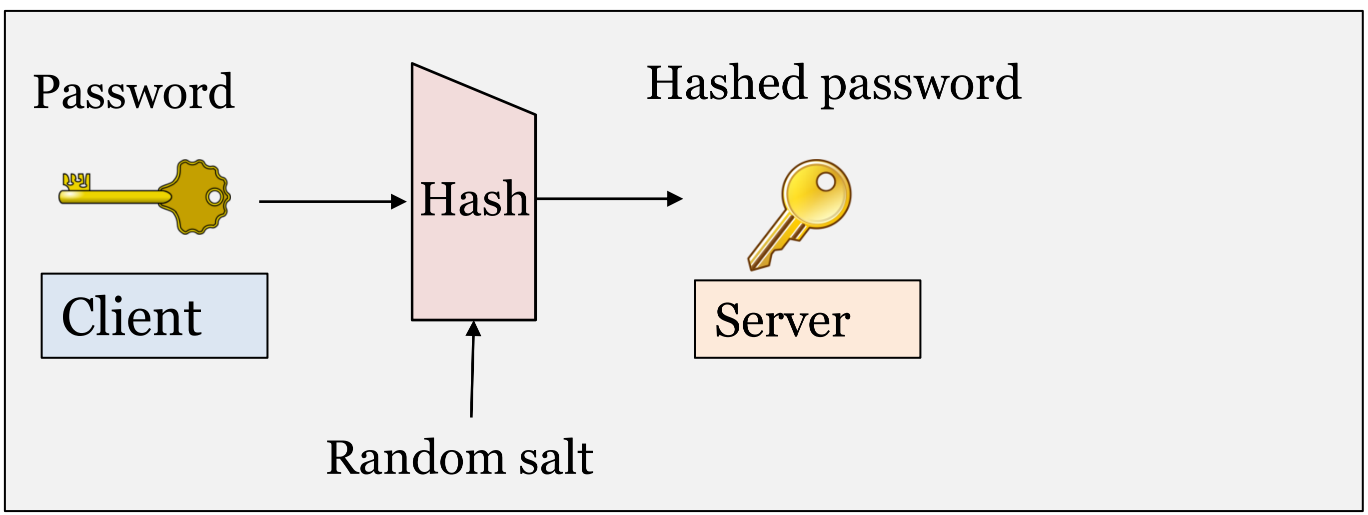
Speed of hackers to crack passwords shows weakness of security scheme used by LinkedIn, researchers say

How Should Servers Store Users' Passwords?

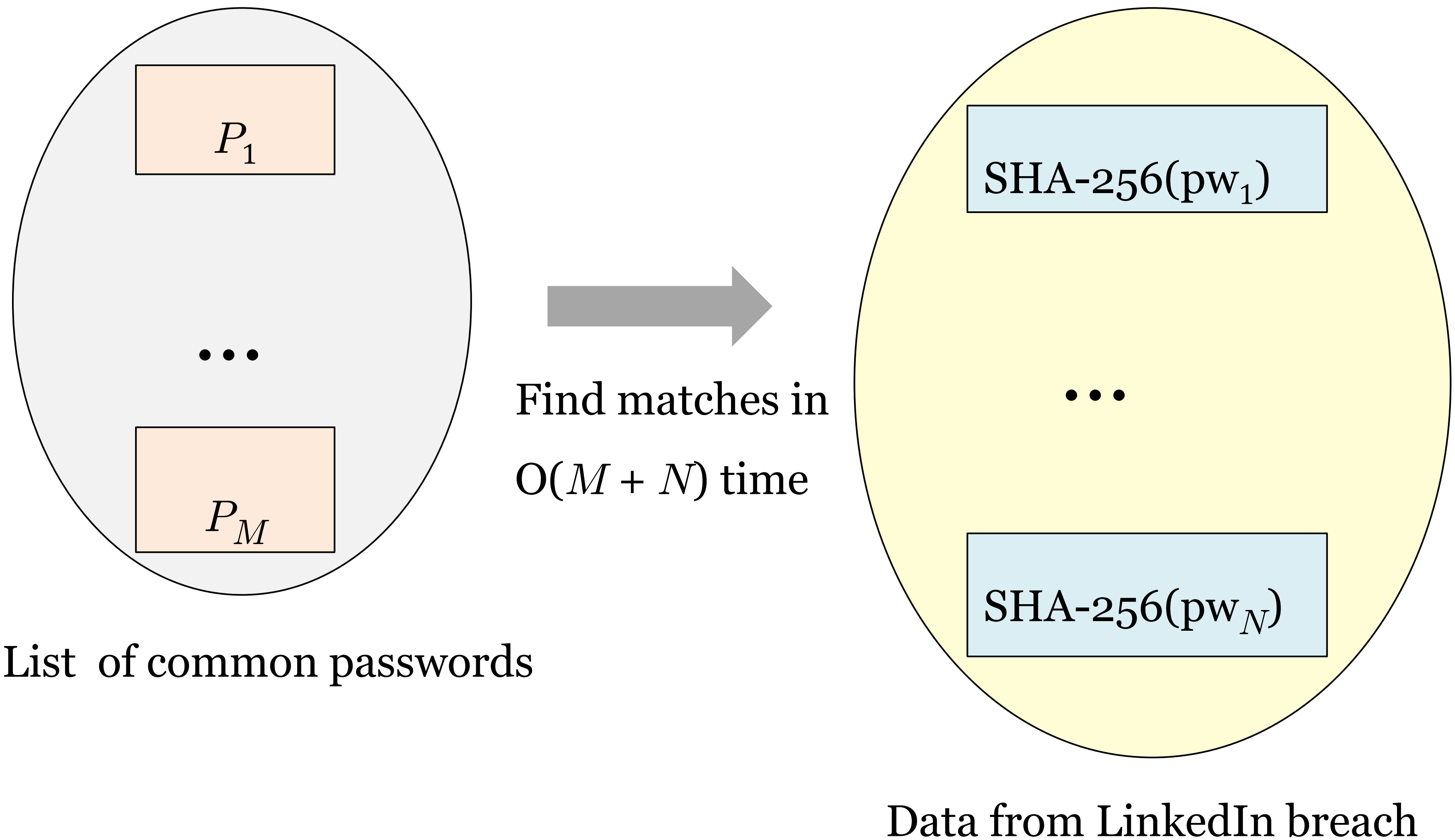
Rule 1: Only store hash outputs of passwords

Even server can't recover the passwords

Rule 2: Use a random salt for each user

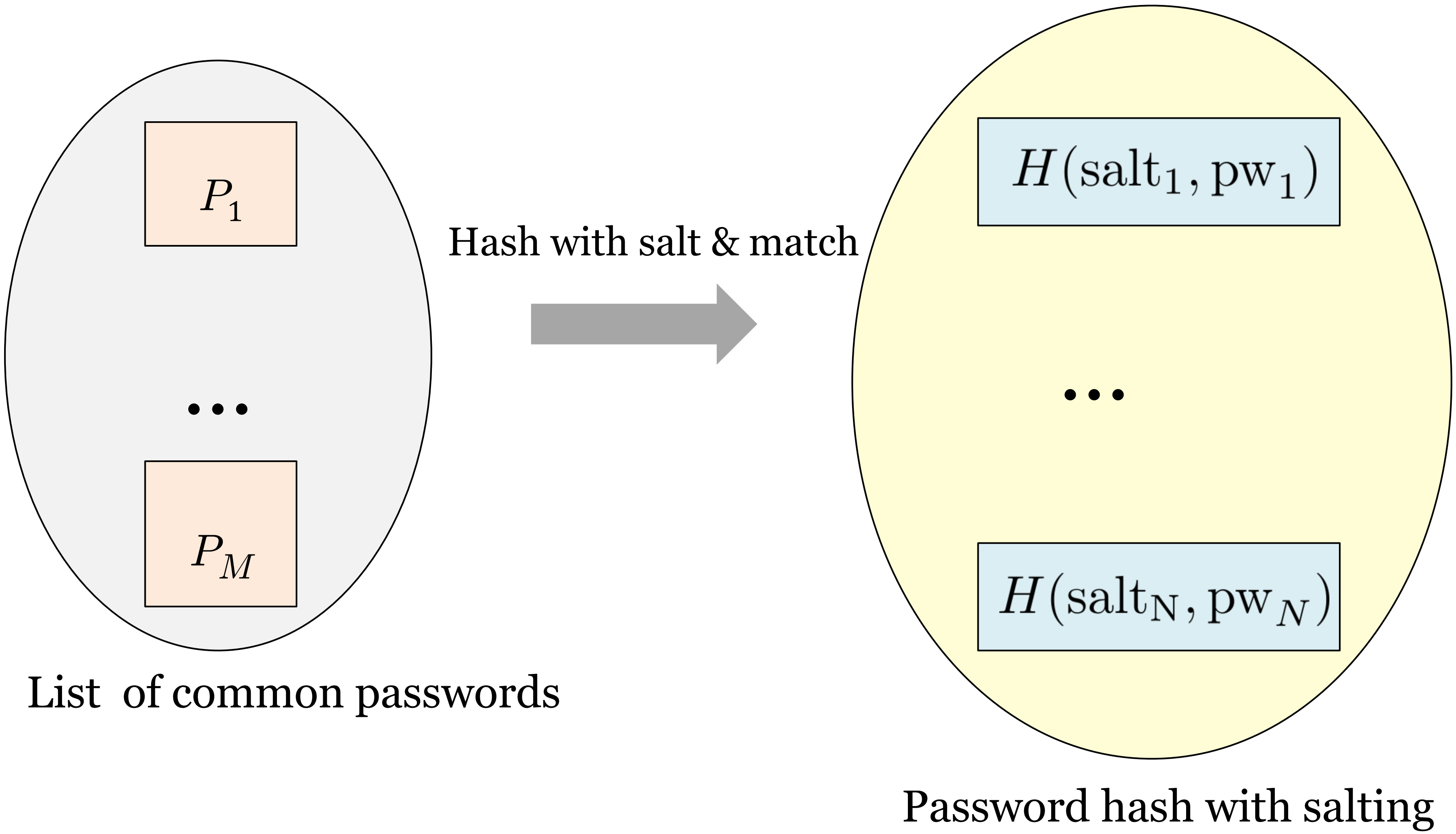


Why Salts: Dictionary Attacks



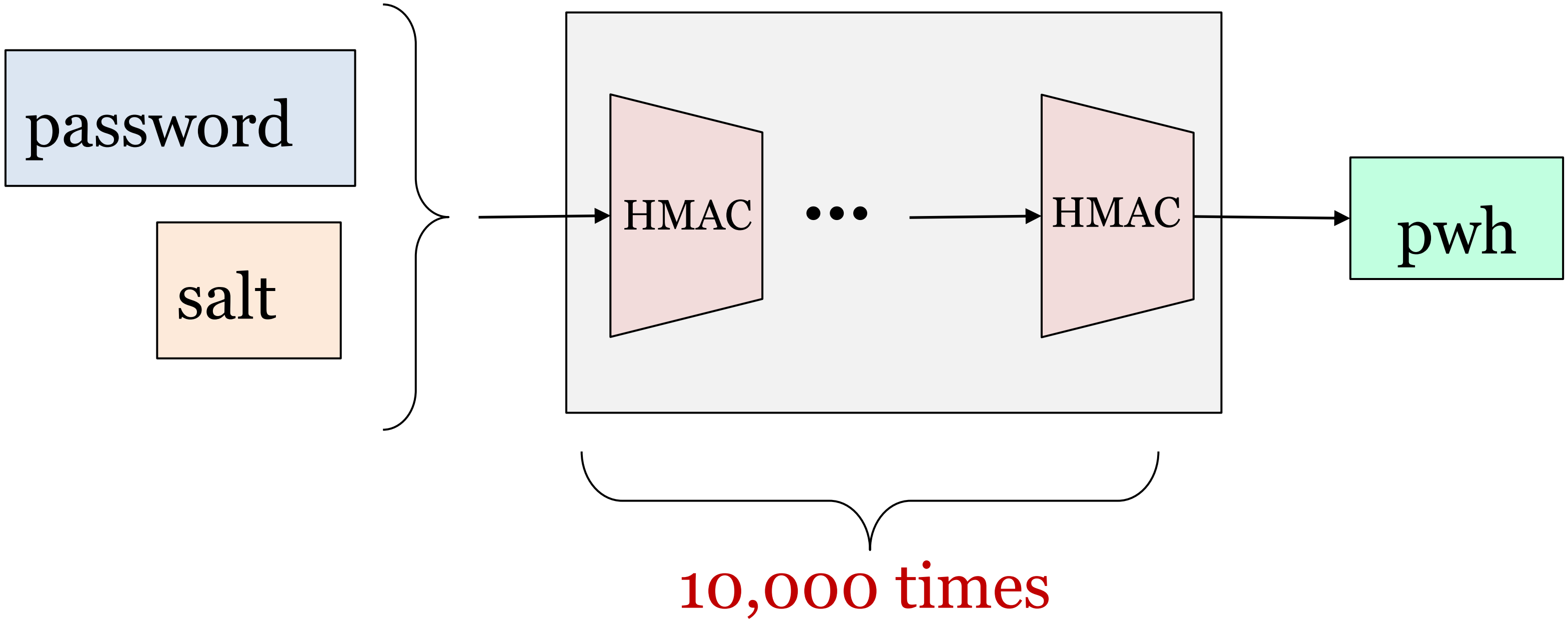
Cost of Dictionary Attacks on Salting

Need $\Theta(Mq)$ calls to H to recover q passwords



Make It Even More Expensive

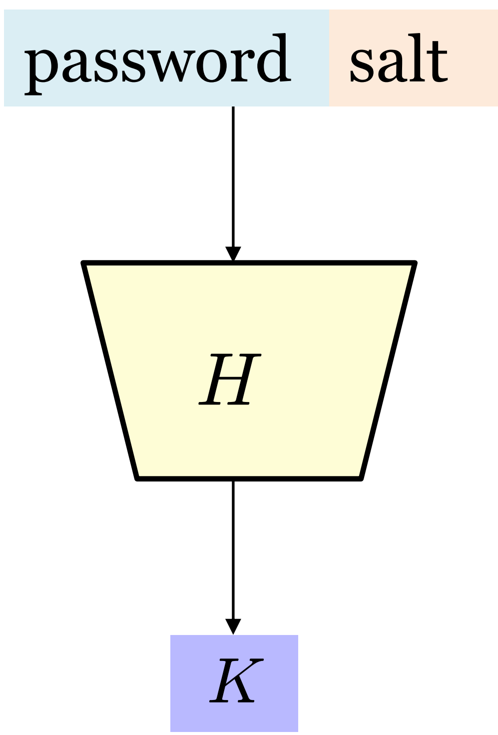
Deliberately Slow Hashing



- Makes no difference for human users.
- Increase the cos of attackers for 10,000 times

Password-Based Encryption

$\mathcal{E}_K(M)$



- Commonly used in practice
- In Hw4, we'll analyze its security